

TOMORROW starts here.



Cisco *live!*

Evolution and Challenges of Data Centre Network and Host-Based Overlays

BRKDCT-2328

Victor Moreno

Distinguished Engineer

Agenda

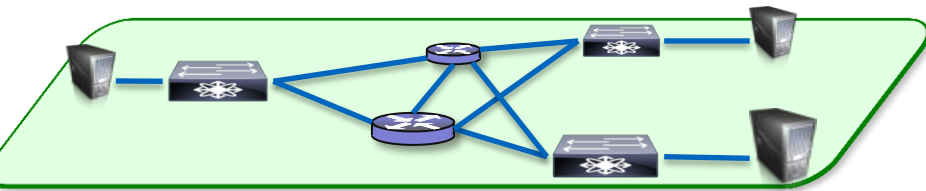
- Overlay Foundational Principles and evolution
- Mapping overlay technologies to the network
- The role of the underlay
- Management and orchestration



Foundational Principles of Network Overlays

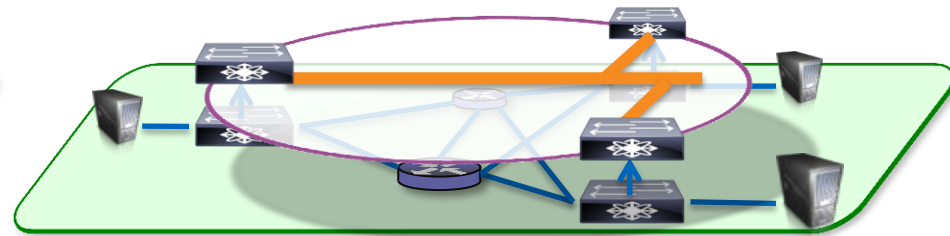
Why Overlays?

Seek well integrated best in class Overlays and Underlays



Robust Underlay/Fabric

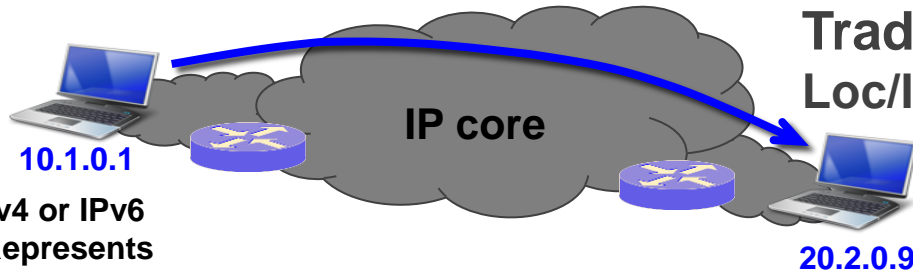
- High Capacity Resilient Fabric
- Intelligent Packet Handling
- Programmable & Manageable



Flexible Overlay Virtual Network

- Mobility – Track end-point attach at edges
- Scale – Reduce core state
 - Distribute and partition state to network edge
- Flexibility/Programmability
 - Reduced number of touch points

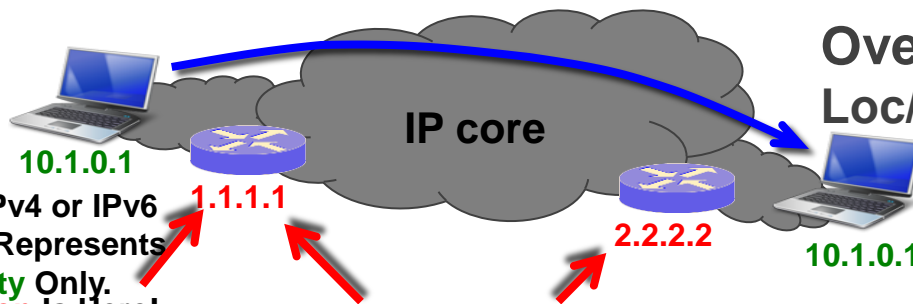
Seminal Idea: Location and Identity Separation



Device IPv4 or IPv6 Address Represents Identity and Location

Traditional Behaviour Loc/ID “Overloaded” Semantic

When the Device Moves, It Gets a New IPv4 or IPv6 Address for Its New Identity and Location



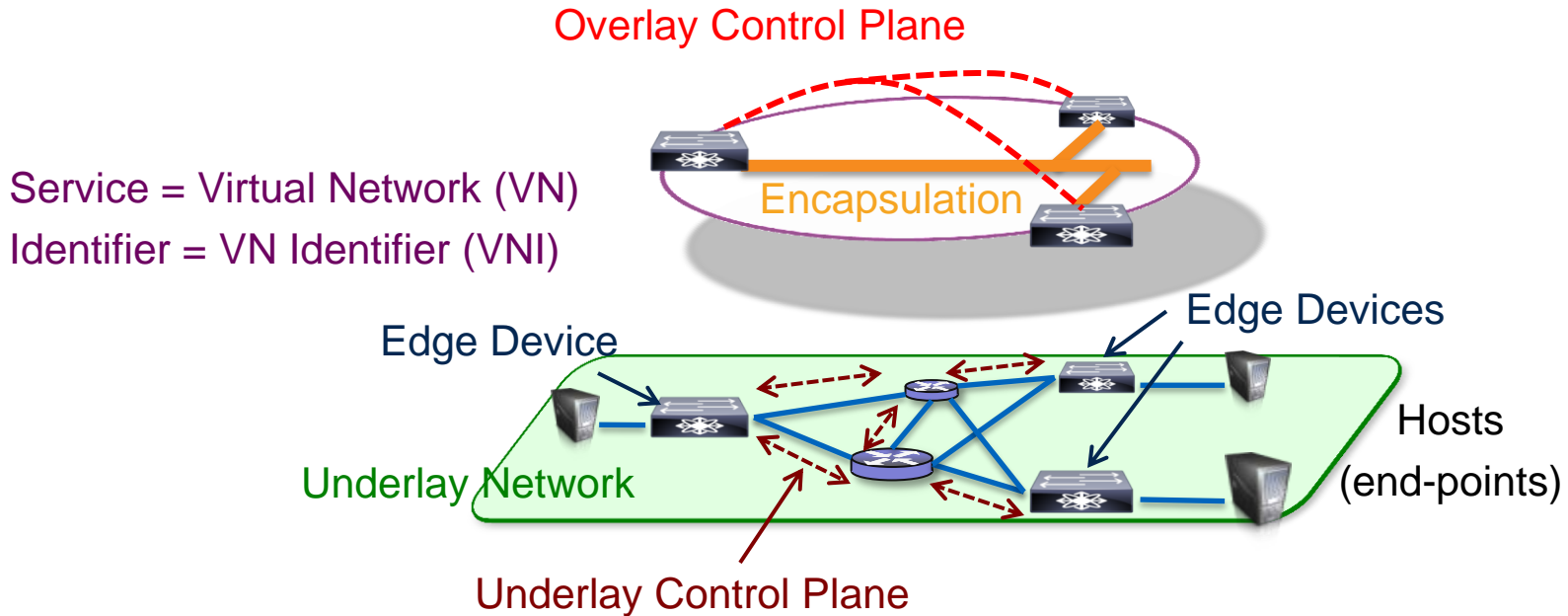
Device IPv4 or IPv6 Address Represents Identity Only. Its Location Is Here!

Only the Location Changes

Overlay Behaviour Loc/ID “Split”

When the Device Moves, Keeps Its IPv4 or IPv6 Address. It Has the Same Identity

Overlay Taxonomy



Overlay Attributes

Service

Layer 2 Service

Layer 3 Service

Edge Device

Host Overlays

Network Overlays

Signalling

Data Plane Learning

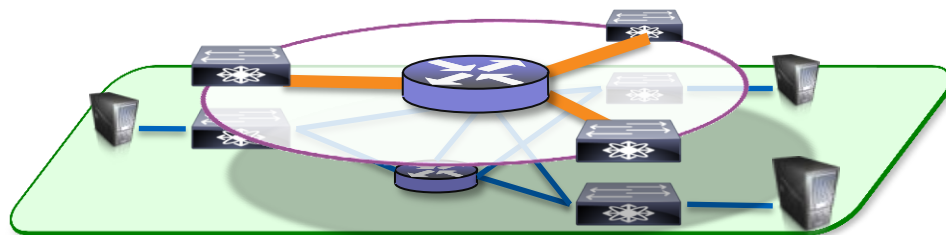
Control Plane Learning

Overlay Service Type Evolution

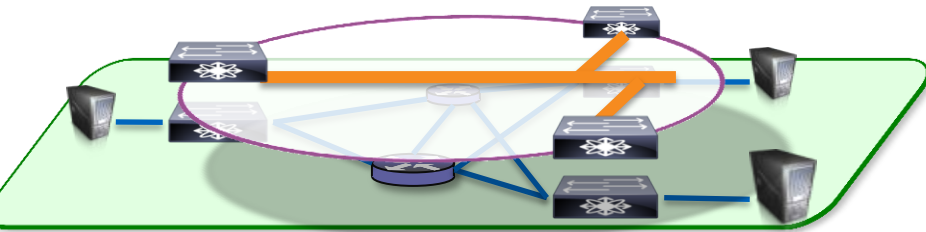
Service

Layer 2 Service

Layer 3 Service

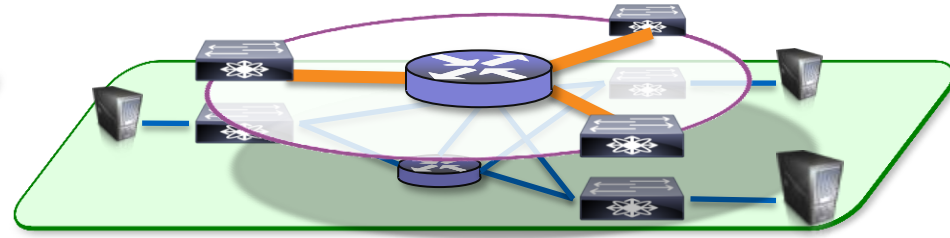


Types of Overlay Service



Layer 2 Overlays

- Emulate a LAN segment
- Transport Ethernet Frames (IP and non-IP)
- Single subnet mobility (L2 domain)
- Exposure to open L2 flooding
- Useful in emulating physical topologies



Layer 3 Overlays

- Abstract IP based connectivity
- Transport IP Packets
- Full mobility regardless of subnets
- Contain network related failures (floods)
- Useful in abstracting connectivity and policy

Hybrid L2/L3 Overlays offer the best of both domains

Layer 2 Overlay Considerations

- **Scale** of the edge devices
 - L2 addresses in Ethernet (MACs) use a flat space which cannot be summarised
- **L2/L3 boundary** scaling
 - Large L2 domains require a large capacity L3 gateway to handle large ARP and MAC tables at a frequent rate of refresh
- **Multi-homing** sites can induce loops in the network
- **Flooding** of L2 protocols, unknown unicasts and broadcast in general can propagate failures across the entire L2 domain

Solved with ...



Layer 3 Overlays



Layer 3 Overlays



Network Overlays



MAC routing

Multi-homing in L2 Overlays

Source learning assumes single attached sites
But network overlays involve edge resiliency

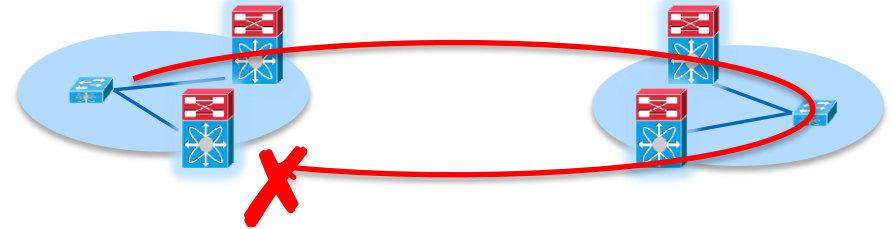
Enhancements are required to address:

- Loop resolution
- Multi-pathing
- Broadcast/Multicast de-duplication

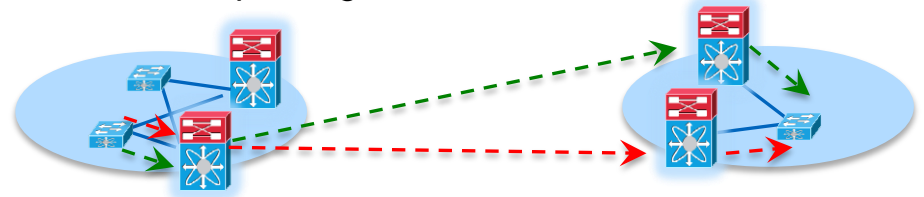
Two Approaches:

- Active-Standby (Data Plane or Control Plane)
 - One active device per VLAN (single attached site)
 - VLAN based load balancing
- Active-Active (Control Plane only)
 - One active device for multi-destination traffic
 - Intra-VLAN load balancing for unicast

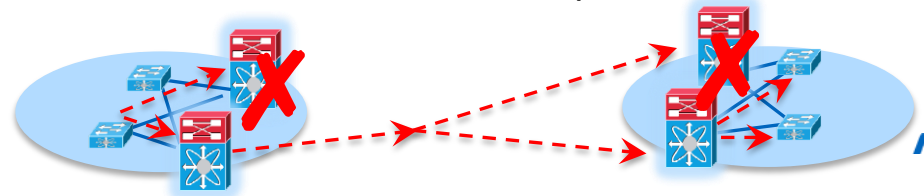
- Loop resolution



- Multi-pathing



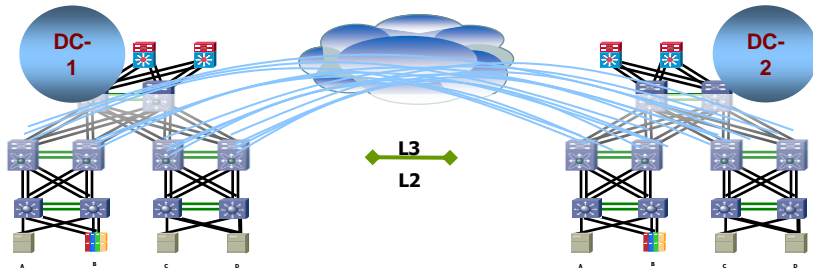
- Broadcast/Multicast de-duplication



Flooding in L2 Overlays

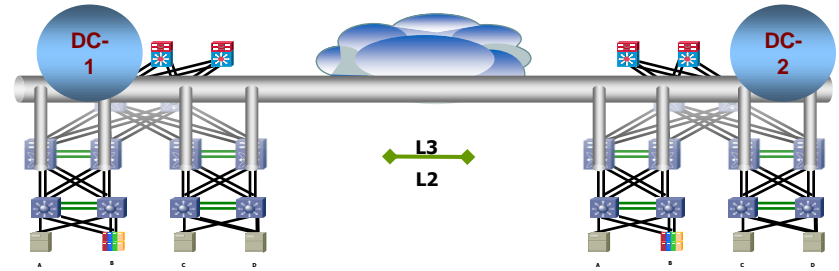
Control Plane Signalling eliminates the need for floods

Data Plane Learning



- Pre-set flood facility
- MAC learning based on flooding
- Flood L2 protocols and unknown unicast
 - ➔ Failure propagation
- Fail Open
- Suitable for small domains (failure scope)

Control Protocol



- No predetermined flood tree
- MAC learning by control protocol
 - ➔ Contain Failures and L2 protocols
 - ➔ Rich information
- Fail Closed
- Better suited for broad scope

Flooded L2 Overlays

MAC Routing

L2 Overlay Evolution

Inter-DC (DCI)

VPLS → A-VPLS
Flood and Learn



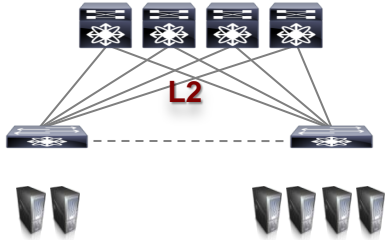
OTV
MAC routing (IS-IS)



EVPN / L2-LISP
MAC routing (BGP/LISP)

Backbone Network

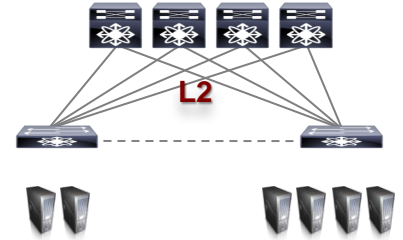
Intra-DC
(Fabric)



FP → TRILL
Flood and Learn



Enhanced Forwarding (hybrid)



Host Overlays

VXLAN
Flood and Learn



Control Plane

Layer 3 Overlay Considerations

- **Scale** of the edge devices
 - Can be improved further by using an on-demand pull model
- **IP Mobility** for subnet disaggregation
 - Members of a subnet may be distributed across locations
 - Any host anywhere
- **Broadcast & Link-local multicast** traffic to be handled as a special case
 - Potentially without even learning MAC addresses



Addressed with ...

On-demand Pull



Layer 2 Semantics
with IP routing



Combined L2/L3
overlay

L3 Overlay Evolution

Edge Device Scale

Push Protocol Model

- IP/BGP MPLS VPNs are highly scalable today
- PE routers must:
 - Hold a large number of prefixes
 - Maintain multiple routing protocol adjacencies
- Mobility and cloud will add pressure in terms of:
 - Prefix granularity and volume
 - Increased number of PEs

Pull Protocol (on-demand) Model

- LISP deployments and footprint are increasing rapidly
- On-demand caching models ease the requirements on the edge devices:
 - Only prefixes being utilised are cached
 - No routing adjacencies are maintained
- A pull model is expected to provide global scalability to enable pervasive cloud models

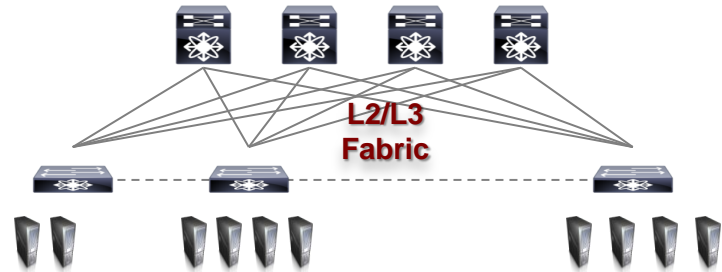
IP Mobility with L3 Overlays

- Granular location information
 - Allow subnet members to move anywhere
- Layer 2 semantics
 - ARP proxy
 - Consistent default Gateway presence
- L3 at the Access
 - Access switch replies to all ARPs with the same MAC address
 - Host routing for all traffic within the fabric
 - Summary prefix outside the fabric

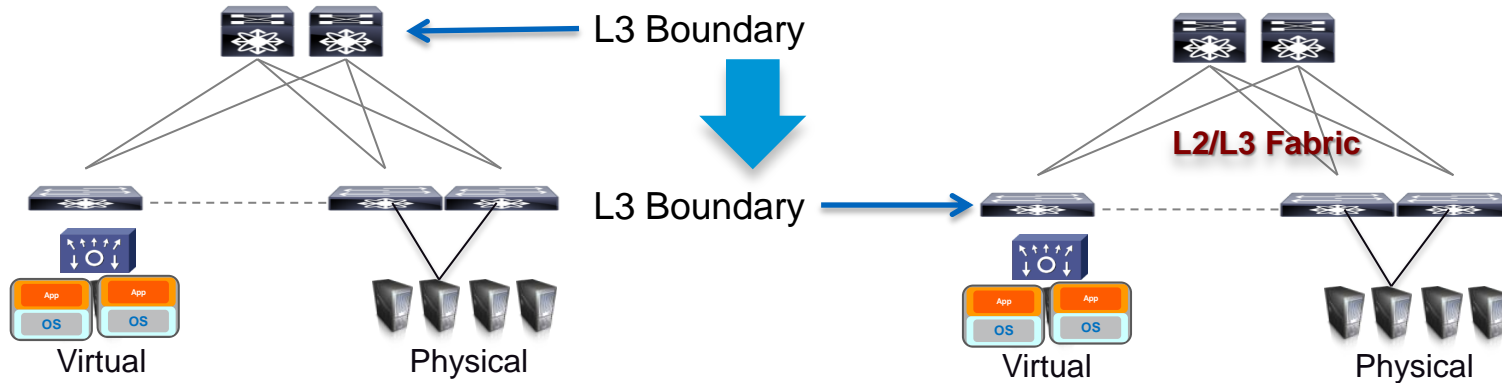


Combined L2/L3 Overlays

- Route all IP traffic including Intra-subnet
- Bridge only broadcast and link-local multicast traffic
- Assumption is that most traffic is IP (if not all)
- Bridge any non-IP traffic present
 - If only handling multi-destination non-IP traffic then no MAC address learning is required
 - If handling non-IP unicast traffic, then MAC address learning is required



Distributed Gateway Function in L2/L3 Overlays



Traditional L2 - centralised L2/L3 boundary

- Always bridge, route only at an aggregation point
- Large amounts of state converge
- Scale problem for large# of L2 segments
- Traditional L2 and L2 overlays

L2/L3 fabric (or overlay)

- Always route (at the leaves), bridge when necessary
- Distribute and disaggregate necessary state
- Optimal scalability
- Enhanced forwarding and L3 overlays

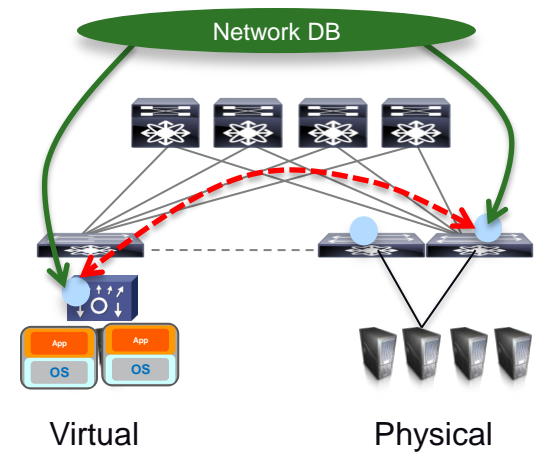
Overlay Edge Device and Data Plane Evolution

Service

Edge Device

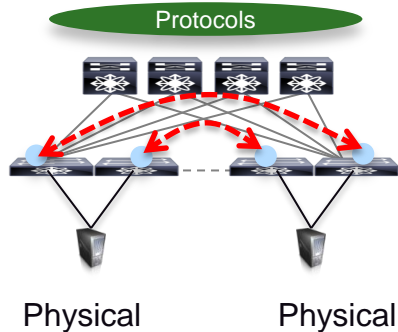
Layer 2 Service
Layer 3 Service

Host Overlays
Network Overlays



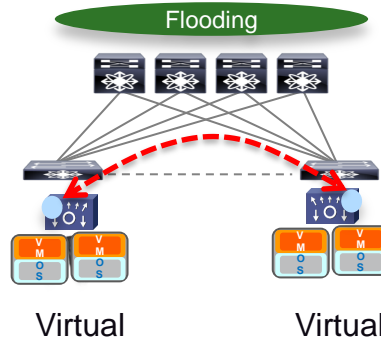
Overlay Network Evolution: Edge Devices

Network Overlays



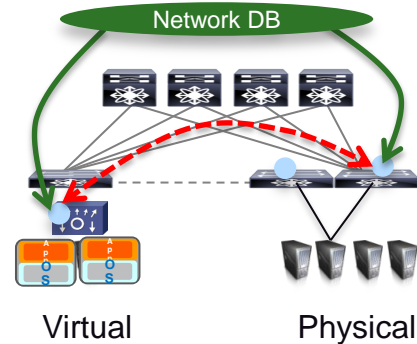
- Router/switch end-points
- Protocols for resiliency/loops
- Traditional VPNs
- OTV, VPLS, LISP, FP

Host Overlays




- Virtual end-points only
- Single admin domain
- VXLAN, NVGRE, STT

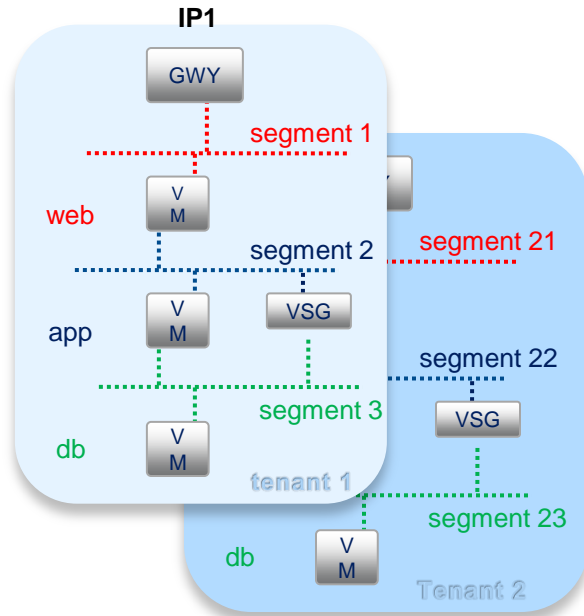
Hybrid Overlays



- Physical and Virtual
- Resiliency + Scale
- x-organisations/federation
- Open Standards

 Tunnel End-points

Host Overlays



Elastic creation of virtual Segments

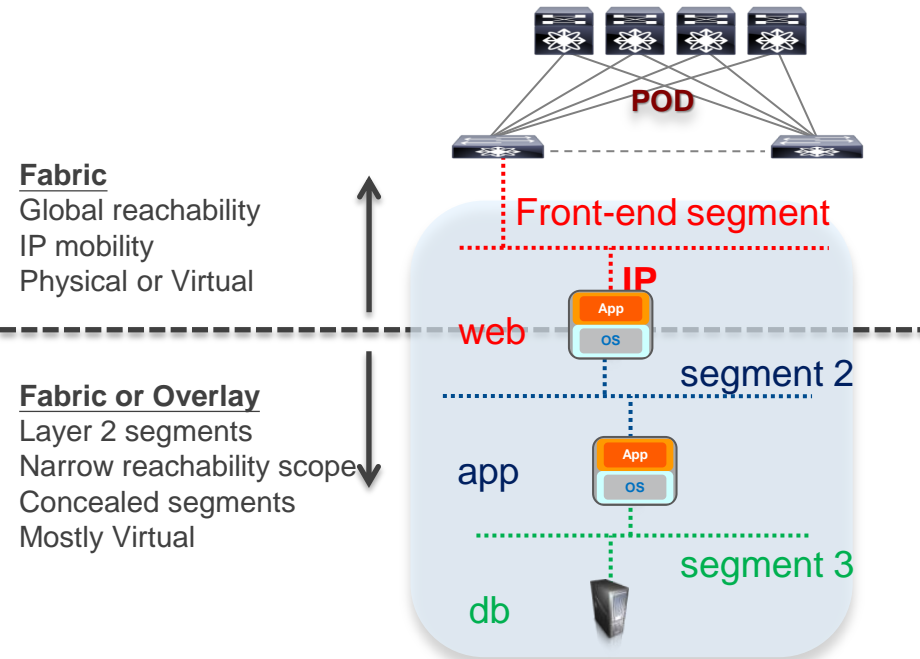
- Mobile: Can be instantiated anywhere
 - Move along with VMs as necessary
- Very large number of segments
 - Do not consume resources in the network core
- Isolated, not reachable from the IP network
 - Front-end segment must be handled by the fabric
- Host overlays are initiated at the hypervisor virtual switch → Virtual hosts only
- GWY to connect to the non-virtualised world
- Variants: VXLAN, NVGRE, STT

Multi-tier Virtual App = VMs + vSegments + GWY

Application: Cloud Services

Segmentation:

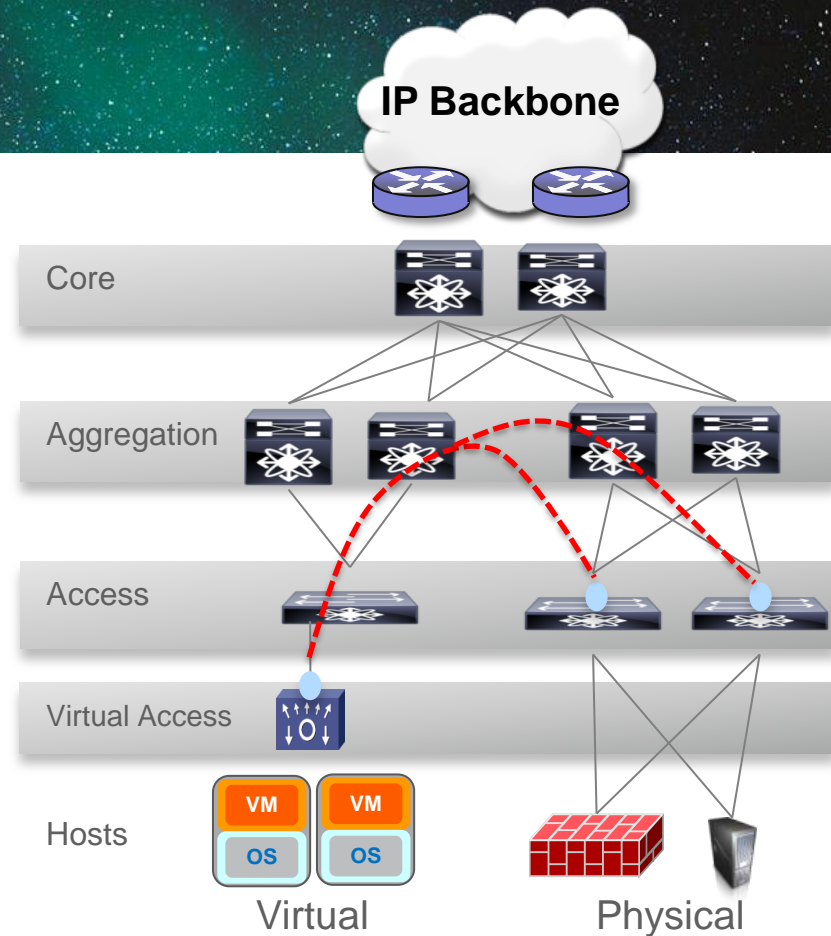
Tiered Applications and Virtualisation



- Many segments to interconnect app tiers
 - Small segments, isolated/concealed between tiers
 - One app will have multiple tier-segments
- Front-end segments provide connectivity to the broader “physical” network
- Tenant Segmentation != app-tier segmentation
 - Tenant Segmentation (Front-end) in the Fabric
 - App-tier segments in an overlay or in the Fabric
- A mix of physical and virtual
 - Web and app layers are commonly virtualised
 - Web layer reachable via the front-end physical network
 - The DB layer and services often run on bare metal

Hybrid Overlays

- Hypervisors introduce an additional tier in the network: The virtual Access (virtual Switch)
- **VMs** connect to the virtual Access
 - **Host overlays** start at the virtual Access
 - Virtualisation based resiliency: **Single attached sites**
- **Physical hosts** connect to the physical Access
 - **Network overlays** start at the physical Access
 - Network resiliency: **Site multi-homing**
- A hybrid overlay allows the combination of physical and virtual resources



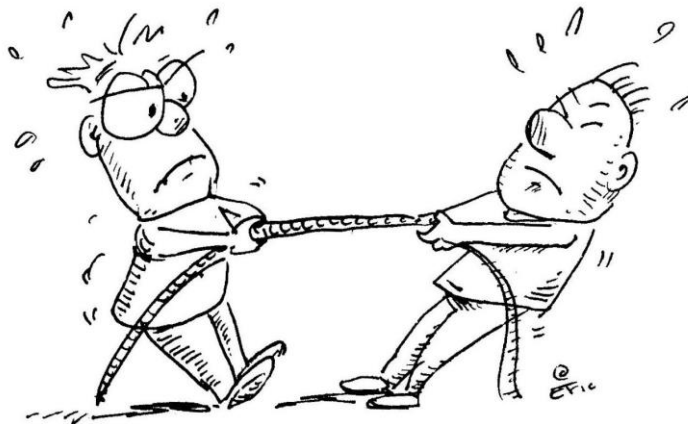
Which Encapsulation?

VXLAN

NVGRE

LISP

MPLS



STT

FabricPath

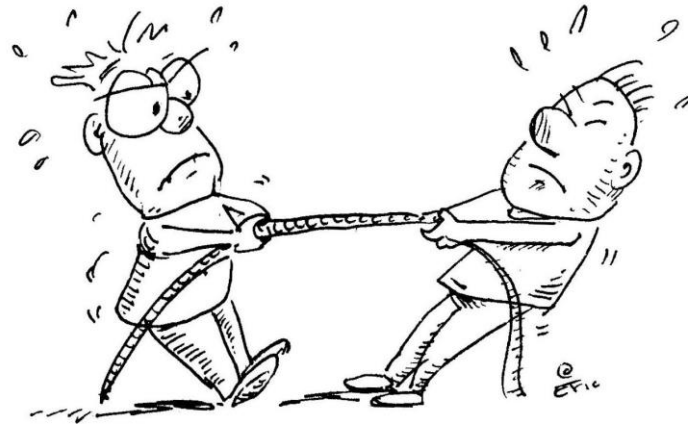
The Multi-protocol Router

TCP/IP

SPX/IPX

AppleTalk

Token Ring

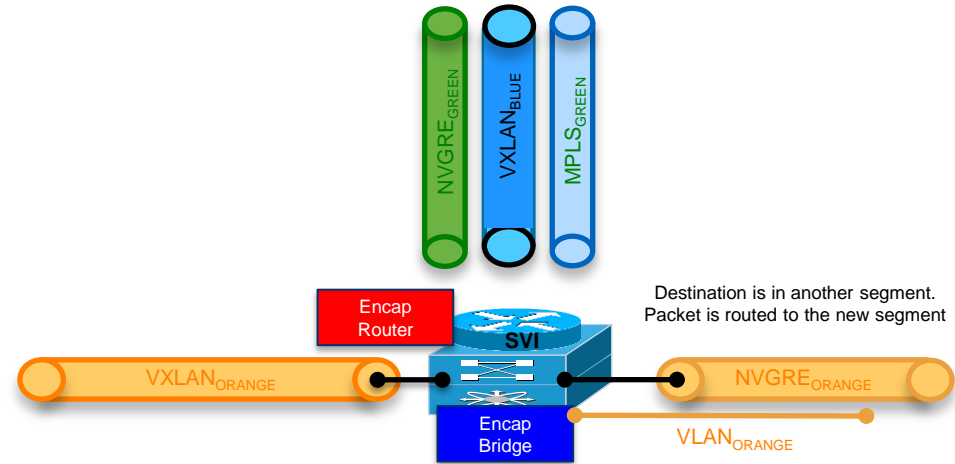


ATM

DECNet

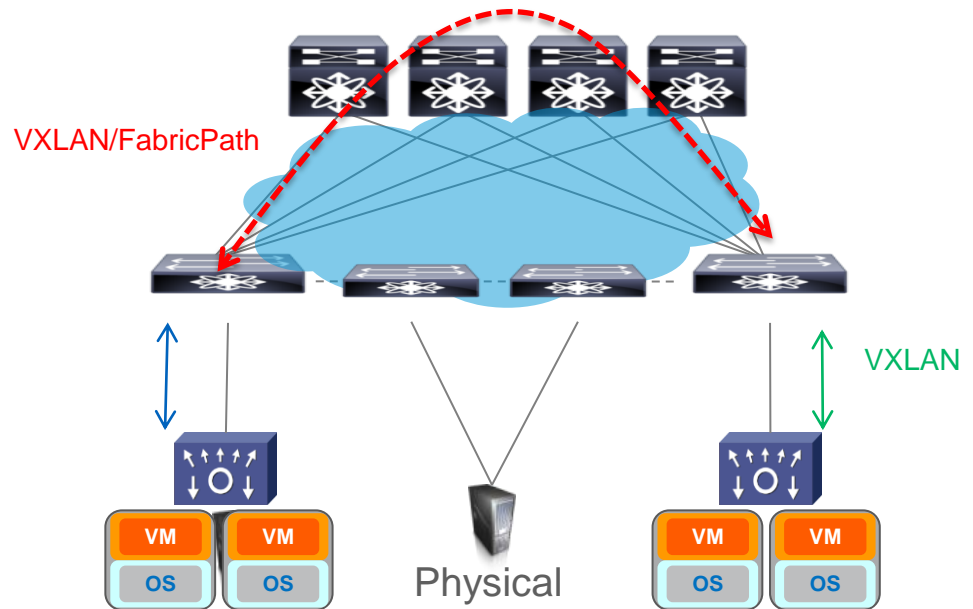
The Multi-encapsulation Gateway

- Multi-encapsulation Gateway:
 - VXLAN, NVGRE, MPLS, LISP, VLAN, OTV
- Bridging (L2 Gateway)
- Routing (L3 Gateway)
- Multiple TEPs in independent VRFs
- Nesting of IP overlays into MPLS VPNs
- Available across the product line

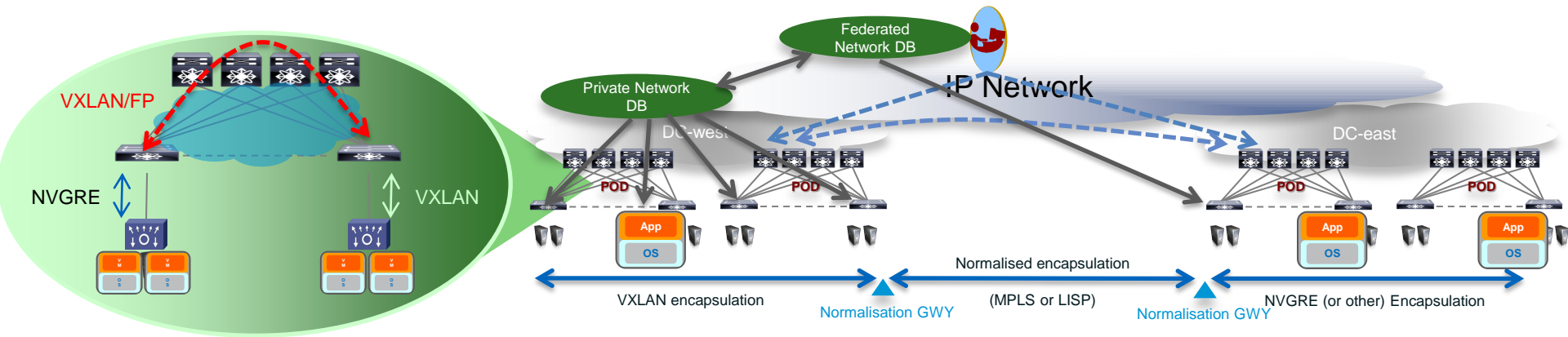


Normalisation: The Encapsulation Doesn't Matter

- Intelligence in the Control Plane
- Capabilities Exchange in Control Plane (negotiate encapsulation)
- Normalise to common encapsulation
- Pervasive Multi-encap Gateways for optimal traffic patterns



Data Plane and Control Plane Normalisation



- Multi-encapsulation Hardware Gateways
- Normalise to a common encapsulation in the Fabric and/or between Data Centres
- Terminate and map multiple types of encapsulation
 - VXLAN, NVGRE, MPLS, OTV, LISP
- Terminate and re-distribute information between overlay control protocols
 - Controllers, BGP, LISP

Encapsulation HW Offload

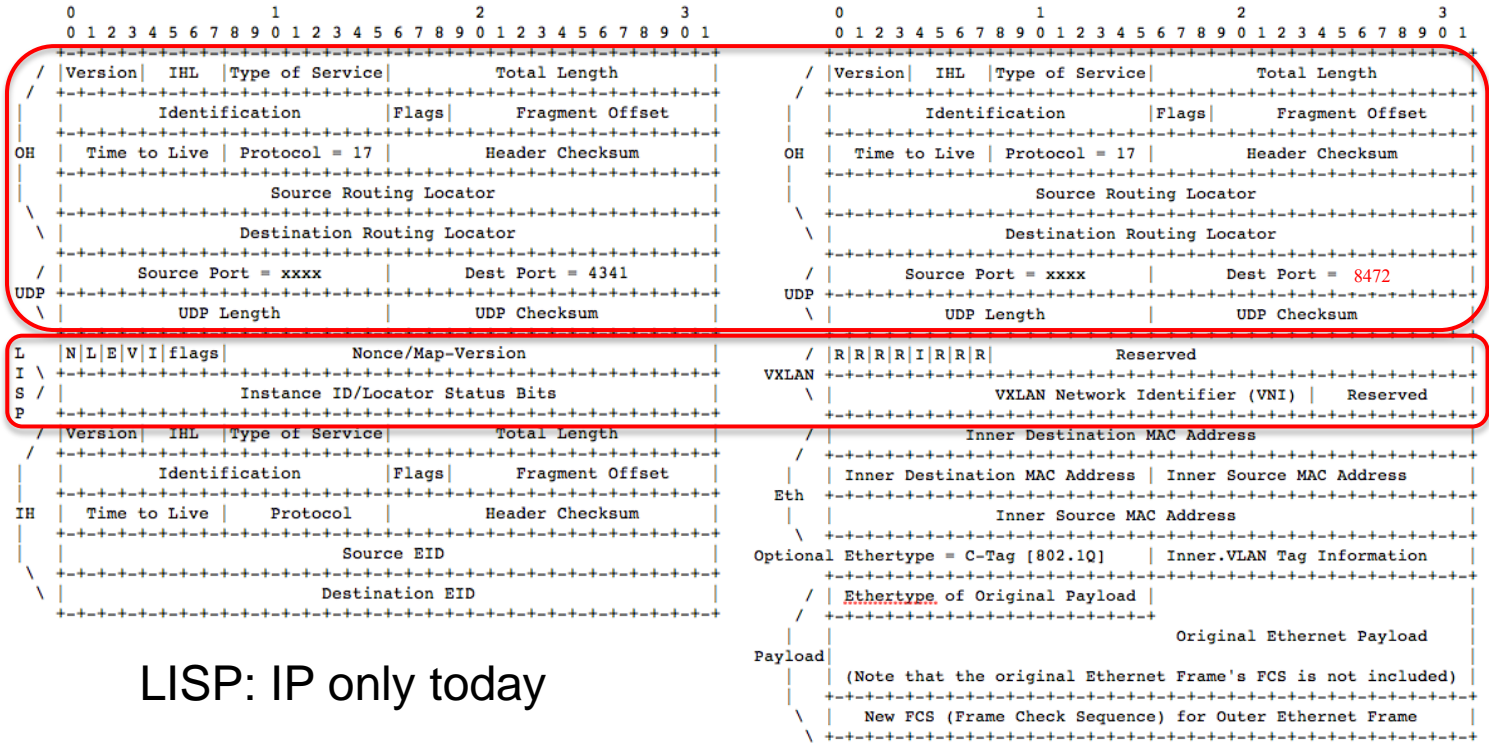
Host Overlays

- Current forwarding penalty for SW encap is about 50% throughput
- STT trick leverages TCP offload engine in existing NICs
 - TCP violation, short lived workaround
 - P2P only, no routing of flows
- VXLAN/NVGRE offload on NICs
 - The way forward for host overlays
 - Disruptive, many touch points
 - Static as ASICs: headers still in flux
 - Cisco 3rd Gen VIC 2HCY14 (stateless offload)

Network Overlays

- ASIC acceleration of overlay encapsulations
 - Cisco F-series ASICs with parser programmability
 - Fast enablement of incremental functions in header reserved fields without replacing HW
- Minimal disruption at the network access
 - Manageable number of touch points
- Encapsulation Normalisation
- Maximise throughput

LISP and VXLAN Headers Today



LISP: IP only today



LISP, OTV and VXLAN Normalisation with Generic Protocol Extension (gpe)

draft-lewis-lisp-gpe-00.txt

| | | | | |
|------|---------------------------------|---|---------------------------------|---|
| | 0 | 1 | 2 | 3 |
| | 0 | 1 | 2 | 3 |
| | 0 | 1 | 2 | 3 |
| / | Version | | IHL | |
| / | Type of Service | | Total Length | |
| / | Identification | | Flags | |
| / | Fragment Offset | | Fragment Offset | |
| OH | Time to Live | | Protocol = 17 | |
| | Header Checksum | | Header Checksum | |
| | Source Routing Locator | | | |
| \ | Destination Routing Locator | | | |
| / | Source Port = xxxx | | Dest Port = 4341 | |
| UDP | UDP Length | | UDP Checksum | |
| \ | UDP Length | | UDP Checksum | |
| / | N L E V I P R R | | Reserved | |
| LISP | Nonce/Map-Version/Protocol-Type | | Nonce/Map-Version/Protocol-Type | |
| \ | Instance ID/Locator-Status-Bits | | | |

draft-quinn-vxlan-gpe-00.txt

| | | | | |
|-------|--------------------------------|---|------------------|---|
| | 0 | 1 | 2 | 3 |
| | 0 | 1 | 2 | 3 |
| | 0 | 1 | 2 | 3 |
| / | Version | | IHL | |
| / | Type of Service | | Total Length | |
| / | Identification | | Flags | |
| / | Fragment Offset | | Fragment Offset | |
| OH | Time to Live | | Protocol = 17 | |
| | Header Checksum | | Header Checksum | |
| | Source Routing Locator | | | |
| \ | Destination Routing Locator | | | |
| / | Source Port = xxxx | | Dest Port = 4789 | |
| UDP | UDP Length | | UDP Checksum | |
| \ | UDP Length | | UDP Checksum | |
| / | R R R R I P R R | | Reserved | |
| VXLAN | Protocol Type | | Protocol Type | |
| \ | VXLAN Network Identifier (VNI) | | Reserved | |

Ethernet or IP Payload: Defined in the Protocol Type
 Common encapsulation for LISP and VXLAN
 L2 and L3 Payloads in both LISP and VXLAN

Header Evolution: Metadata in Overlay Headers

- Segmentation (VRFs, VPNs, Segments)
- L2 and L3 Payloads
- Policy
- Service Chaining
- Underlay integration (load balancing, traffic engineering)

LISP, OTV and VXLAN GPE Plus Network Service Header

draft-quinn-nsh-01
standardises the definition for the vPath 3.0 header

Base Service Header:

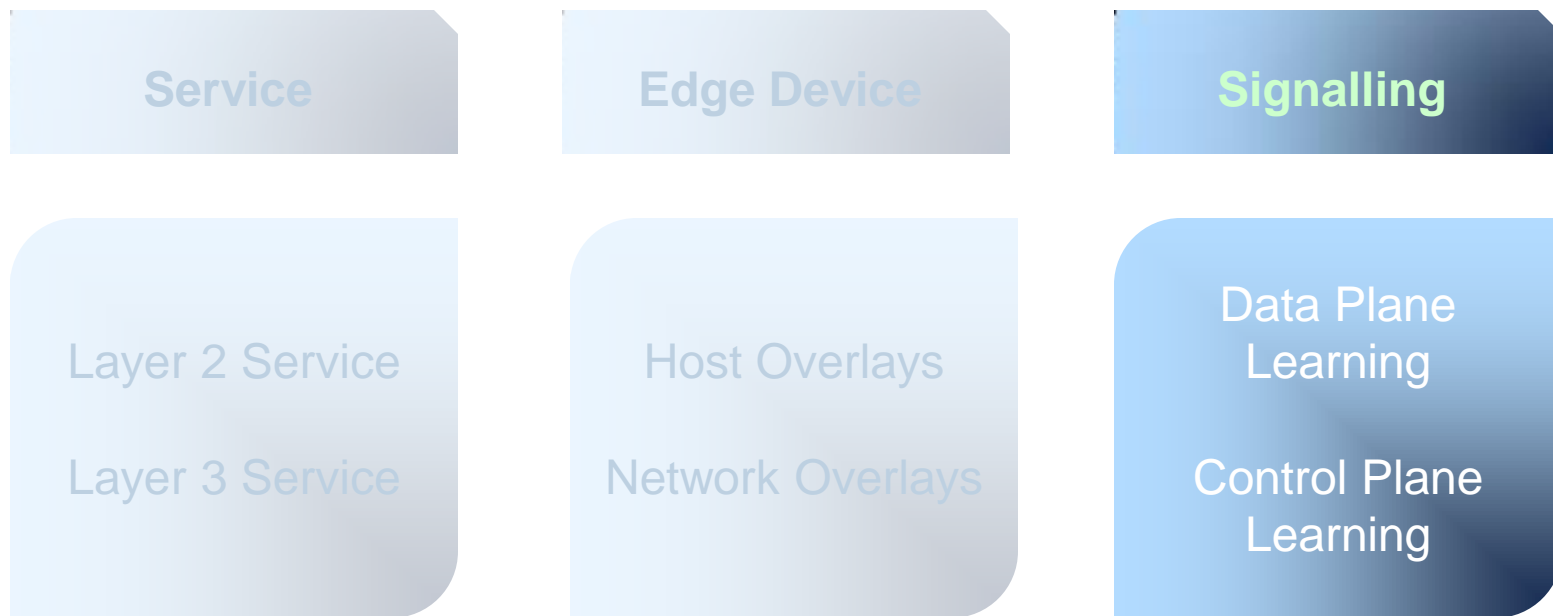
| | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | |
|-------------------|---|---|---|---|---|---|---|---|---|---------------|---|---|---|---|---|---|---|---|---|---------------|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|
| 0 | | | | | | | | | | 1 | | | | | | | | | | 2 | | | | | | | | | | 3 | | | | | | | | | |
| 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9 | 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9 | 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9 | 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9 |
| O C R R R R R R R | | | | | | | | | | Protocol Type | | | | | | | | | | Service Index | | | | | | | | | | | | | | | | | | | |
| | | | | | | | | | | Service path | | | | | | | | | | Reserved | | | | | | | | | | | | | | | | | | | |

Protocol Type =
0xNSH

Protocol Type =
IP

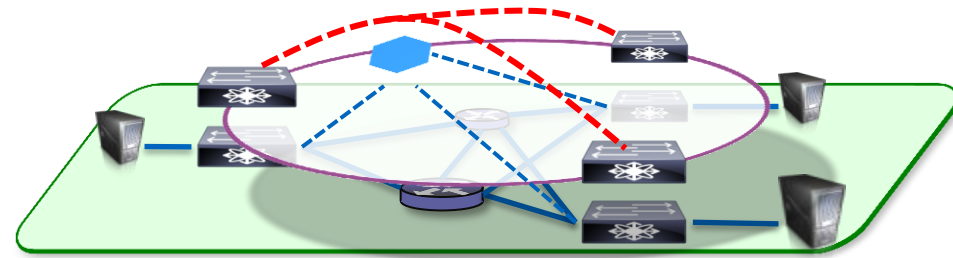
| | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | |
|------------|---|---|---|---|---|---|---|---|---|---------------------------------|---|---|---|---|---|---|---|---|---|----------------------|---|---|---|---|---|---|---|---|---|---------------------------------|---|---|---|---|---|---|---|---|---|
| 0 | | | | | | | | | | 1 | | | | | | | | | | 2 | | | | | | | | | | 3 | | | | | | | | | |
| 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9 | 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9 | 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9 | 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9 |
| / Version | | | | | | | | | | IHL Type of Service | | | | | | | | | | Total Length | | | | | | | | | | | | | | | | | | | |
| | | | | | | | | | | Identification | | | | | | | | | | Flags | | | | | | | | | | Fragment Offset | | | | | | | | | |
| OH | | | | | | | | | | Time to Live Protocol = 17 | | | | | | | | | | Header Checksum | | | | | | | | | | | | | | | | | | | |
| | | | | | | | | | | Source Routing Locator | | | | | | | | | | | | | | | | | | | | | | | | | | | | | |
| | | | | | | | | | | Destination Routing Locator | | | | | | | | | | | | | | | | | | | | | | | | | | | | | |
| | | | | | | | | | | Source Port = xxxx | | | | | | | | | | Dest Port = 4341 | | | | | | | | | | | | | | | | | | | |
| UDP | | | | | | | | | | UDP Length | | | | | | | | | | UDP Checksum | | | | | | | | | | | | | | | | | | | |
| | | | | | | | | | | N L E V I P R R | | | | | | | | | | Reserved | | | | | | | | | | Nonce/Map-Version/Protocol-Type | | | | | | | | | |
| LISP | | | | | | | | | | Instance ID/Locator-Status-Bits | | | | | | | | | | | | | | | | | | | | | | | | | | | | | |
| | | | | | | | | | | Base Header | | | | | | | | | | | | | | | | | | | | | | | | | | | | | |
| | | | | | | | | | | Context Header | | | | | | | | | | | | | | | | | | | | | | | | | | | | | |
| NSH | | | | | | | | | | Context Header | | | | | | | | | | | | | | | | | | | | | | | | | | | | | |
| | | | | | | | | | | Context Header | | | | | | | | | | | | | | | | | | | | | | | | | | | | | |
| | | | | | | | | | | Context Header | | | | | | | | | | | | | | | | | | | | | | | | | | | | | |
| | | | | | | | | | | Context Header | | | | | | | | | | | | | | | | | | | | | | | | | | | | | |
| | | | | | | | | | | Version | | | | | | | | | | IHL Type of Service | | | | | | | | | | Total Length | | | | | | | | | |
| | | | | | | | | | | Identification | | | | | | | | | | Flags | | | | | | | | | | Fragment Offset | | | | | | | | | |
| IH | | | | | | | | | | Time to Live Protocol | | | | | | | | | | Header Checksum | | | | | | | | | | | | | | | | | | | |
| | | | | | | | | | | Source EID | | | | | | | | | | | | | | | | | | | | | | | | | | | | | |
| | | | | | | | | | | Destination EID | | | | | | | | | | | | | | | | | | | | | | | | | | | | | |

Overlay Signalling Evolution



Overlay Signalling

- Service Discovery
 - Edge devices in an overlay need to discover each other
- Address Advertising and Tunnel Mapping
 - Edge devices must exchange host reachability information
 - Map end-point to location
- Tunnel Management
 - Maintain and manage connections between edge devices



Overlay Signalling

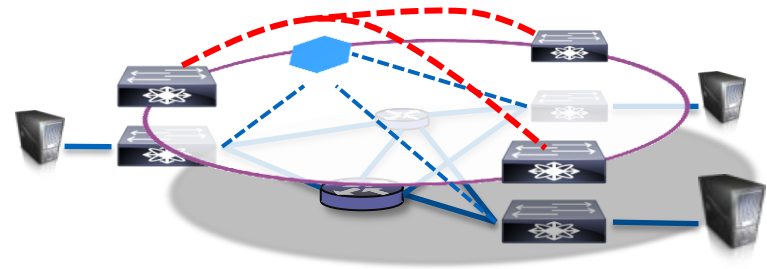
Data Plane Learning

- Based on gleaning information from data plane events
 - Example: Source Learning on bridges
- Provides the following:
 - Address advertisement/mapping (very effectively)
 - Some tunnel management is possible
 - Does not provide Service Auto-discovery
- Requires a flood facility for data plane events to propagate:
 - Multicast tree
 - Unicast replication group at the head-end
- Flood facility can be manually configured on every device (e.g. join a mcast group or configure a list of unicast destinations)
- Usually is supplemented with a control protocol for Service Discovery (specially if using unicast replication)

Overlay Signalling

Control Plane

- Provides:
 - Service Discovery
 - Address Advertising/Mapping
 - Tunnel Management
 - Extensions for multi-homing and advanced services can be provided



Protocol or Controller:

- **Routing Protocol** amongst Edge Devices
 - BGP, IS-IS, LISP
- Central database on a **Controller**
 - Distributed Virtual Switches (OVS, N1Kv/VSM)

Push or Pull:

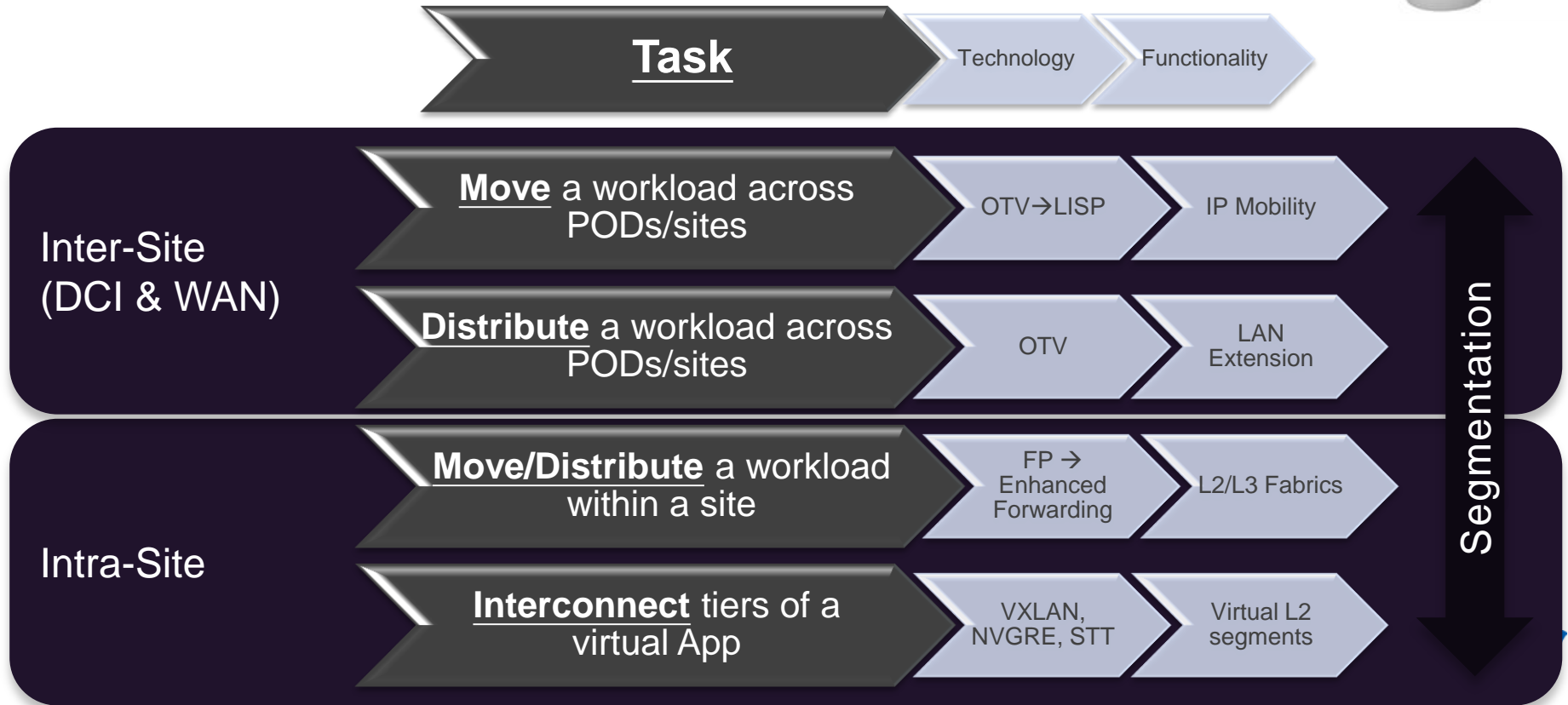
- **Push** all information to all Edge Devices
 - BGP, IS-IS, Controllers
- **Pull** and cache on demand @ ED
 - LISP, DNS, Controllers



Overlays Evolve to Meet Network Challenges

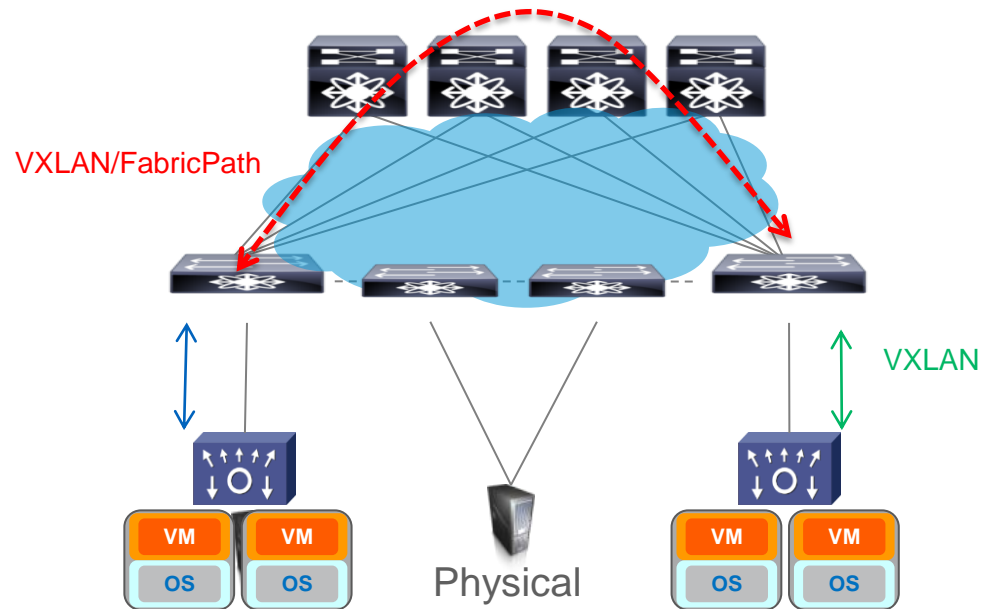
Handling Workloads in the Data Centre

Choosing the right tools ...



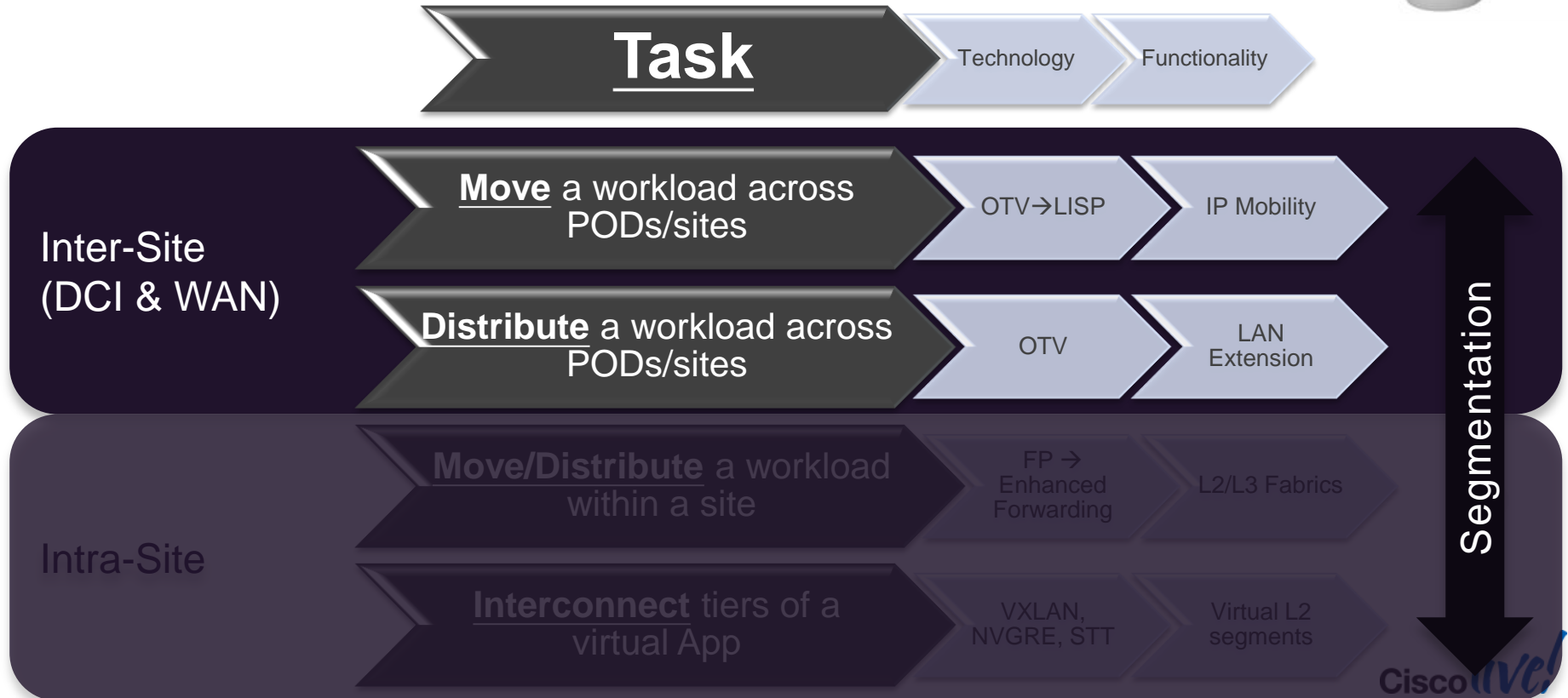
DC-Fabric: Normalised L2/L3 Network Overlays

- Terminate the encapsulation from the host overlay
- Translate to a normalised encapsulation in the fabric
- Seamlessly allow physical and virtual to connect to the fabric
- Fabric overlay provides L2 and L3 services with mobility and segmentation

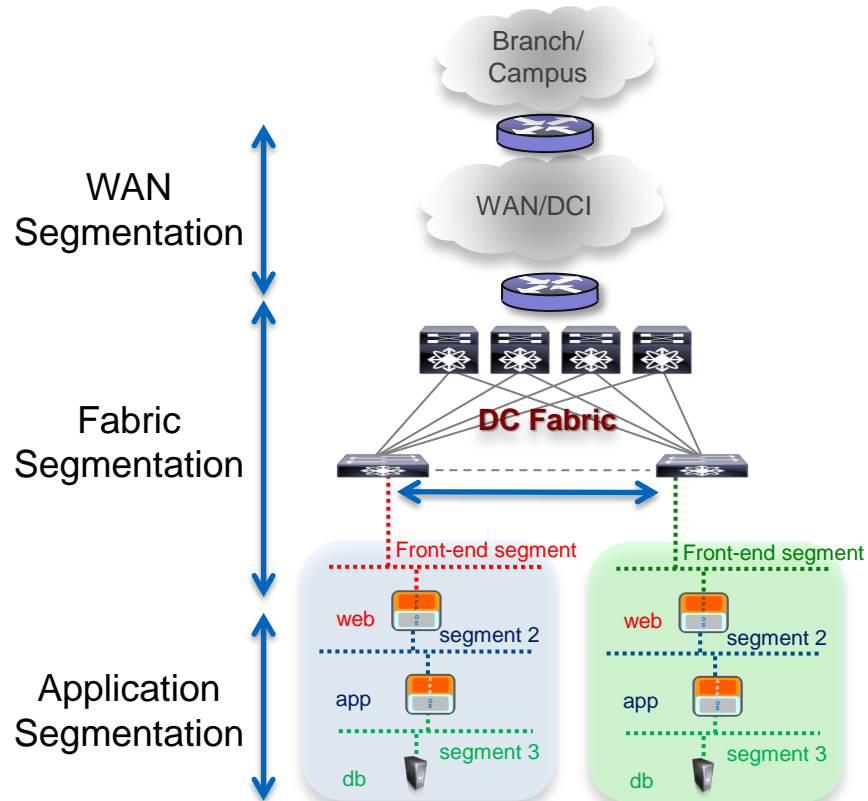


Handling Workloads in the Data Centre

Choosing the right tools ...

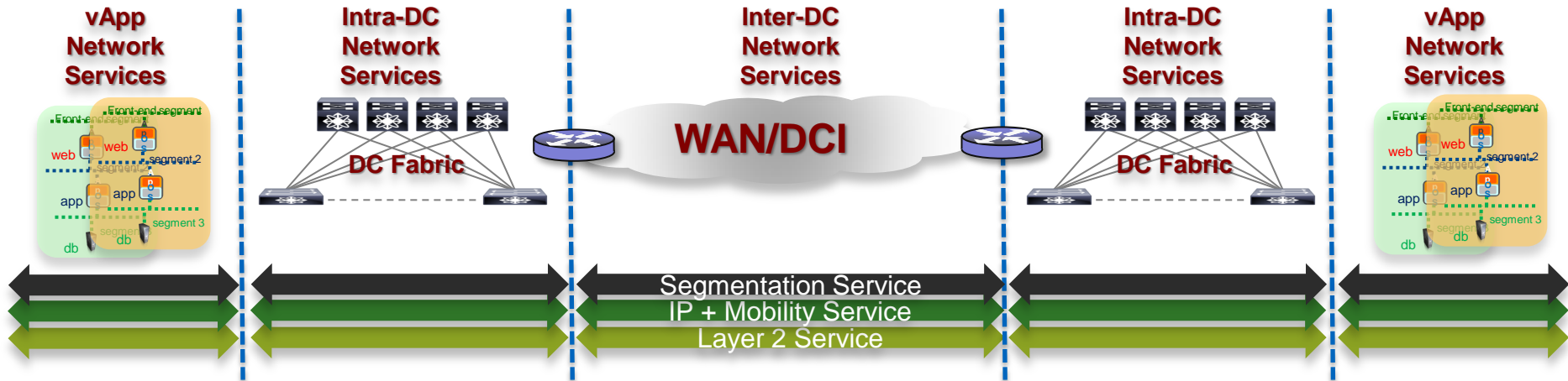


Segmentation End-to-end



- Segmentation at many levels
- Must be given continuity
 - Across the different network places
 - Across organisations and administrative boundaries
- All relevant technologies include the required segmentation semantics
- The network maps the segments together to provide a scalable and interoperable e2e segmentation solution

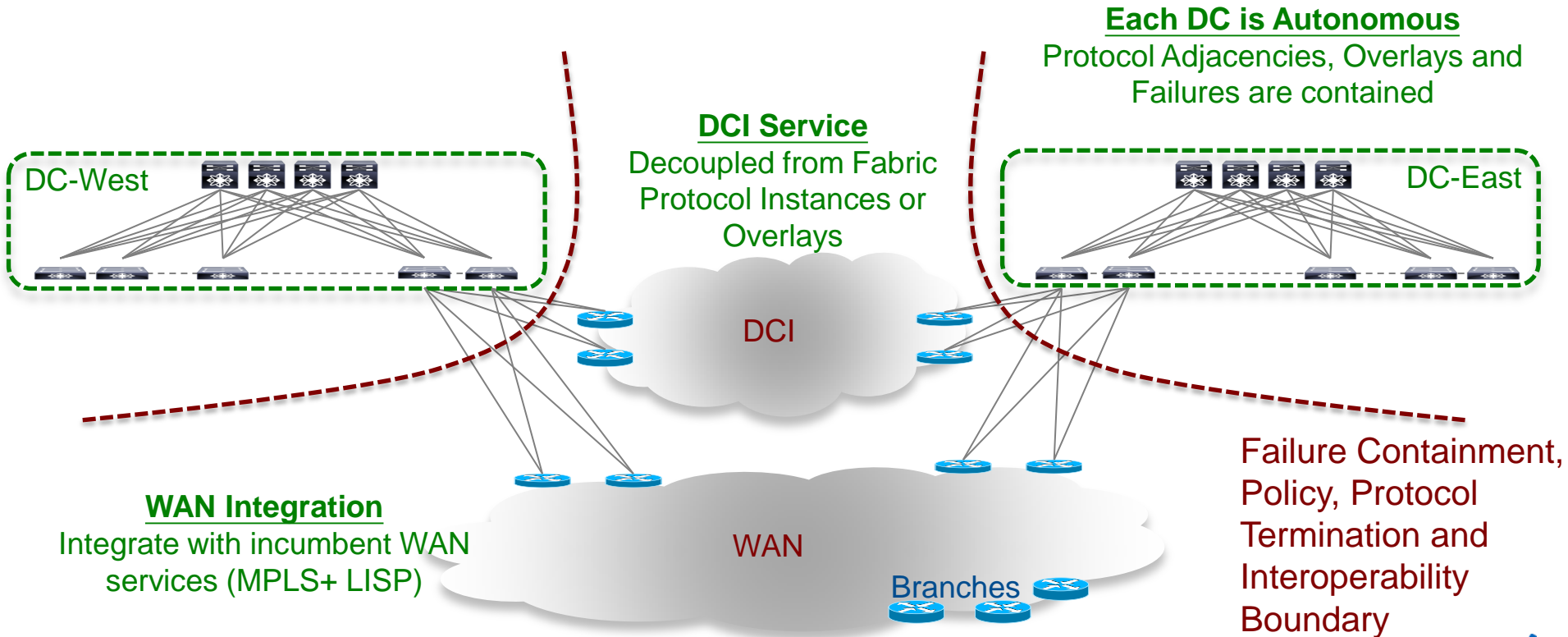
Failure Domain Scope



Core Principles of Network Resiliency/Scale applied to Overlay Services

- Clearly delineated Fault Boundaries and service domains
- Control Plane Hierarchy and Federation within and across domains
- Data Plane Boundaries
- Administrative Domain Delineation and Federation

DCI and WAN Integration

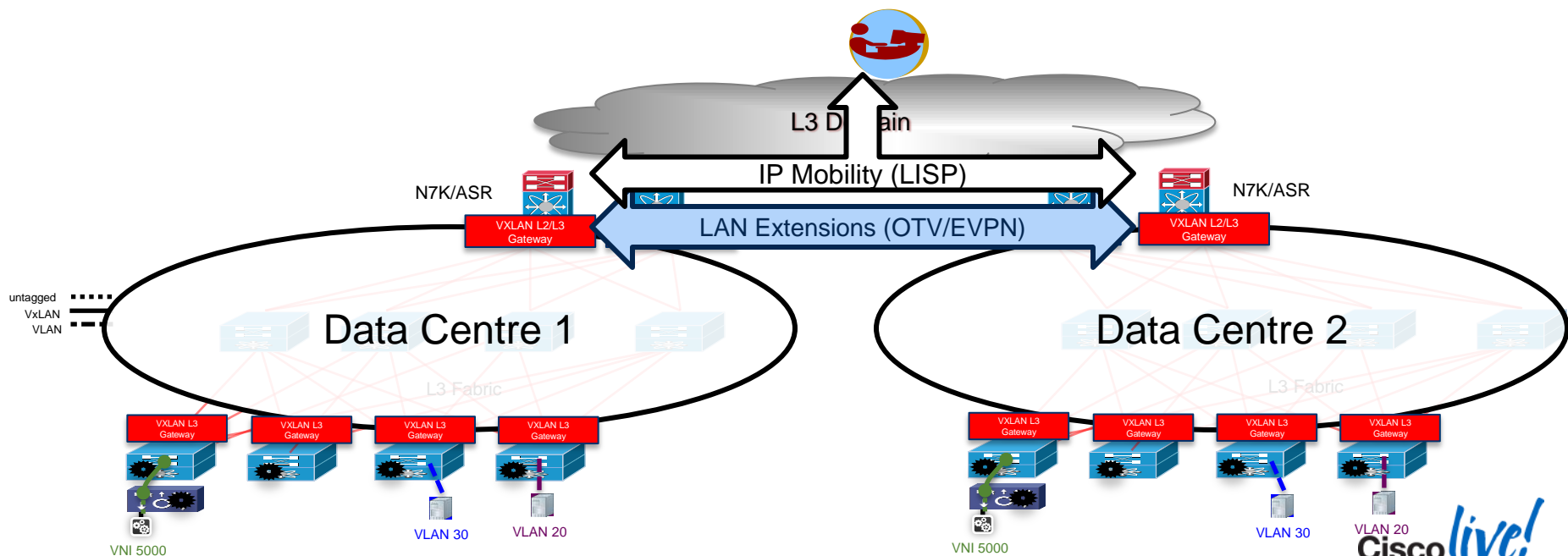


End-to-end Segmentation and Mobility

Interconnecting Multiple Data Centres

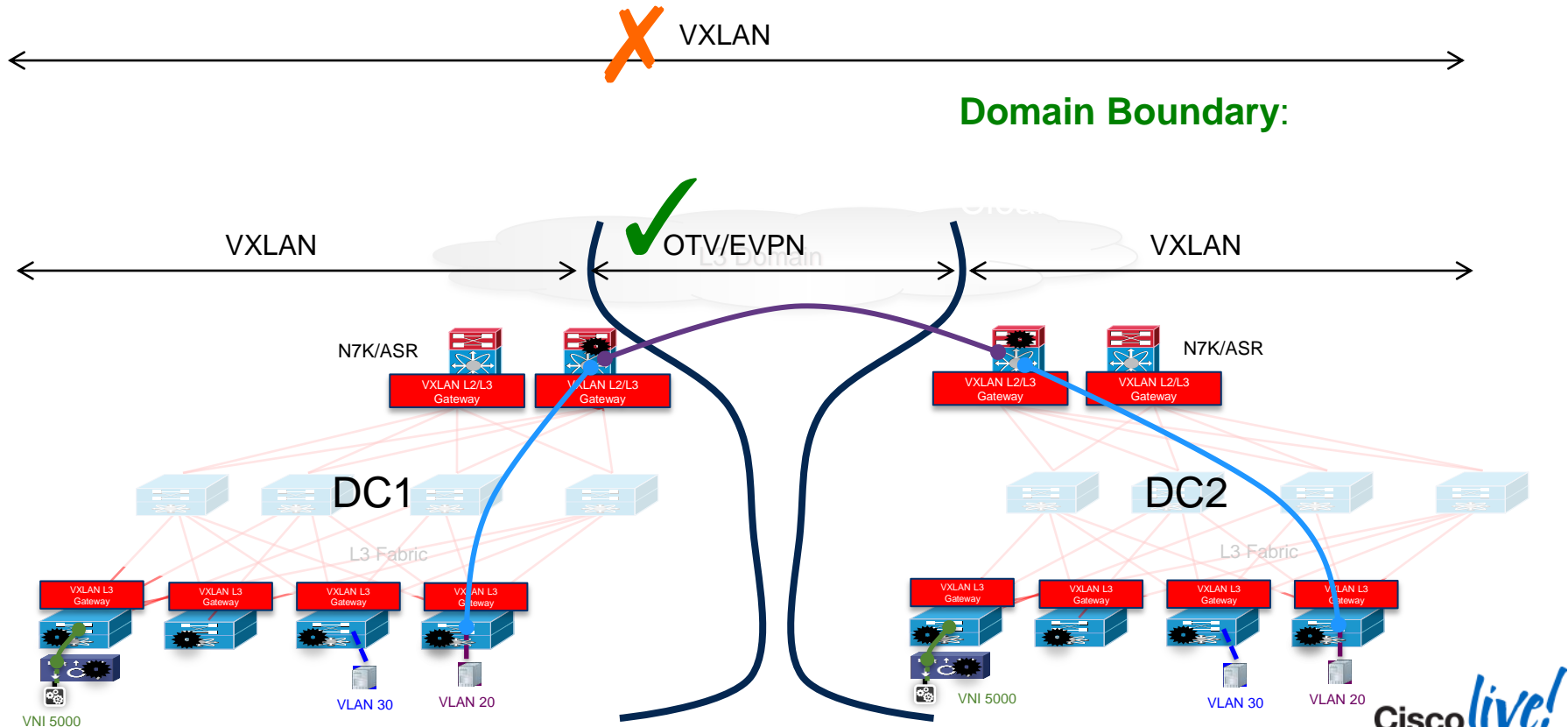
LAN Extensions and IP mobility

Ethernet extensions between independent fabrics
IP traffic is forwarded via the optimal path (no hair-pinning)



Interconnecting Multiple Data Centres

LAN Extensions



Interconnecting Multiple Data Centres

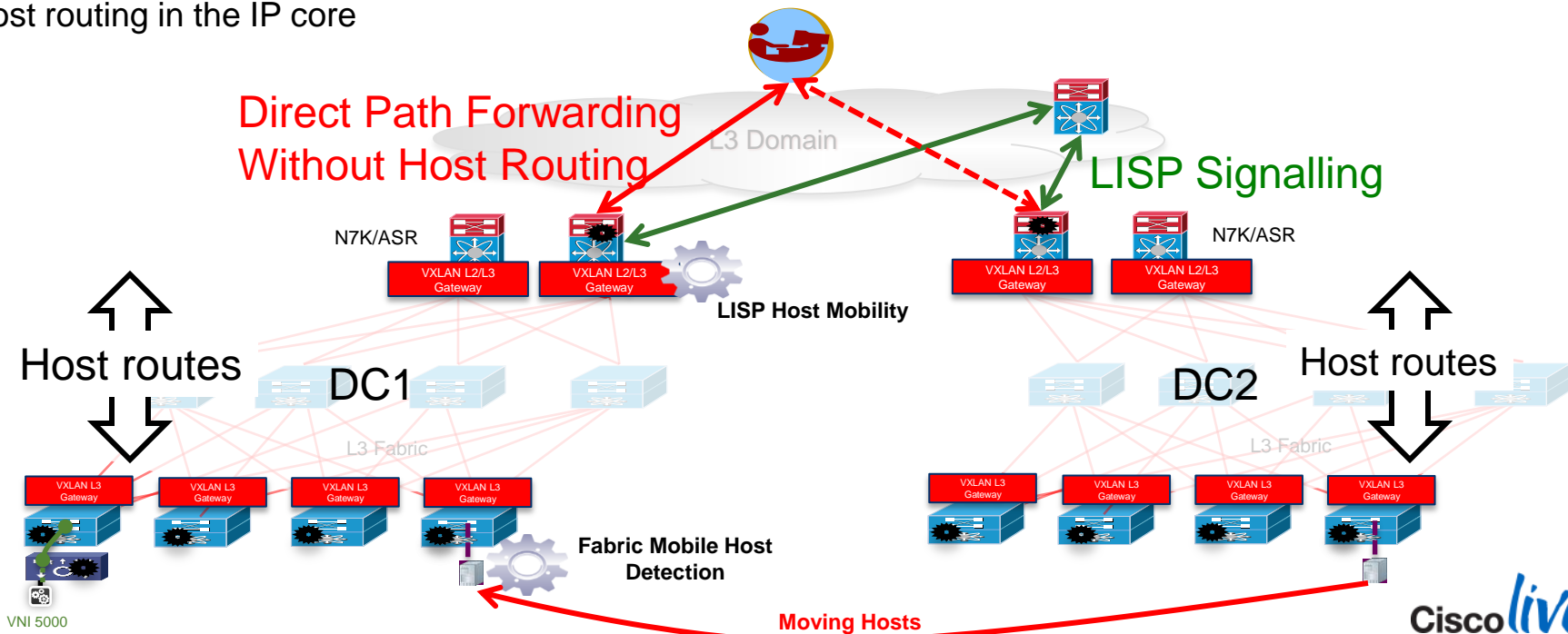
LISP IP Mobility for optimised routing

LISP Mobility:

- LISP registrations and notifications
- LISP encapsulation from client sites
- No host routing in the IP core

LISP Signalling:

Relay mobility state between sites



VNI 5000

BRKDCT-2328

© 2014 Cisco and/or its affiliates. All rights reserved.

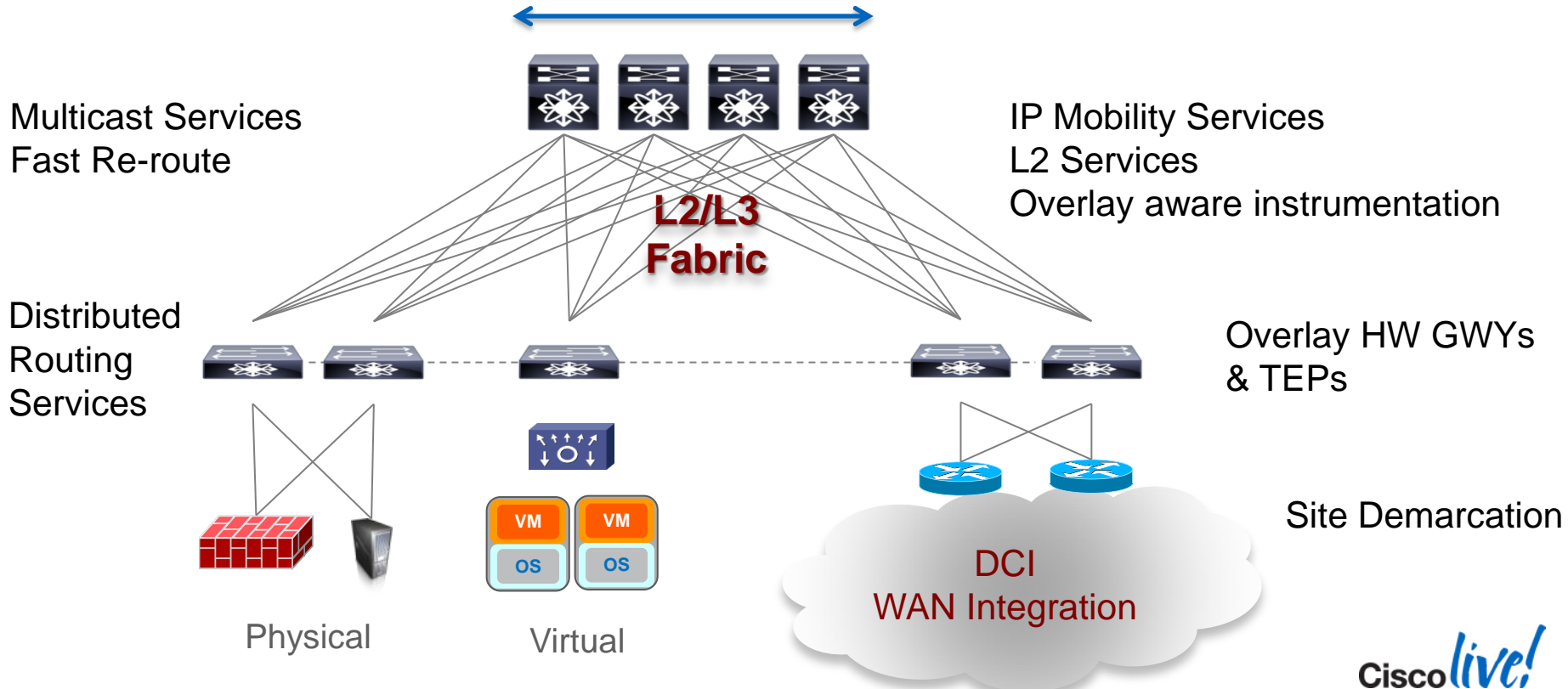
Cisco Public

Cisco *live!*



Role of the Underlay

Fabric Relevance to a Hybrid Overlay



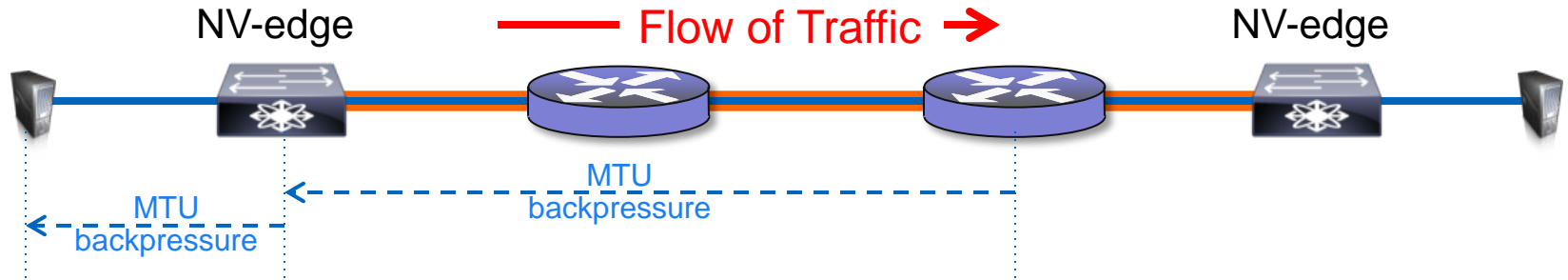
Encapsulation and Effective Throughput



1500bytes/packet (10Gbps) → 1542 bytes/packet (10.1 Gbps)
64bytes/packet (10Gbps) → 106 bytes/packet (10.3 Gbps)

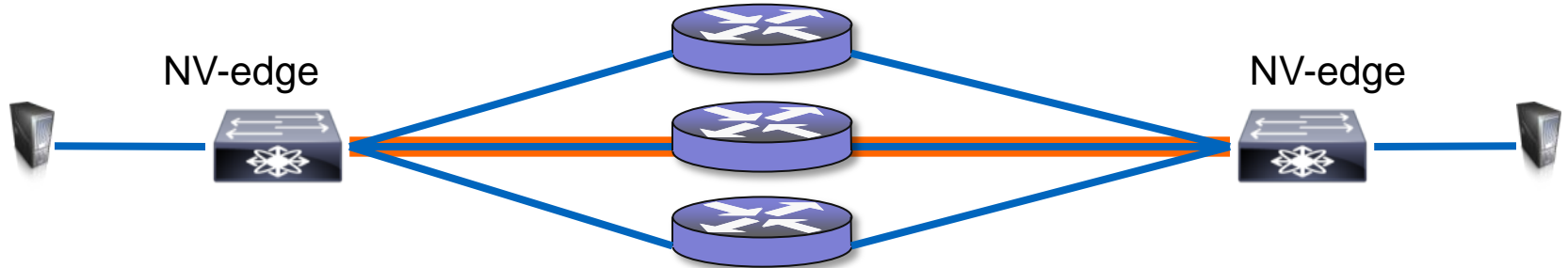
- Encapsulation adds bits to the traffic being sent
- When receiving traffic at full line rate, the encapsulated traffic will exceed the line-rate BW of the egress interface
 - Packet drops
 - Diminished effective throughput
- The uplink BW should be greater than the downlink BW to avoid congestion by encapsulation
 - This is naturally done in the network

MTU Issues



- Encapsulated traffic may exceed max MTU of the path
- When traffic is encapsulated with the Don't Fragment (DF) bit set:
 - If MTU is exceeded: IGMP unreachable message (datagram-too-big) is sent back to the encapsulating NV-edge
 - Encapsulating NV-edge will lower the tunnel MTU accordingly
 - Subsequent packets from the source will trigger an ICMP unreachable message from the NV-edge back to the server (if the traffic from the source has the DF bit set)
- If the DF bit is not set, the device sensing the MTU is exceeded should attempt to fragment the traffic

Multi-pathing and Entropy



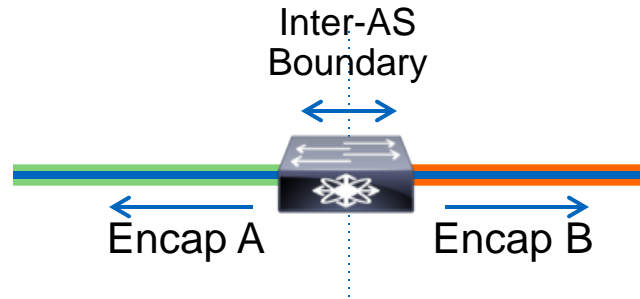
- Tunnel Polarisation: All encapsulated flows tend to look like a single flow between a pair of edge devices
 - Encapsulated traffic always hashes to a single path
- Adding entropy to the encapsulation header can depolarise the tunnels
 - Use all available paths
- UDP headers: Variable UDP source port
- GRE headers: Variable key field
- MPLS headers: Variable LSP label

Instrumentation and Overlay Awareness



- Infrastructure awareness of encapsulated traffic:
 - Outer/Encapsulation header
 - Overlay shim header
 - Internal/Payload header
 - Payload
- Overlay aware Switching & Routing infrastructure:
 - ACLs, QoS, Netflow
- Network Analysis Module (NAM) inspects encapsulated traffic

Data Plane and Control Plane Normalisation

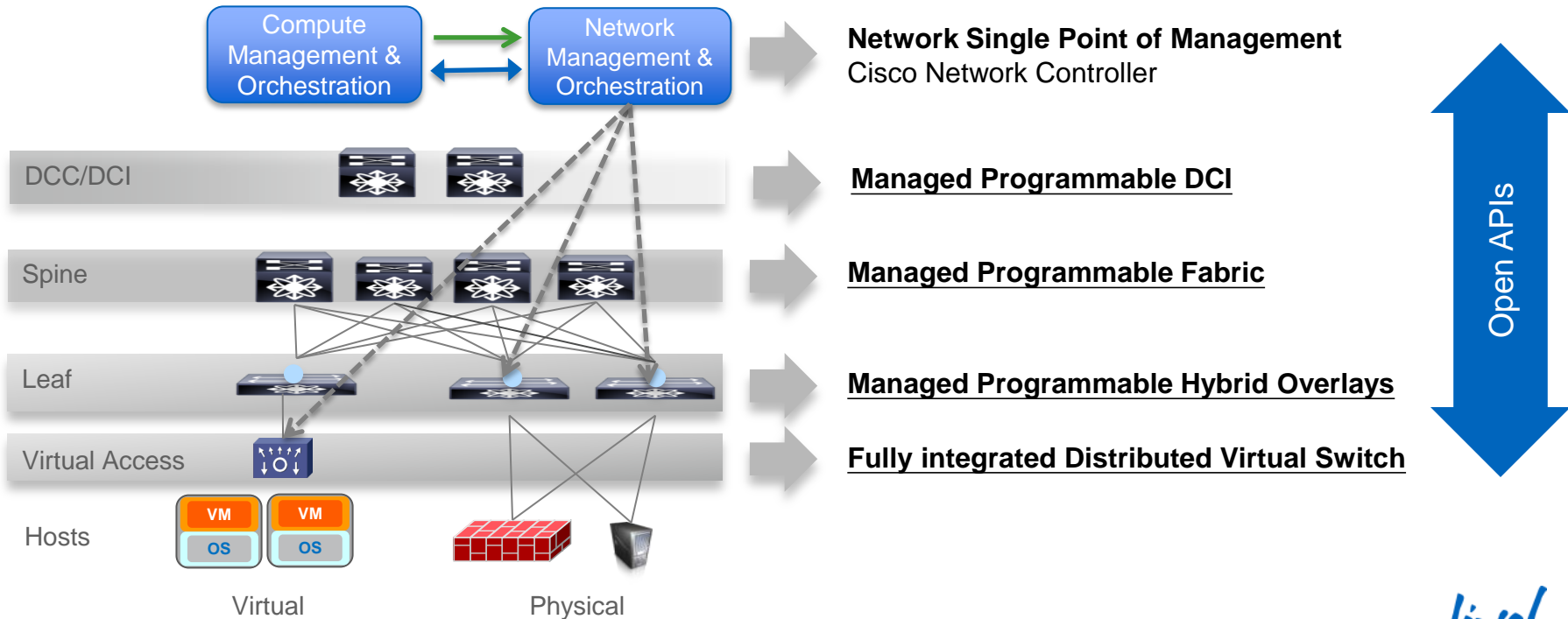


- Multi-protocol overlay gateway
- Terminate and map multiple types of encapsulation
 - VXLAN, NVGRE, MPLS, OTV, LISP
- Terminate and re-distribute information between overlay control protocols
 - Controllers, BGP, LISP



Management and Orchestration

Data Centre Fabric Management



Operational Delineation

Network
Management

↔
Programmatic APIs

Compute
Management



Network team



Nexus OS CLI
DCNM GUI
APIs: One-PK/Openflow

- Network Policies aligned with Compute
- Integrated network management:
 - Fabric + Overlays
- Programmatic APIs



Infrastructure
team



vCenter / UCS Manager
Interface

- Build Servers and Create Port profiles
- Create VMs and Assign Port profiles to VM

No hand-off required between Infrastructure and Network Admins for adding new servers

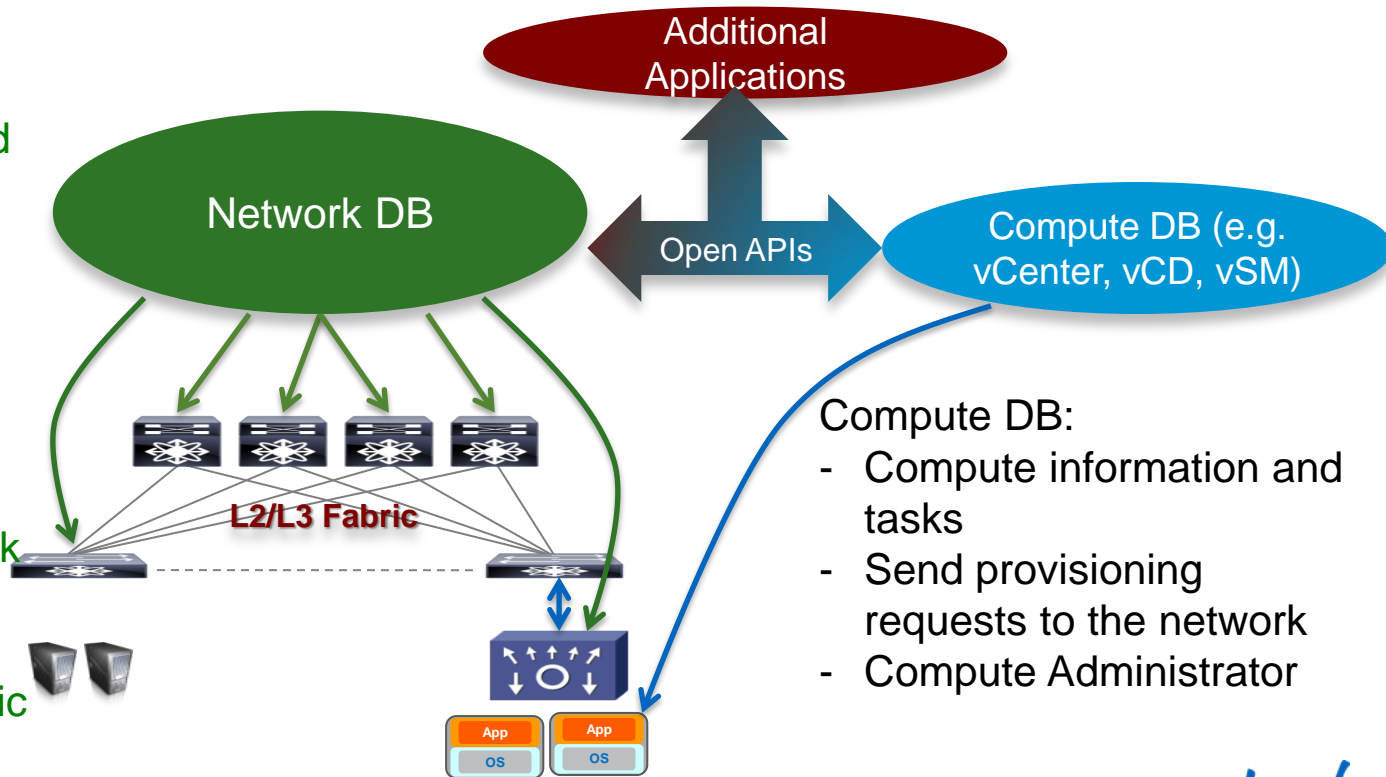
Orchestration, Management and Programmability

Network DB:

- Network information and tasks
- Network Administrator

Single point of Management:

- Broker communications between compute orchestrator and network
- Act on provisioning requests from compute
- Manage underlying fabric and overlay

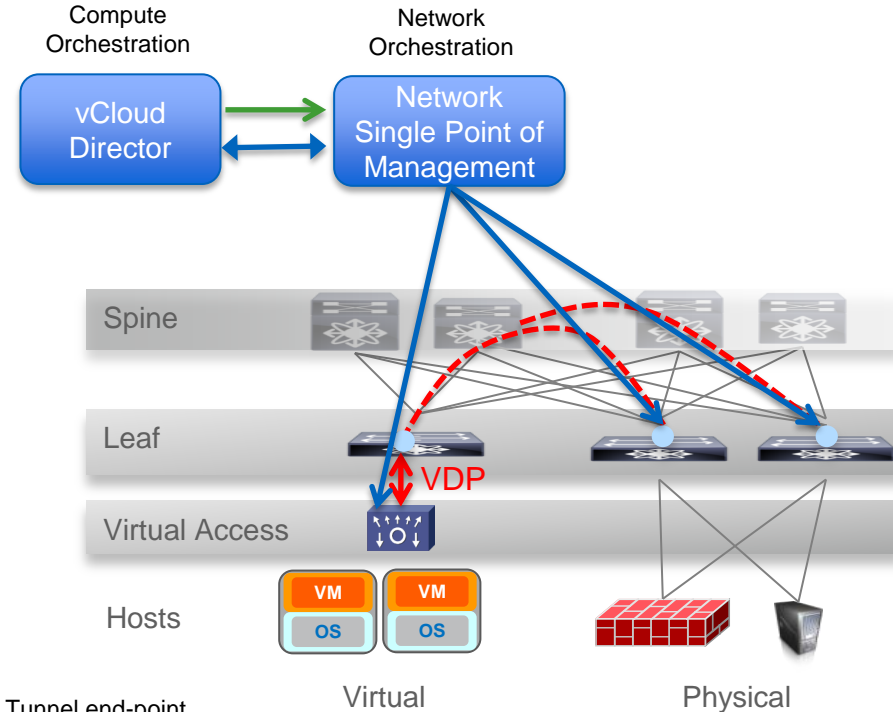


Compute DB:

- Compute information and tasks
- Send provisioning requests to the network
- Compute Administrator

Distributed Virtual Switching

- Compute Controller (e.g. vCloud Director) integrated overlay provisioning
 - Integrates physical and virtual endpoints
- Overlay encap/decap can be offloaded to network hardware
 - VDP = VSI Discovery Protocol (IEEE 802.1Qbg)
 - Cisco, HP, IBM, Brocade, Qlogic, Emulex, Broadcom, Mellanox, others ...



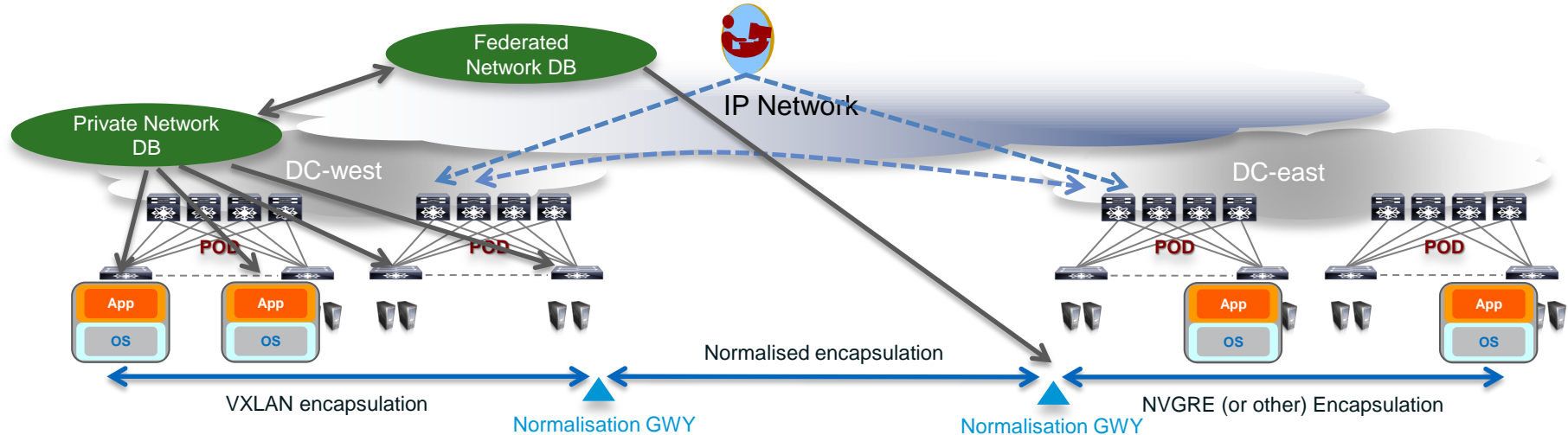
Federated/Normalised Overlays Vision

Inter-DC and Intra-DC – LISP/BGP Protocol + Any encapsulation

Virtual and Physical Hosts

Layer 2 and Layer 3

Internet Scale





Q & A

Complete Your Online Session Evaluation

Give us your feedback and receive a Cisco Live 2014 Polo Shirt!

Complete your Overall Event Survey and 5 Session Evaluations.

- Directly from your mobile device on the Cisco Live Mobile App
- By visiting the Cisco Live Mobile Site www.ciscoliveaustralia.com/mobile
- Visit any Cisco Live Internet Station located throughout the venue

Polo Shirts can be collected in the World of Solutions on Friday 21 March 12:00pm - 2:00pm



Learn online with Cisco Live!

Visit us online after the conference for full access to session videos and presentations.

www.CiscoLiveAPAC.com



CISCO TM