

TOMORROW starts here.



Cisco *live!*

LISP – A Next-Generation Networking Architecture

BRKRST-3045

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LISP - A Next Generation Routing Architecture

Agenda

- LISP Overview
- LISP Operations
- LISP Deployment Examples
- LISP Status
- LISP Summary
- LISP References



LISP Overview

Locator/ID Split and LISP

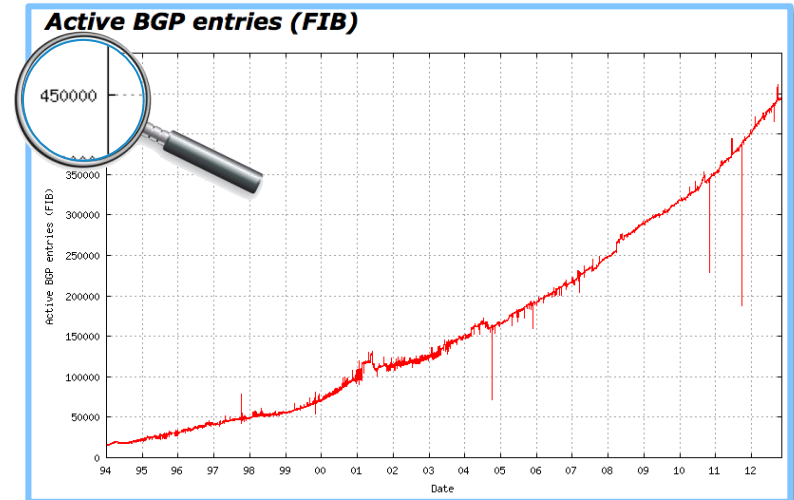
Routing and Addressing Architecture of the Internet Protocol

- Addresses today combine **location** and **identity** semantics in a single 32-bit or 128-bit number
- Separating Location and Identity changes this...
 - Provide a clear separation **at the Network Layer** between **what we are looking for** vs. **how best to get there**
 - ***Translation vs. Tunnelling is a key question***
- **Network Layer Identifier**: **WHO** you are in the network
 - long-term binding to the thing that they name, does not change often at all
- **Network Layer Locator**: **WHERE** you are in the network
 - Think of the source and destination “addresses” used in routing and forwarding
- **WHERE** you are can change! **WHO** you are should be the same!

LISP Overview

Original Motivation...

- An IP address “overloads” location and identity
 - Today... “addressing follows topology”
 - Efficient aggregation is only available for Provider Assigned (PA) addresses
 - Ingress Traffic Engineering usually requires Provider Independent (PI) addresses and the injection of “more specifics” :: this limits route aggregation compactness
 - IPv6 does not fix this
- Route scaling issues drive system costs higher
 - Forwarding plane (FIB) requires expensive memory
 - Route scaling “drivers” are also seen in Data Centres and for Mobility :: not just the Internet DFZ

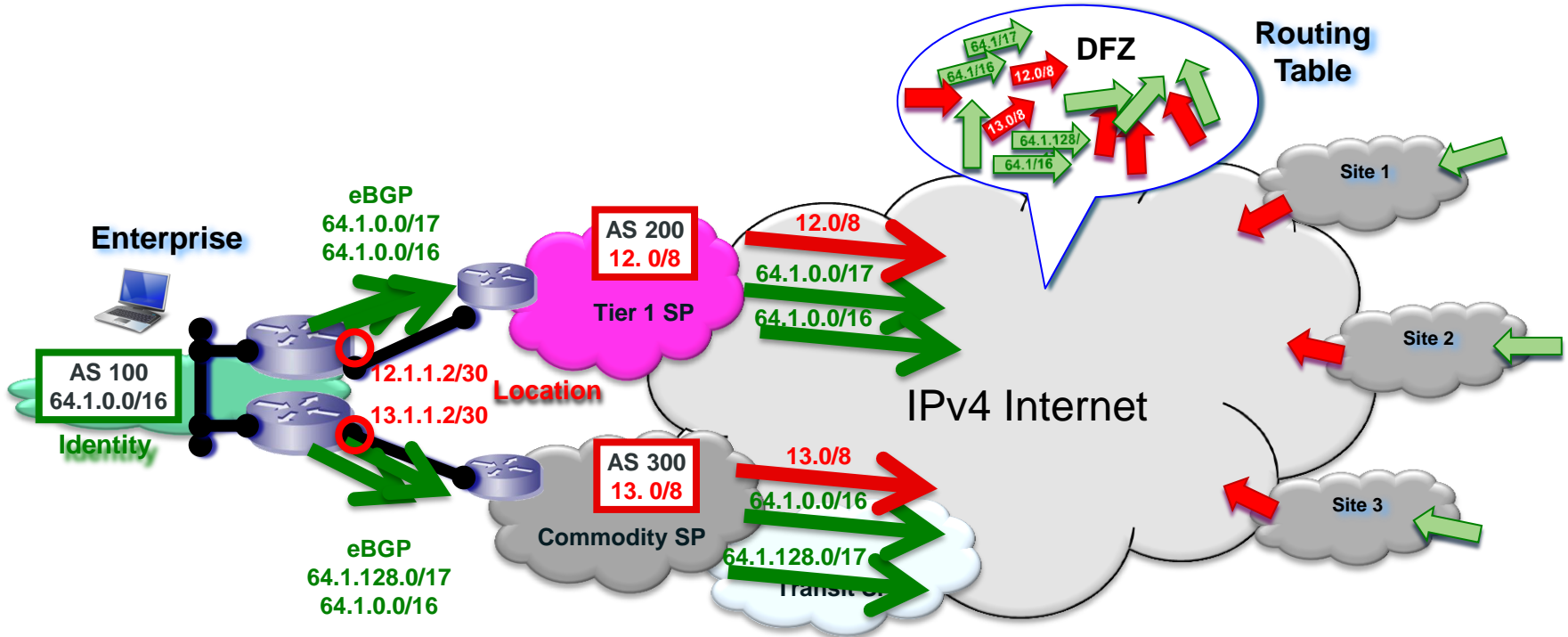


“... routing scalability is the most important problem facing the Internet today and must be solved ...”

Internet Architecture Board (IAB)
October 2006 Workshop (written as RFC 4984)

LISP Overview

Identity and Location :: an Overloaded Concept in Routing Today...

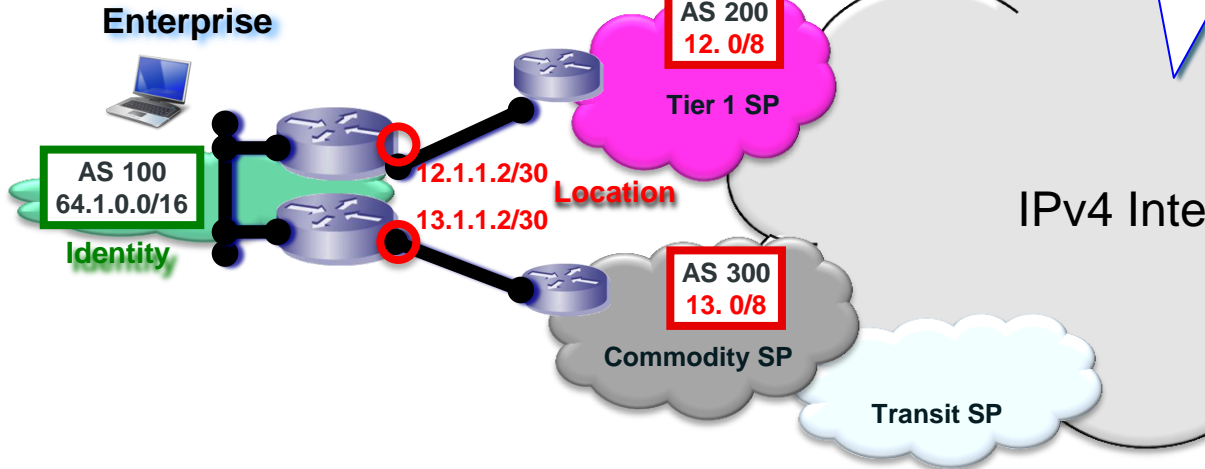


LISP Overview

Identity and Location :: an Overloaded Concept in Routing Today...

- Let's put **ID address** and **Locator address** in different databases
- Let's create a "level of **indirection**" between **ID** and **LOCATION** in the network!

LISP Mapping System



Clear Separation at the Network Layer::
•who/what you are looking for
VS. ...
•how to best get there

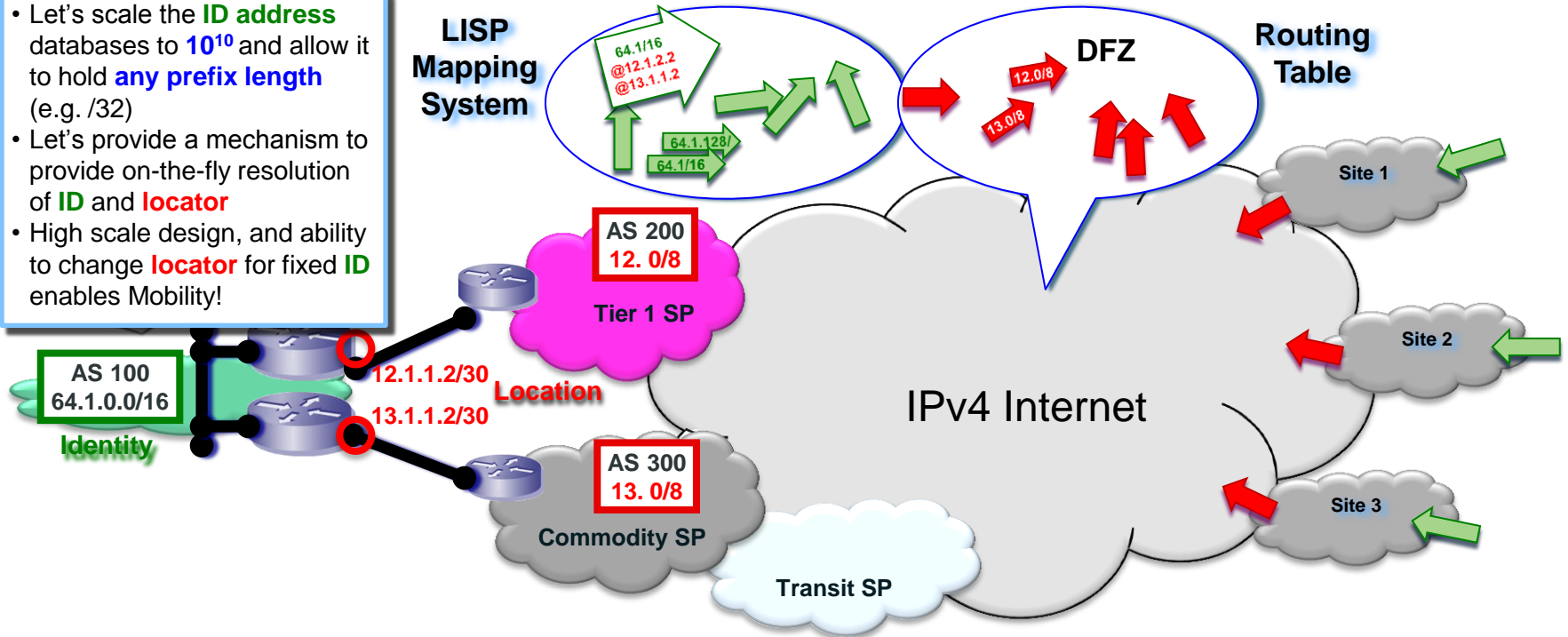
Two Approaches::
•Translations (e.g. NAT)
VS. ...
•Tunnels (e.g. GRE, IPsec, MPLS)

What is needed is Locator/ID Separation on a GLOBAL Scope, and that doesn't carry all routing in the Forwarding Plane!

LISP Overview

Identity and Location :: an Overloaded Concept in Routing Today...

- Let's scale the **ID address** databases to 10^{10} and allow it to hold **any prefix length** (e.g. /32)
- Let's provide a mechanism to provide on-the-fly resolution of **ID** and **locator**
- High scale design, and ability to change **locator** for fixed **ID** enables Mobility!



LISP Overview

LISP :: A Routing Architecture – Not a Feature

IDENTITY

LOCATION

level of indirection

LISP **changes** the current ROUTING Architecture

LISP Overview

Locator/ID Separation :: The Mapping System is the Key

- A Mapping Systems is “key” for a Locator/ID separation architecture
 - Mapping systems provide the control plane for the architecture
 - Mapping systems represent the great opportunity for these architecture to excel
 - Most of the time, users/operators think about the data plane
 - The control plane is where the magic happens!
- Some general components of a mapping system to be aware...
These affect how the system scales much differently than routing
 - state :: must scale to large numbers (such as 10^{10}) of hosts
 - rate :: must be small globally; damp reachability and mobility from globally impacting the system
 - latency :: must be low enough not to harm existing applications
 - scope :: must allow for both a global and a private scope for mapping

LISP Overview

Locator/ID Separation :: Changing the Routing Architecture

- Locator/ID Separation “architecture” helps solve other current network problems
- IPv4/IPv6 Co-existence at the “ID” and “Locator” spaces
 - IPv4 and IPv6 can be implemented at the “ID” and/or “locator” spaces for simple integration
 - In reality, **anything** can be an “ID” and carried over traditional cores (IPv4 and IPv6)
 - e.g. RFID, VIN#, Geo-Location, MAC-Addr, etc. etc. etc.
- Scaling IP Mobility is very similar to scaling Internet Multihoming
 - Mobility moves an “ID” (unique address) from one network “location” to another network “location”
 - Multihoming connects an “ID” (a unique address) to multiple networks “locations” at the same time
 - With both Mobility and Multihoming, the network must keep more specific state “globally” about where something is located

LISP Overview

LISP :: A Routing Architecture – Not a Feature

- Uses pull vs. push routing
 -
 -
- An over-the-top technology
 - Address Family agnostic
 - Incrementally deployable
 - End systems can be unaware of LISP
- Deployment simplicity
 - No host changes
 - Minimal CPE changes
 - Some new core infrastructure components
- LISP use-cases are complimentary
 - Simplified multi-homing with Ingress traffic Engineering; no need for BGP
 - Address Family agnostic support
 - Virtualisation support
 - End-host mobility without renumbering
- Enables IP Number Portability
 - Never change host IP's; No renumbering costs
 - No DNS changes; “name == EID” binding
 - Session survivability
- An Open Standard
 - Being developed in the IETF (RFC 6830-6836)
 - No Cisco Intellectual Property Rights



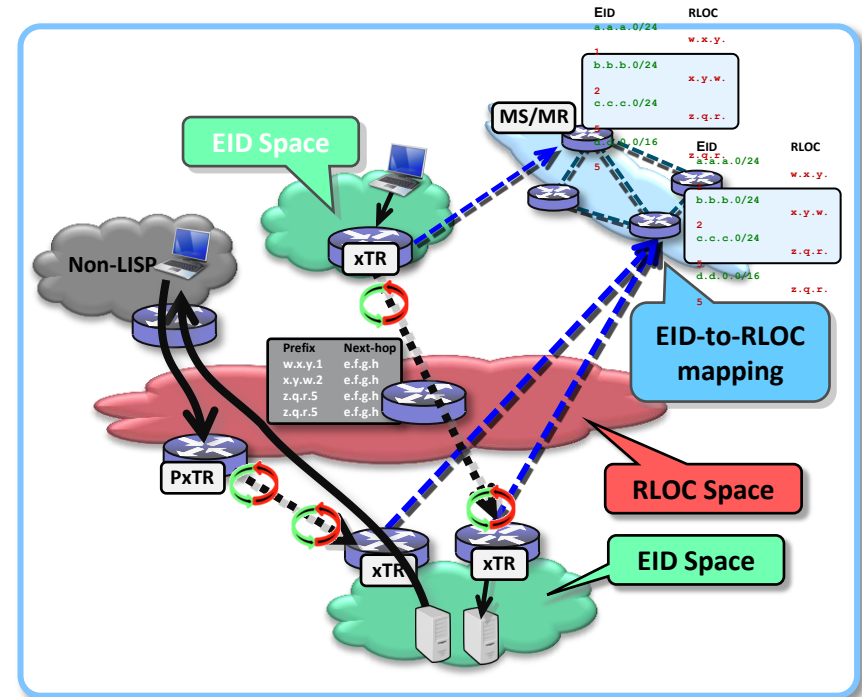
LISP Operations

LISP Operations

Main Attributes of LISP

- LISP namespaces
 - **EID (Endpoint Identifier)**
 - **RLOC (Routing Locator)** is the IP address of the LISP router for the host
 - **EID-to-RLOC mapping** is the distributed architecture that maps **EIDs** to **RLOCs**

- Network-based solution
- No host changes
- Minimal configuration
- No DNS changes
- Address Family agnostic
- Incrementally deployable (support LISP and non-LISP)
- Support for mobility



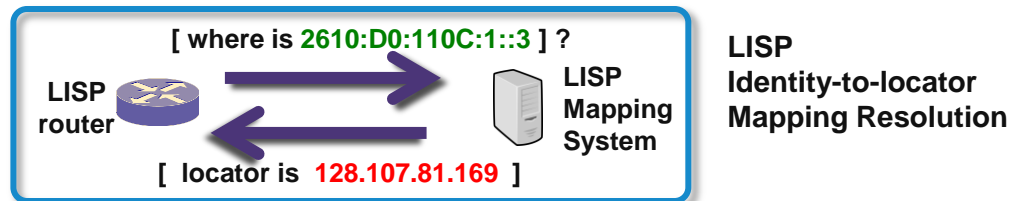
LISP Operations

LISP :: Mapping Resolution “Level of Indirection” DNS Analog

- LISP “Level of Indirection” is analogous to a DNS lookup
 - DNS resolves IP addresses for URL Answering the “WHO IS” question

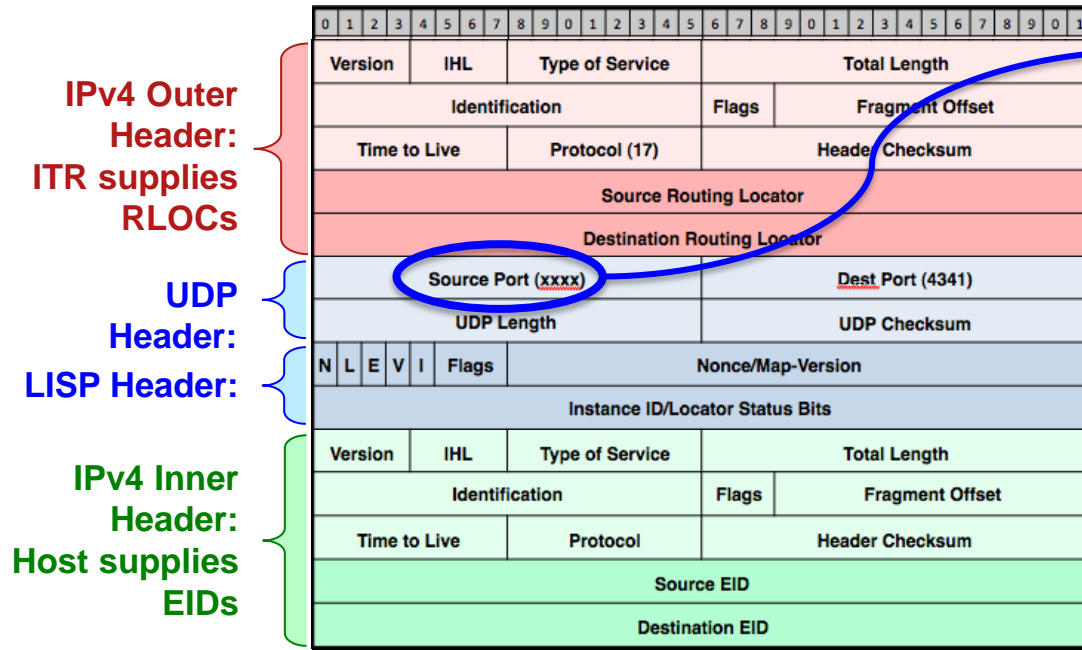


- LISP resolves locators for queried identities Answering the “WHERE IS” question



LISP Operations

LISP IPv4 EID / IPv4 RLOC Data Packet Header Example

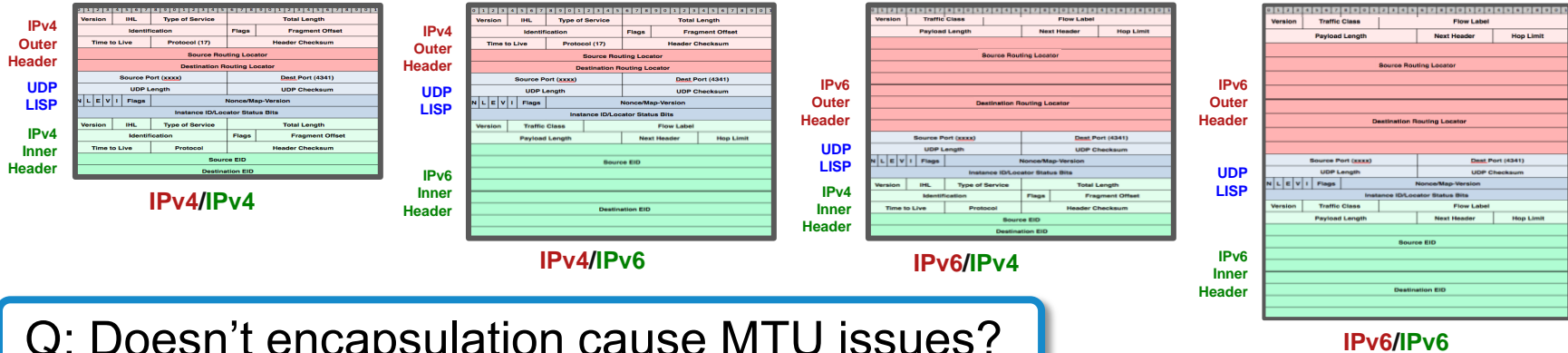


Q: How does the UDP source port get selected ?

A: It's either a 3-tuple or 5-tuple hash of the inner (EID) header

LISP Operations

LISP Encapsulation Combinations – IPv4 and IPv6 Supported



Q: Doesn't encapsulation cause MTU issues?

A: It can... But preparation limits issues...

- Encapsulation overhead is 36B IPv4 and 56B IPv6
- LISP supports “stateful” (PMTUD) and “stateless” (fragmentation) options
- Tunnel/MTU issues are well known (GRE, IPsec, etc.) and are usually operationally tractable

LISP Operations

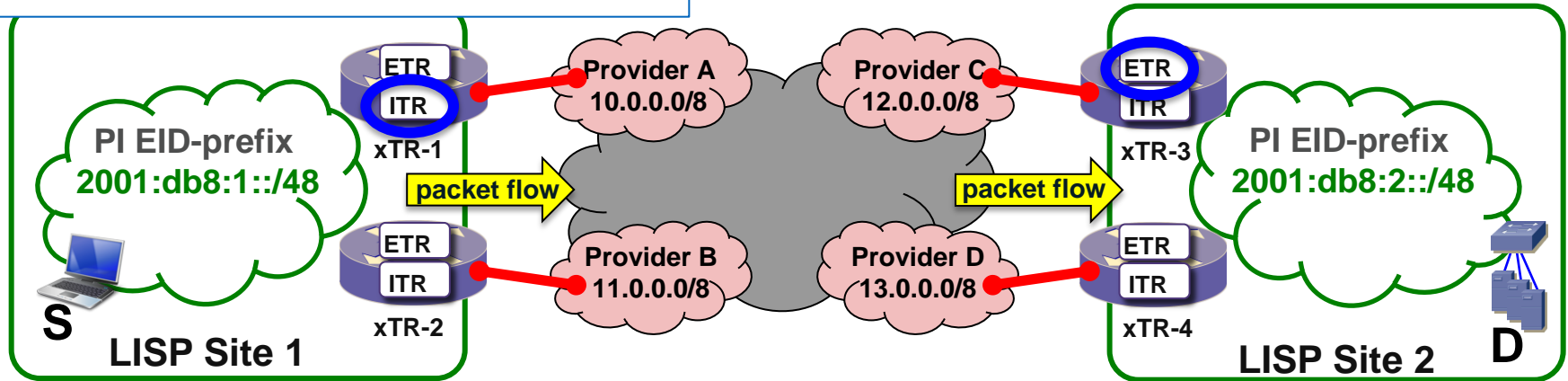
LISP Data Plane :: Ingress/Egress Tunnel Router (xTR)

ITR – Ingress Tunnel Router

- Receives packets from site-facing interfaces
- Encap to remote LISP sites, or native-fwd to non-LISP sites

ETR – Egress Tunnel Router

- Receives packets from core-facing interfaces
- De-cap and deliver packets to local **EIDs** at site



LISP Operations

LISP Data Plane :: Unicast Packet Flow

Map-Cache Entry

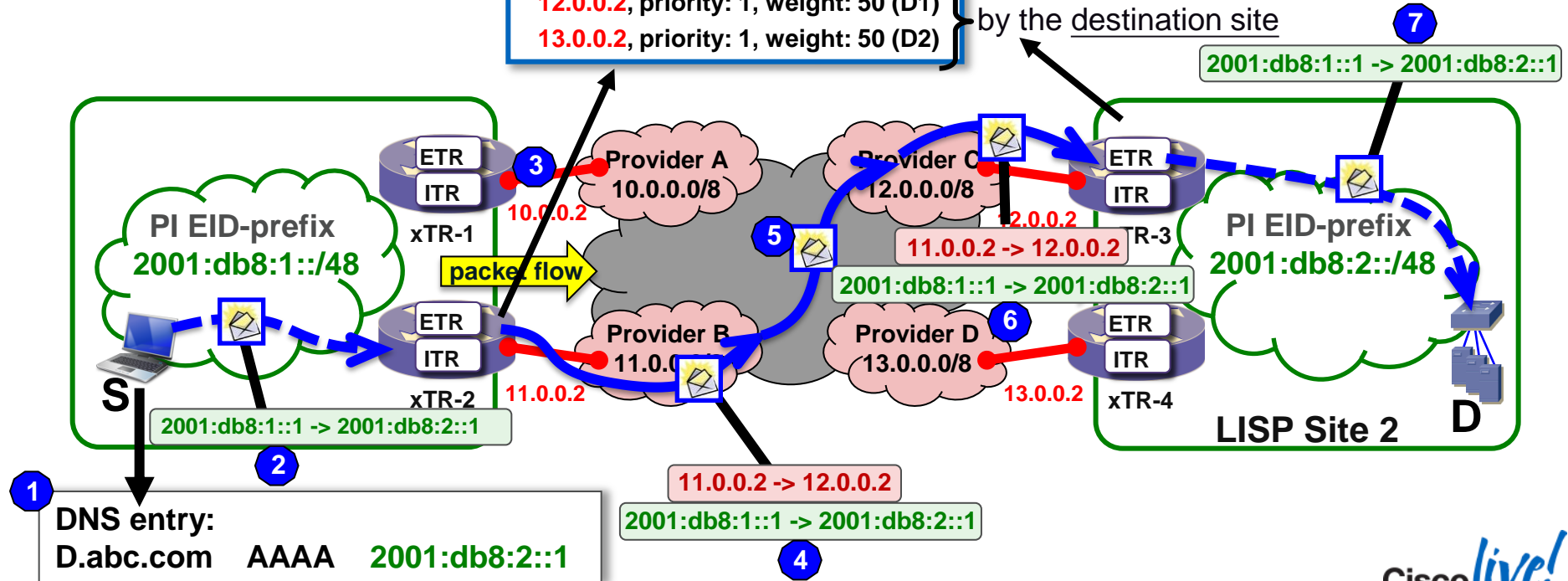
EID-prefix: **2001:db8:2::/48**

Locator-set:

12.0.0.2, priority: 1, weight: 50 (D1)

13.0.0.2, priority: 1, weight: 50 (D2)

This policy controlled
by the destination site

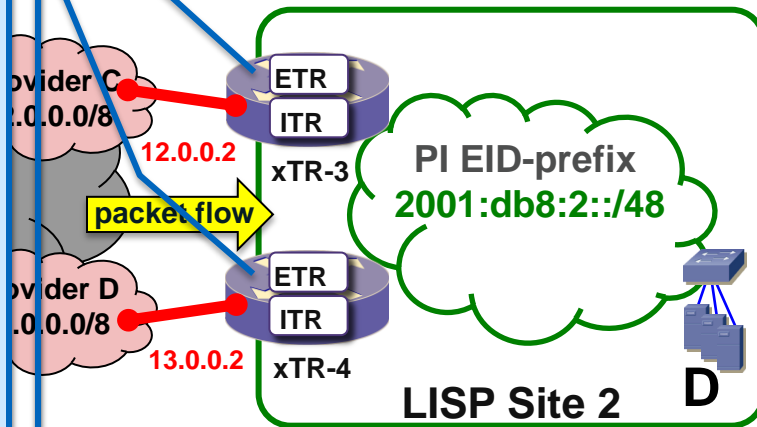


LISP Operations

LISP Data Plane :: Ingress/Egress Tunnel Router (xTR)

```
!  
router lisp  
  locator-set SITE2  
    12.0.0.2 priority 1 weight 50  
    13.0.0.2 priority 1 weight 50  
  exit  
!  
eid-table default instance-id 0  
  database-mapping 2001:db8:2::/48 locator-set SITE2  
  exit  
!  
ipv6 itr map-resolver 66.2.2.2  
ipv6 itr  
ipv6 etr map-server 66.2.2.2 key S3cr3t-2  
ipv6 etr  
exit  
!  
ip route 0.0.0.0 0.0.0.0 12.0.0.1  
ip route 0.0.0.0 0.0.0.0 13.0.0.1  
!
```

Identical config on both xTRs!



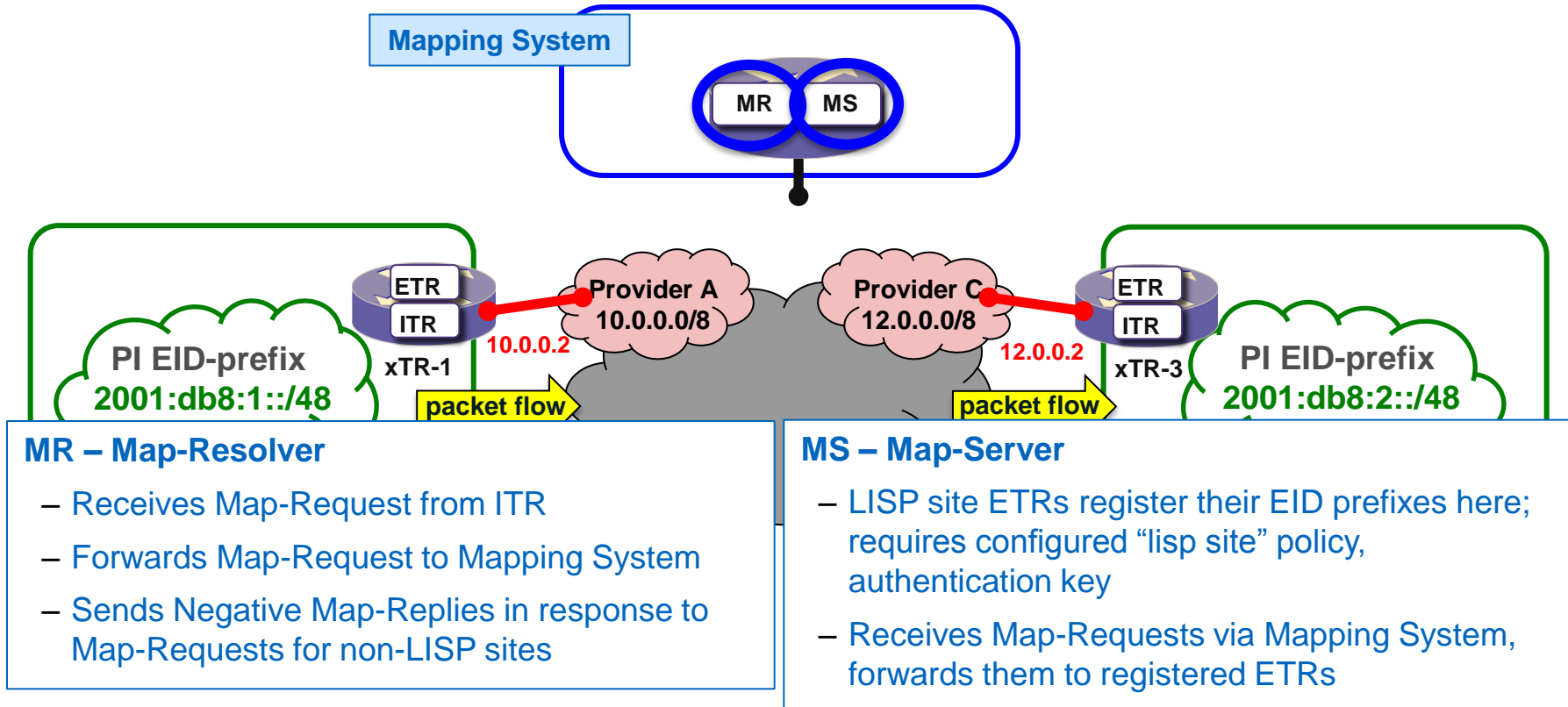
LISP Operations

LISP Control Plane :: Introduction...

- LISP Control Plane Provides On-Demand Mappings
 - Control Plane is separate from the Data Plane (UDP 4342 vs UDP 4341)
 - Map-Resolver and Map-Server (similar to DNS Resolver and DNS Server)
 - LISP Control Plane Messages for EID-to-RLOC
 - Distributed databases and map-caches hold mappings

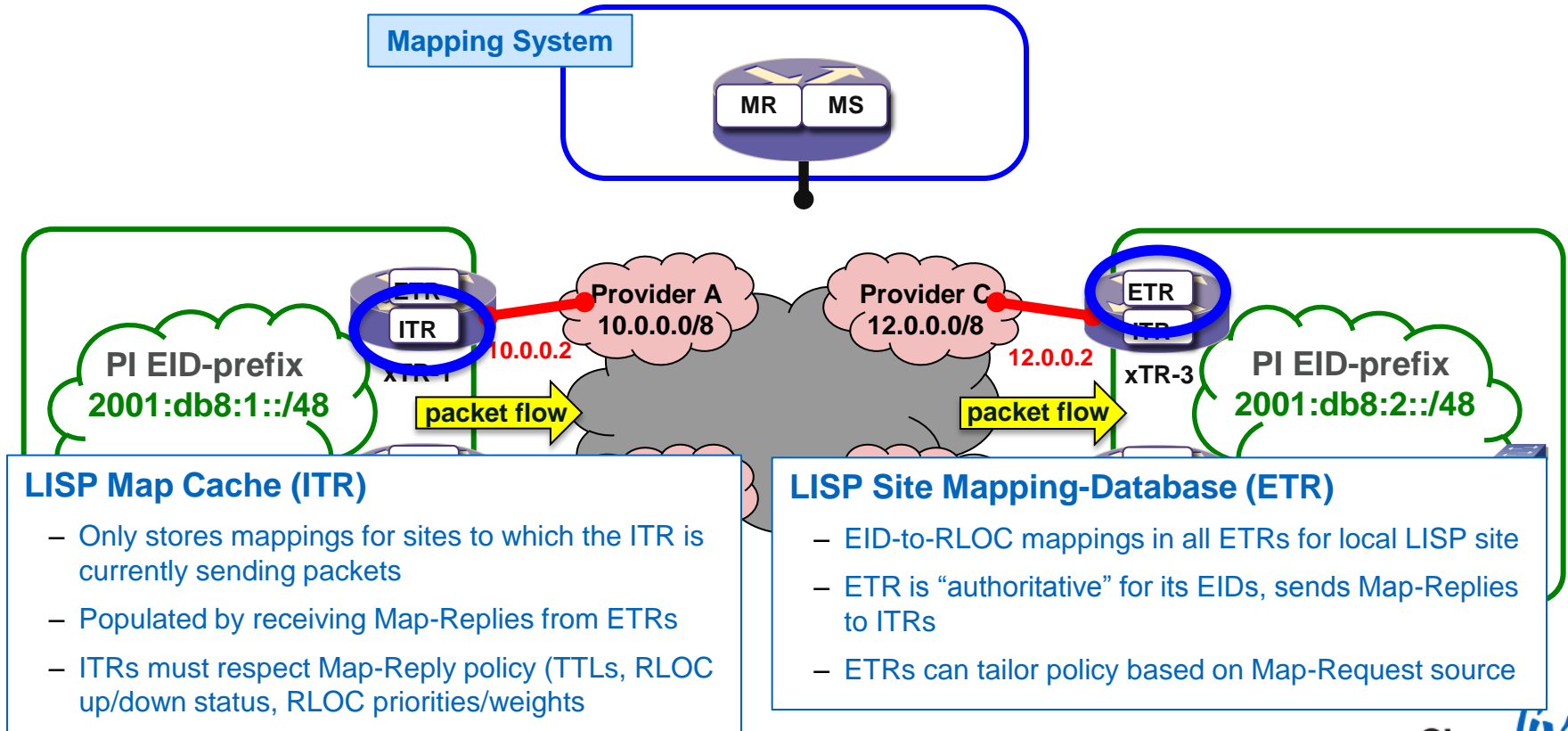
LISP Operations

LISP Control Plane :: Map-Server/Map-Resolver (MS/MR)



LISP Operations

LISP Control Plane :: Map-Server/Map-Resolver (MS/MR)



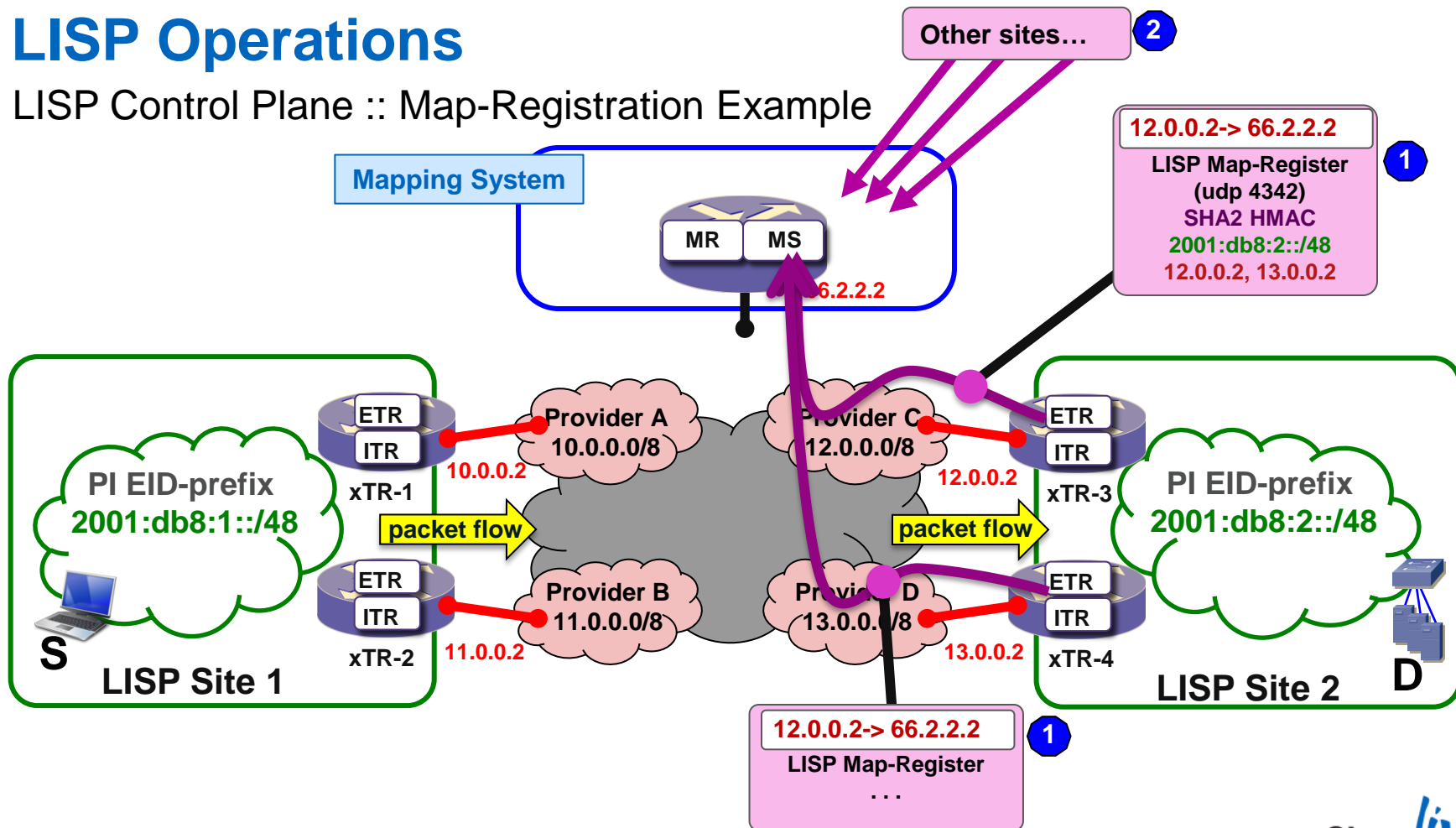
LISP Operations

LISP Control Plane :: Control Plane Messages...

- Control Plane Control Plane **EID** Registration
 - **Map-Register message**
 - Sent by ETR to MS to register its associated **EID** prefixes
 - Specifies the **RLOC(s)** to be used by the MS when forwarding Map-Requests to the ETR
- Control Plane “Data-triggered” mapping services
 - **Map-Request message**
 - Sent by an ITR when it needs an **EID/RLOC** mapping, to test an **RLOC** for reachability, to refresh a mapping before TTL expiration, or in response to a Solicit Map-Request (SMR)
 - **Map-Reply message**
 - Sent by an ETR in response to a valid map-request to provide the **EID/RLOC** mapping and site ingress policy for the requested **EID**
 - **Map-Notify message**

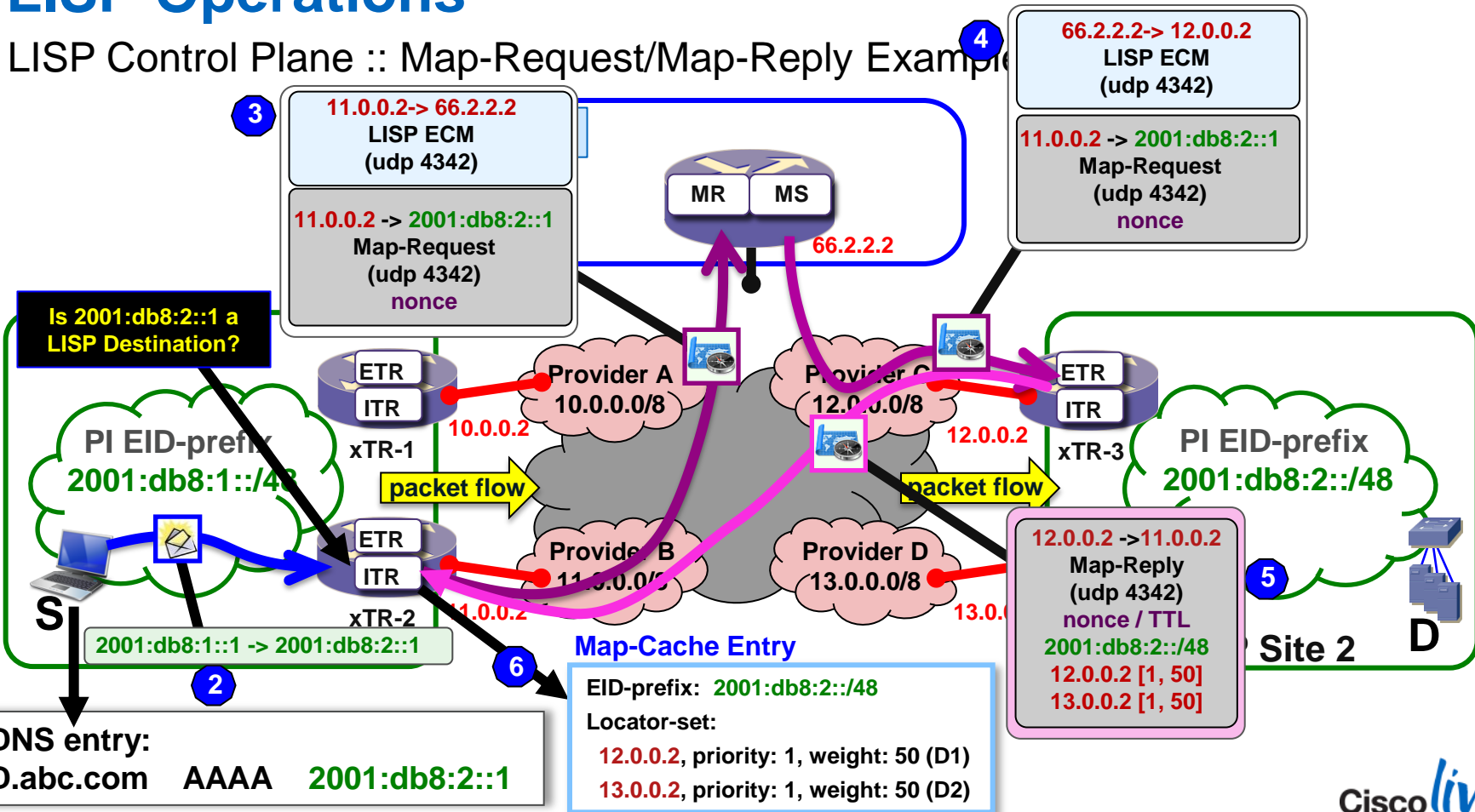
LISP Operations

LISP Control Plane :: Map-Registration Example



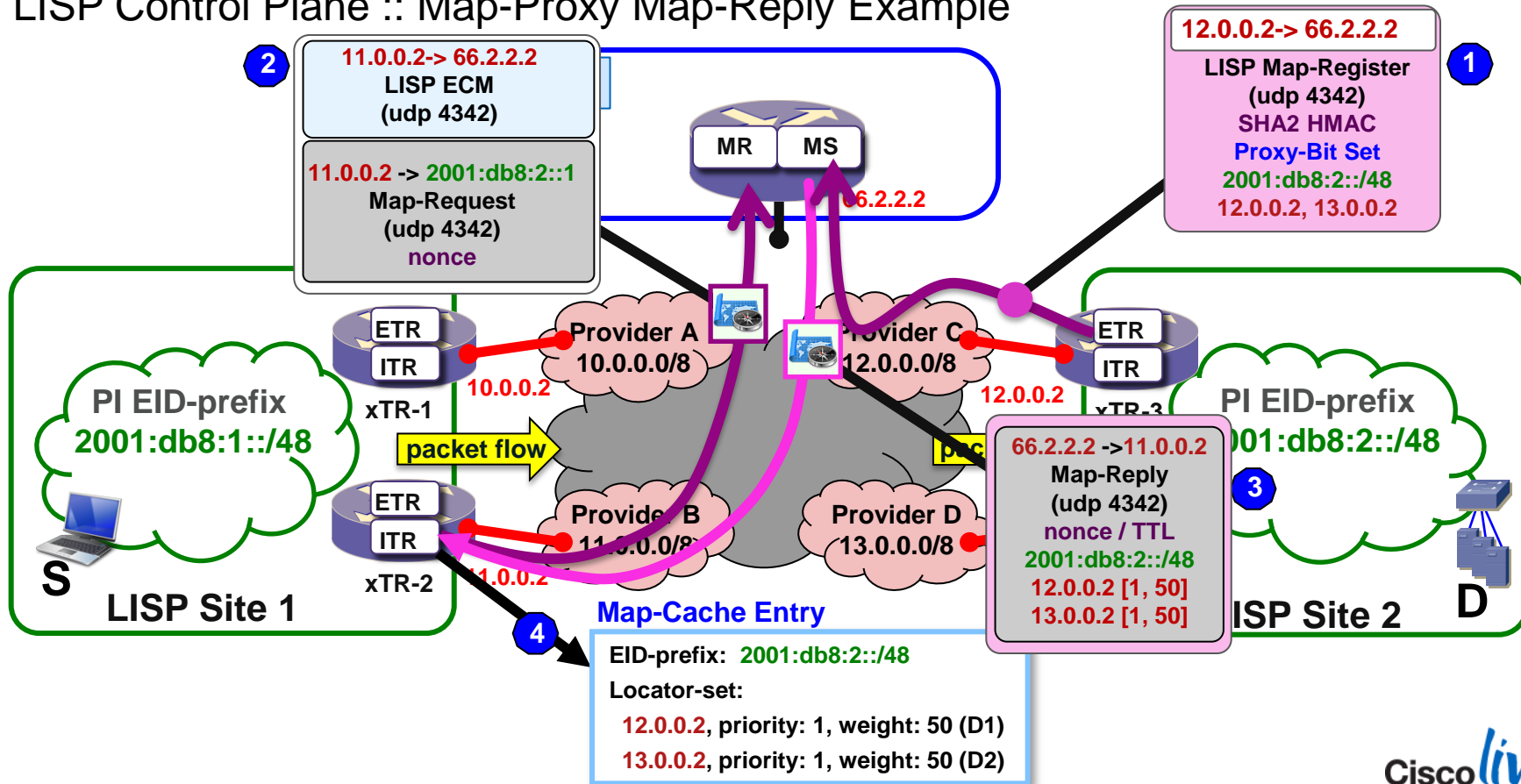
LISP Operations

LISP Control Plane :: Map-Request/Map-Reply Example



LISP Operations

LISP Control Plane :: Map-Proxy Map-Reply Example

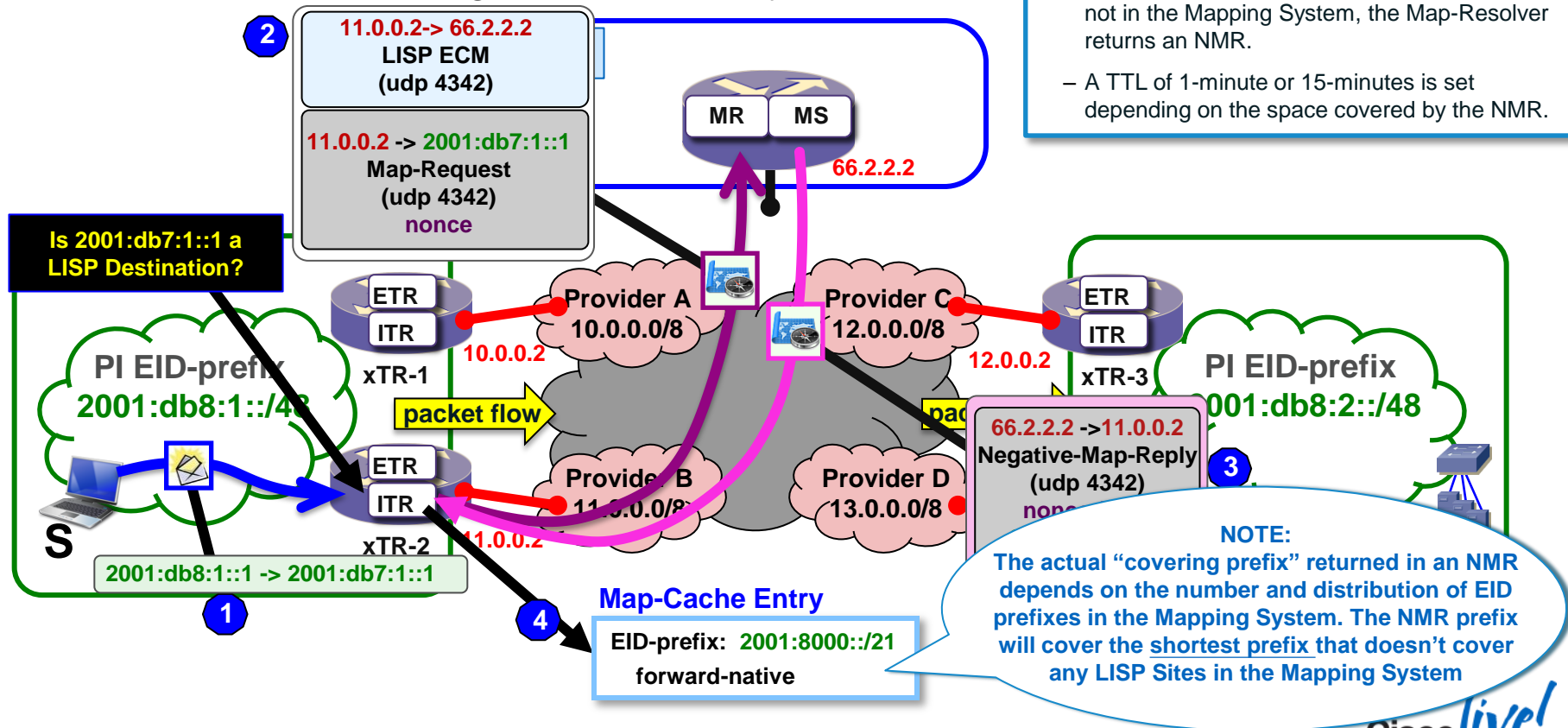


LISP Operations

LISP Control Plane :: Negative Map-Reply Example

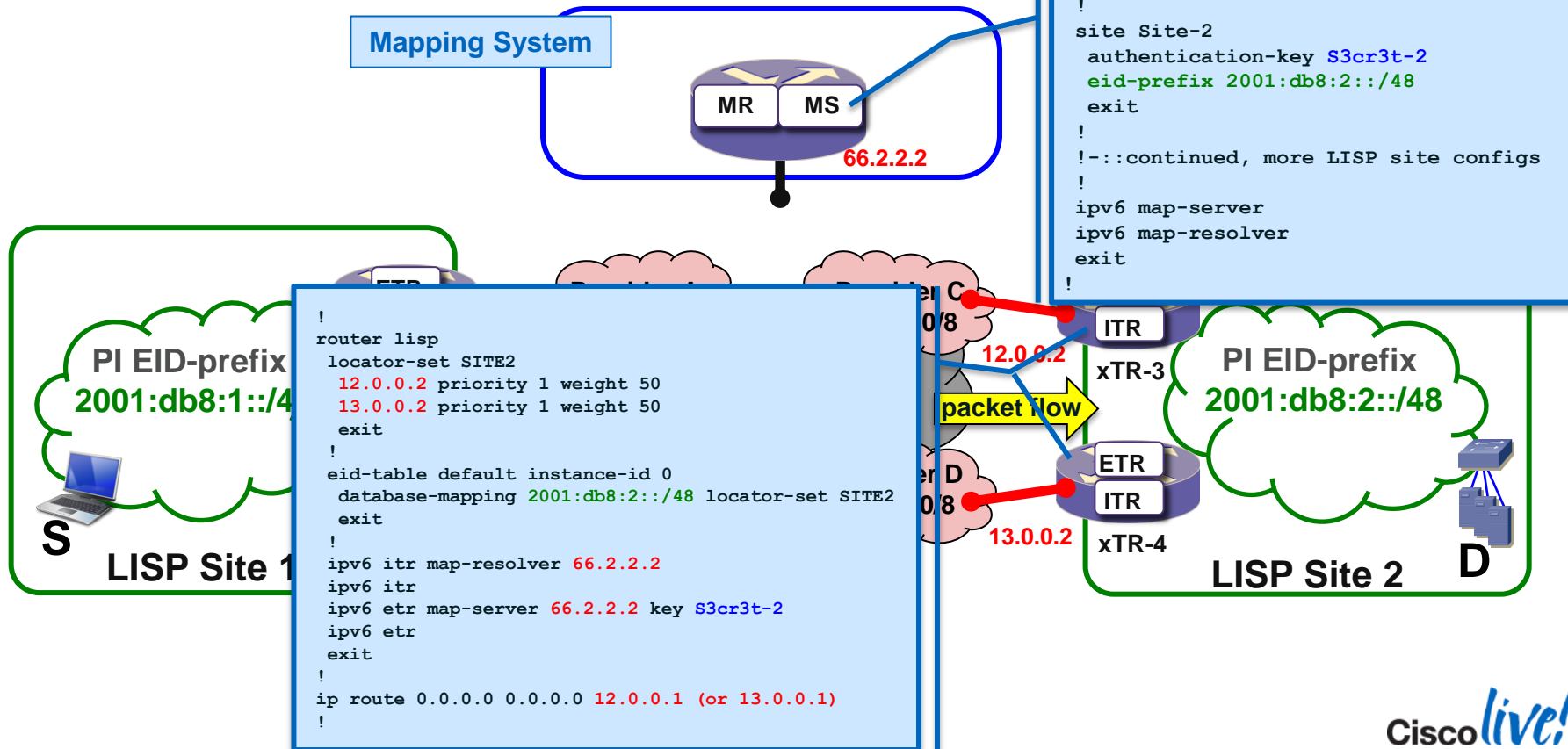
Notes:

- When an ITR queries for a destination that is not in the Mapping System, the Map-Resolver returns an NMR.
- A TTL of 1-minute or 15-minutes is set depending on the space covered by the NMR.



LISP Operations

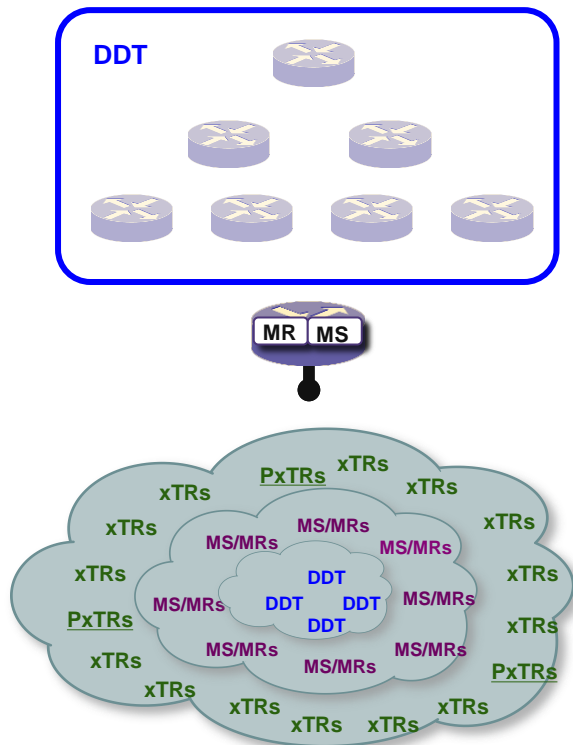
LISP Control Plane :: Map-Server/Map-Resolver (MS/MR)



LISP Operations

The LISP Beta Network uses DDT today...

LISP Control Plane :: Mapping System Scaling...



Scaling the LISP Mapping System

- Deploy multiple “stand-alone” Map-Servers” and register each LISP Site to all of them (up to eight)
- Deploy Map-Resolvers in an “Anycast” manner
- Or, deploy a “hierarchical” Mapping System - DDT

DDT – Delegated Distributed Tree

- Hierarchy for Instance IDs and for EID Prefixes
- DDT Map-Resolvers sends (ECM) Map-Requests
- DDT Nodes Return Map-Referral messages
- DDT Resolvers resolve the Map-Server’s RLOC iteratively
- Conceptually, similar to DNS (IN-ADDR hierarchy) but different prefix encoding, messages, etc.

LISP Deployment Overview

Private and Public LISP Deployment Models...

Private Model

- “Private” LISP deployment support single Enterprises or Entities
- LISP Enterprise deploys:
 - xTRs
 - Mapping System, if required
 - Proxy System, if required

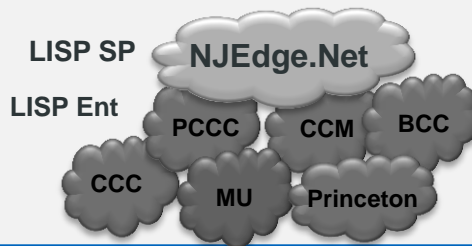
Private Enterprise Examples



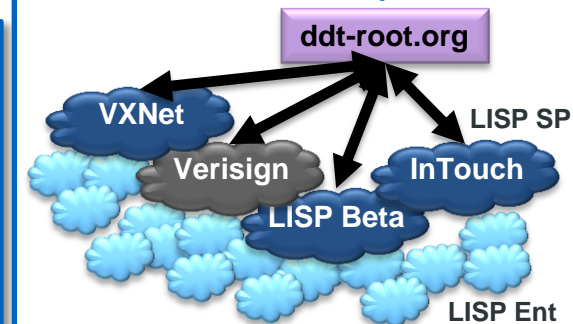
Public Model

- “Public” LISP deployment supports the needs of multiple Enterprises
- LISP Service Provider deploys “shared” Mapping System and Proxy System
- LISP Enterprises subscribe to LISP SP, and deploy their own xTRs

Stand-Alone Example



Global Examples



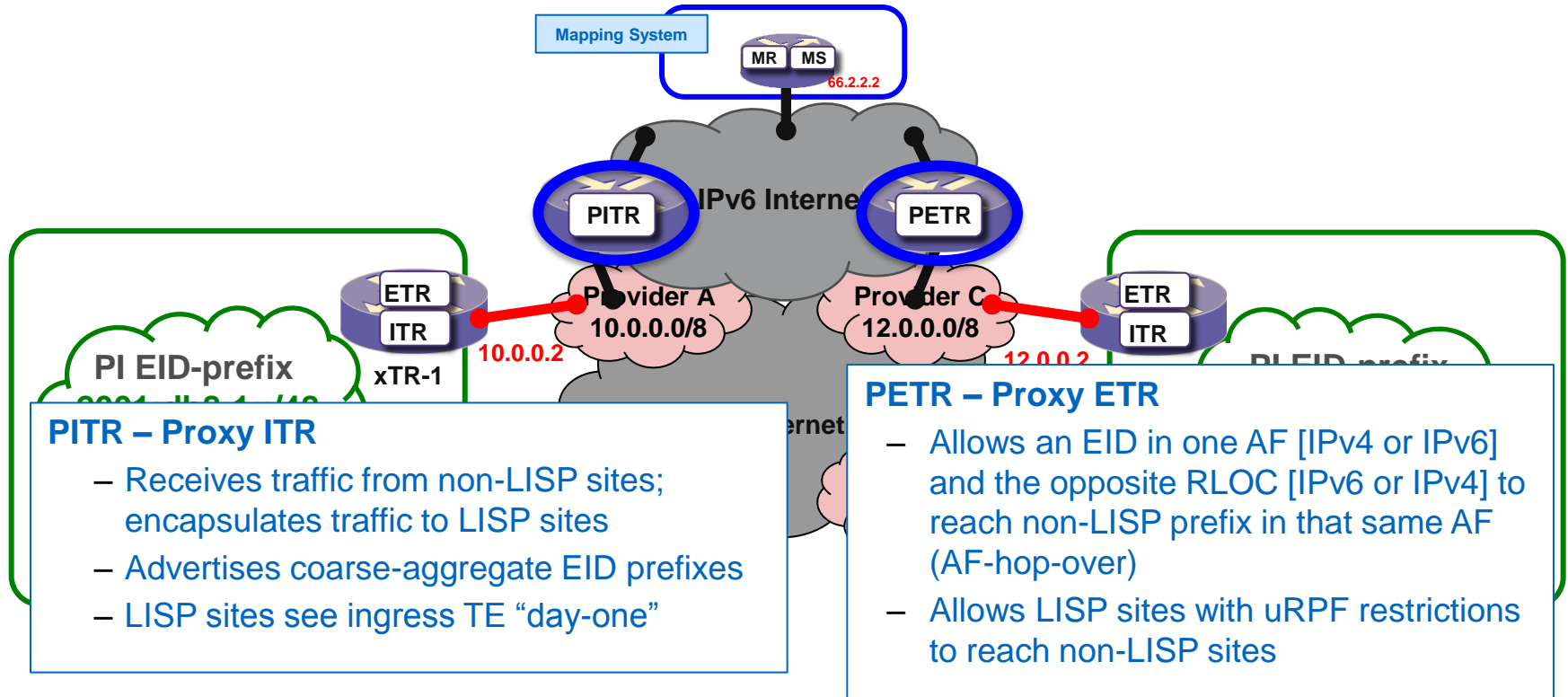
LISP Operations

LISP Internetworking :: Day-One Incremental Deployment

- Early Recognition
 - Up-front recognition of an incremental deployment plan
 - LISP will not be widely deployed day-one
- Interworking for:
 - **LISP-sites** to **non-LISP sites** (e.g. the rest of the Internet)
 - **non-LISP sites** to **LISP-sites**
- Proxy-ITR/Proxy-ETR are deployed today
 - Infrastructure LISP network entity
 - Creates a monetised service opportunity for infrastructure players

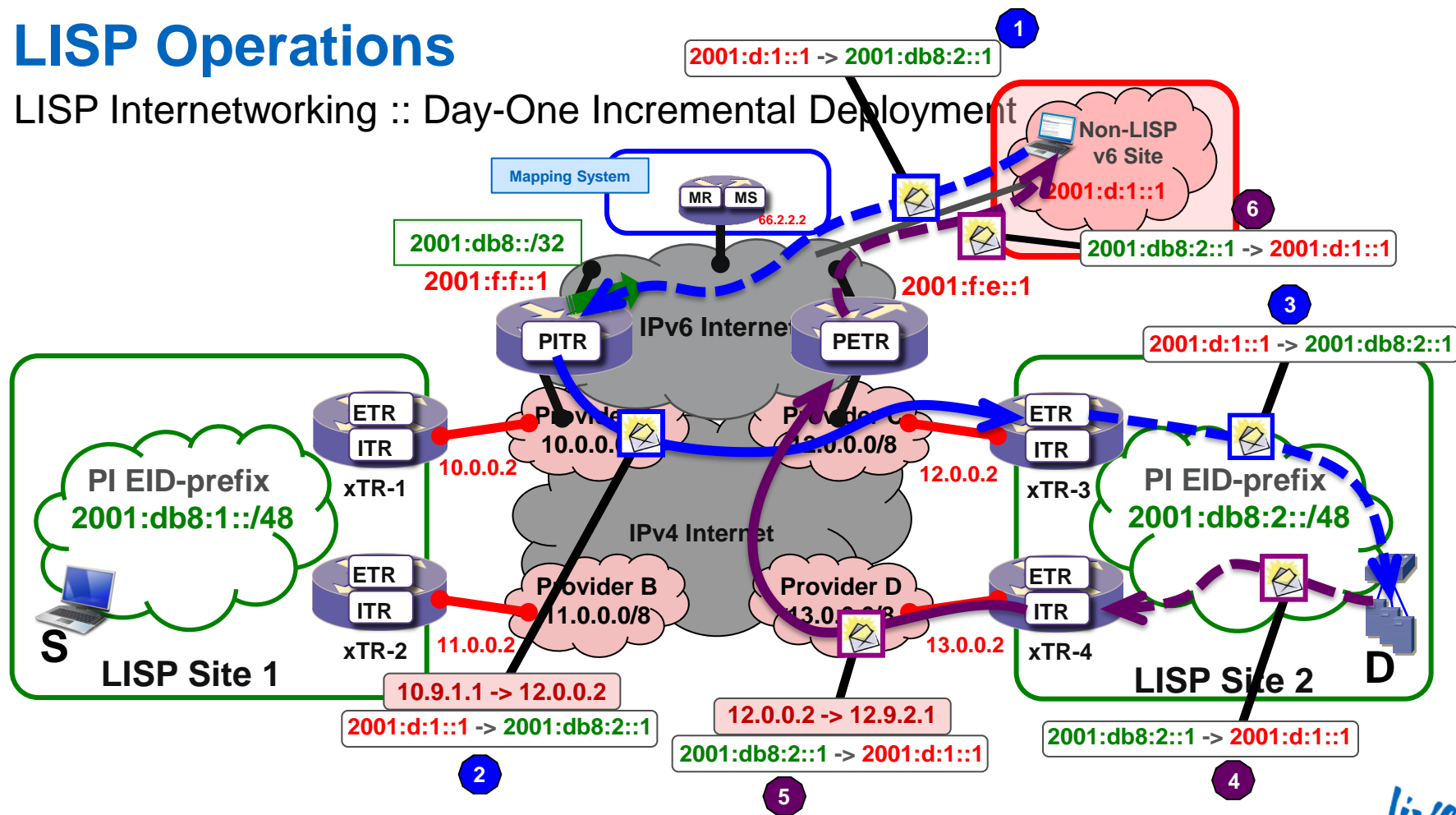
LISP Operations

LISP Internetworking :: Day-One Incremental Deployment



LISP Operations

LISP Internetworking :: Day-One Incremental Deployment





LISP Deployment Examples

LISP Deployment Examples

LISP Deployment Examples...

1. Efficient Multihoming and Multi-AF (IPv4 and IPv6)
2. Efficient Virtualisation and High-Scale VPNs
3. Data Centre/Host Mobility
4. LISP-Mobile Node



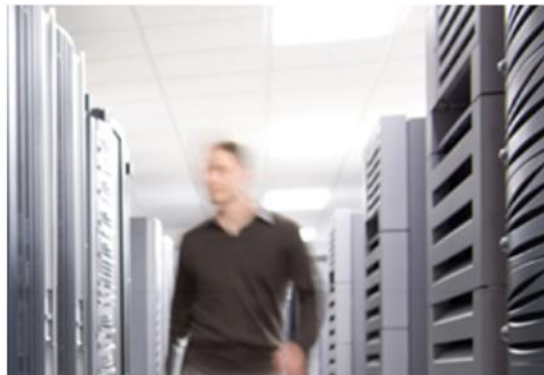
These examples highlight functionality integrated in LISP.

All use-case – multi-homing, v6 transition, virtualisation, and mobility work together!

LISP Deployment Examples

LISP Deployment Examples...

1. Efficient Multihoming and Multi-AF (IPv4 and IPv6)
2. Efficient Virtualisation and High-Scale VPNs
3. Data Centre/Host Mobility
4. LISP-Mobile Node



LISP Multihoming/Multi-AF Details

LISP Operations

LISP Encapsulation – Any IPv4 and IPv6 Combination Supported

IPv4
Outer
Header

0	1	2	3	4	5	6	7	8	9	0	1	2	3	4	5	6	7	8	9	0	1	2	3	4	5	6	7	8	9	0	1
Version				IHL				Type of Service				Total Length																			
Identification								Flags		Fragment Offset																					
Time to Live				Protocol (17)				Header Checksum																							
Source Routing Locator																															
Destination Routing Locator																															

0	1	2	3	4	5	6	7	8	9	0	1	2	3	4	5	6	7	8	9	0	1	2	3	4	5	6	7	8	9	0	1
Version				Traffic Class				Flow Label																							
Payload Length								Next Header				Hop Limit				Source Routing Locator															
Destination Routing Locator																															

IPv6
Outer
Header

UDP
LISP

0	1	2	3	4	5	6	7	8	9	0	1	2	3	4	5	6	7	8	9	0	1	2	3	4	5	6	7	8	9	0	1
Source Port (xxxx)										Dest Port (4341)																					
UDP Length										UDP Checksum																					
N	L	E	V	I	Flags		Nonce/Map-Version																								
Instance ID/Locator Status Bits																															

IPv6/IPv6
IPv6/IPv4
IPv4/IPv4
IPv4/IPv6

IPv6
Inner
Header

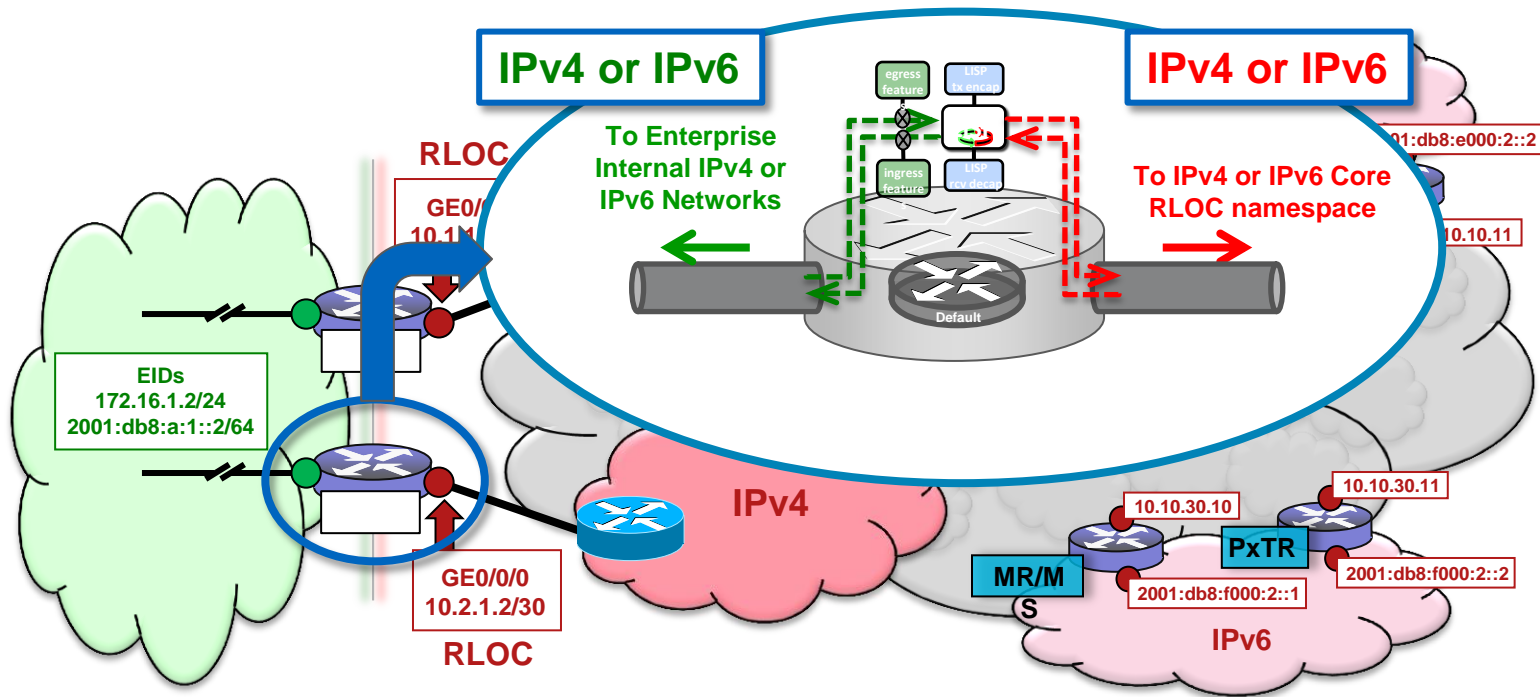
0	1	2	3	4	5	6	7	8	9	0	1	2	3	4	5	6	7	8	9	0	1	2	3	4	5	6	7	8	9	0	1
Version				Traffic Class				Flow Label																							
Payload Length								Next Header				Hop Limit				Source EID															
Destination EID																															
payload																															

0	1	2	3	4	5	6	7	8	9	0	1	2	3	4	5	6	7	8	9	0	1	2	3	4	5	6	7	8	9	0	1
Version				IHL				Type of Service				Total Length																			
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payload																															

IPv4
Inner
Header

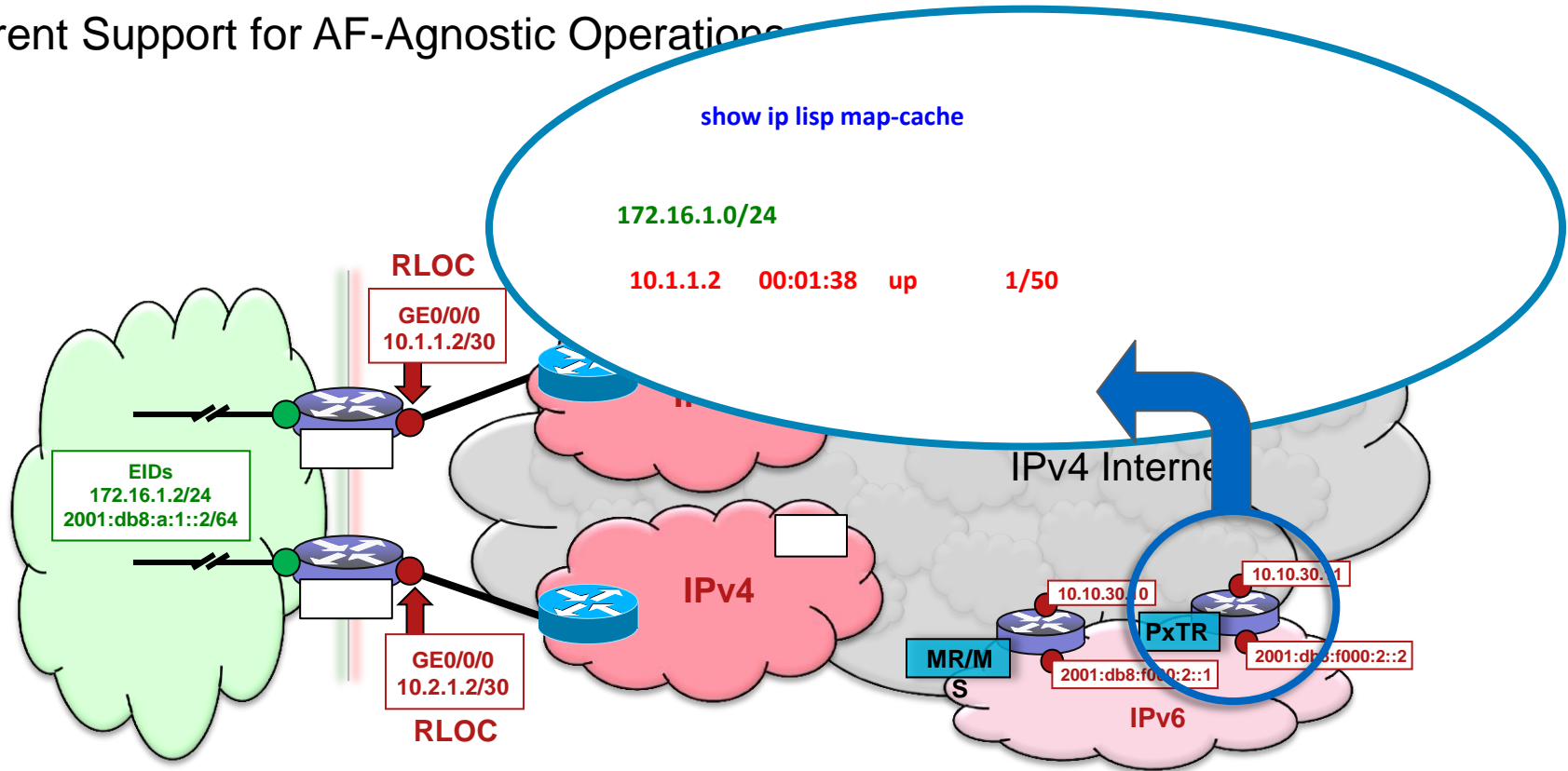
LISP Multihoming and Multi-AF

Inherent Support for AF-Agnostic Operations



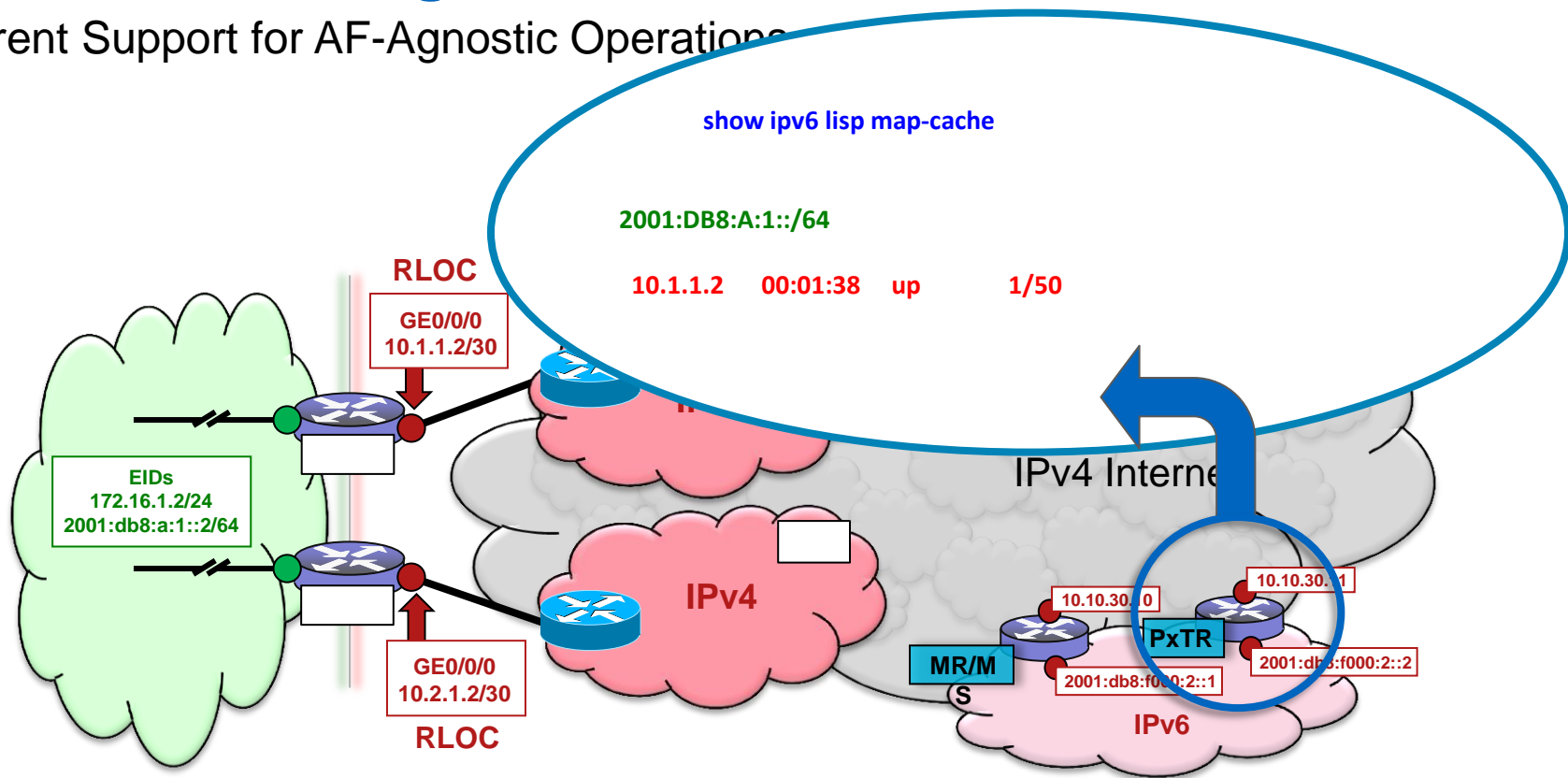
LISP Multihoming and Multi-AF

Inherent Support for AF-Agnostic Operations



LISP Multihoming and Multi-AF

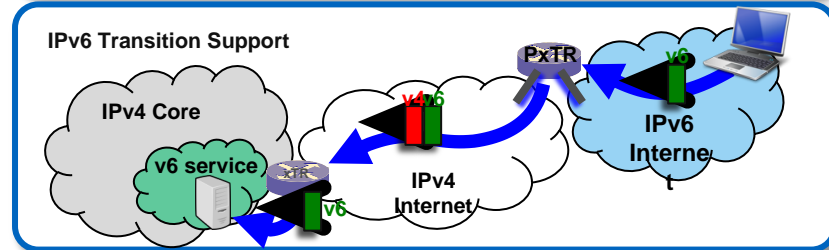
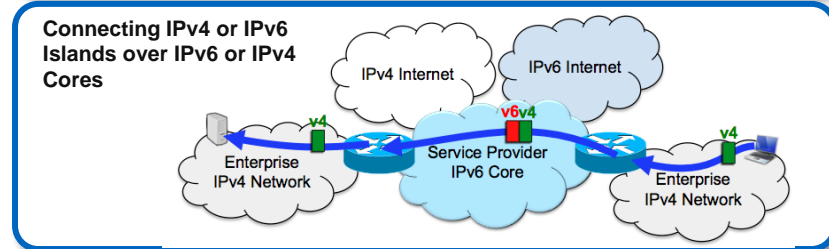
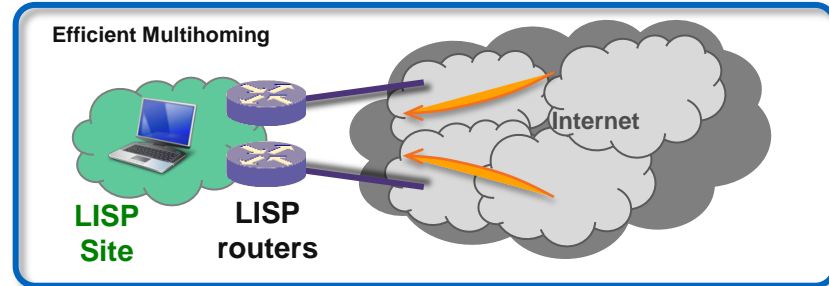
Inherent Support for AF-Agnostic Operations



LISP Multihoming and Multi-AF

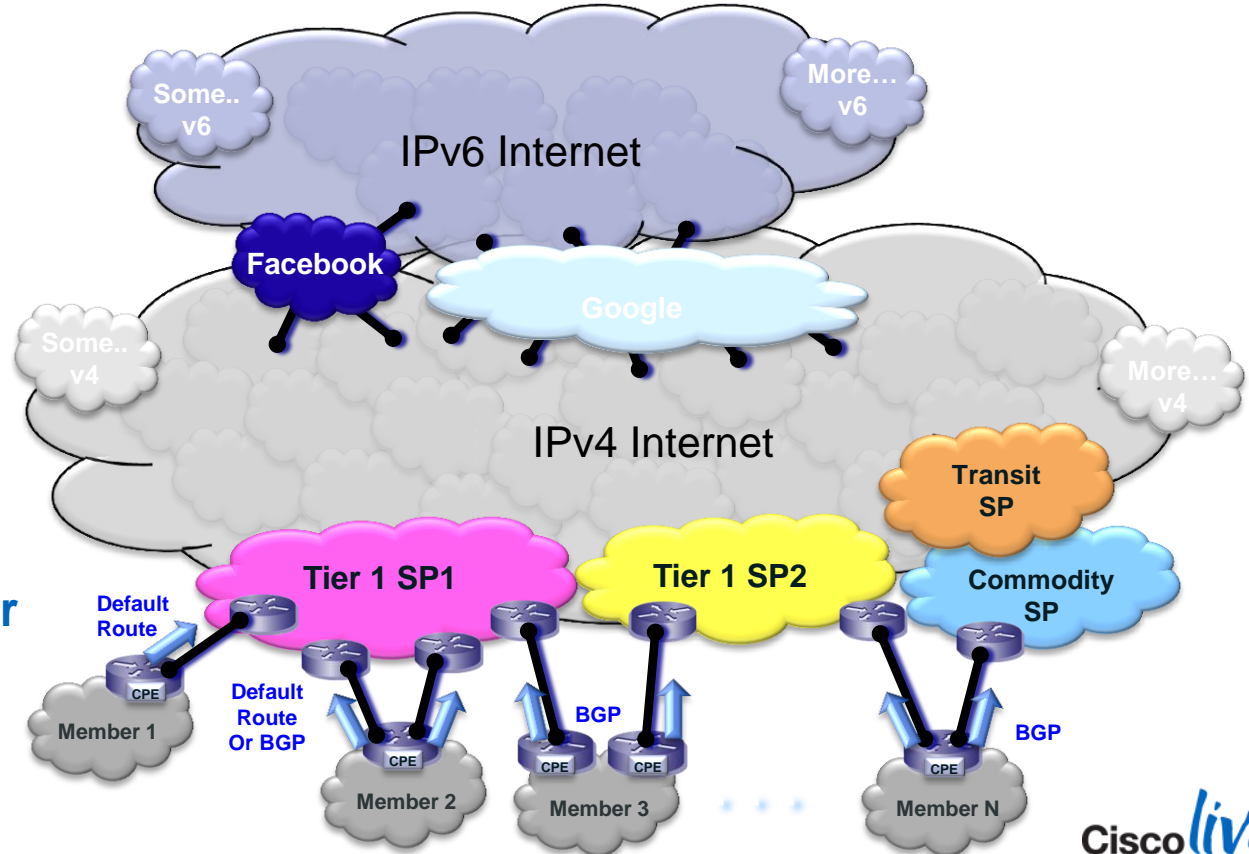
Efficient Multi-Homing and Multi-AF Support

- Needs:
 - Site connectivity to multiple providers for resiliency
 - Low OpEx/CapEx solution for Ingress TE
 - Rapid IPv6 deployment, minimal disruption
- LISP Solution:
 - LISP provides a streamlined solution for handling multi-provider connectivity and policy without BGP complexities
 - LISP encapsulation is Address Family agnostic, allowing for IPv6 over an IPv4 core, or IPv4 over an IPv6 core
- Benefits:
 - OpEx-friendly multi-homing across different providers
 - Simple policy management
 - Ingress Traffic Engineering that actually “works”
 - Minimal configuration
 - No core network changes



LISP Multihoming and Multi-AF

Efficient Multi-Homing and Multi-AF -- Customer Example



Constituent Member Topologies...

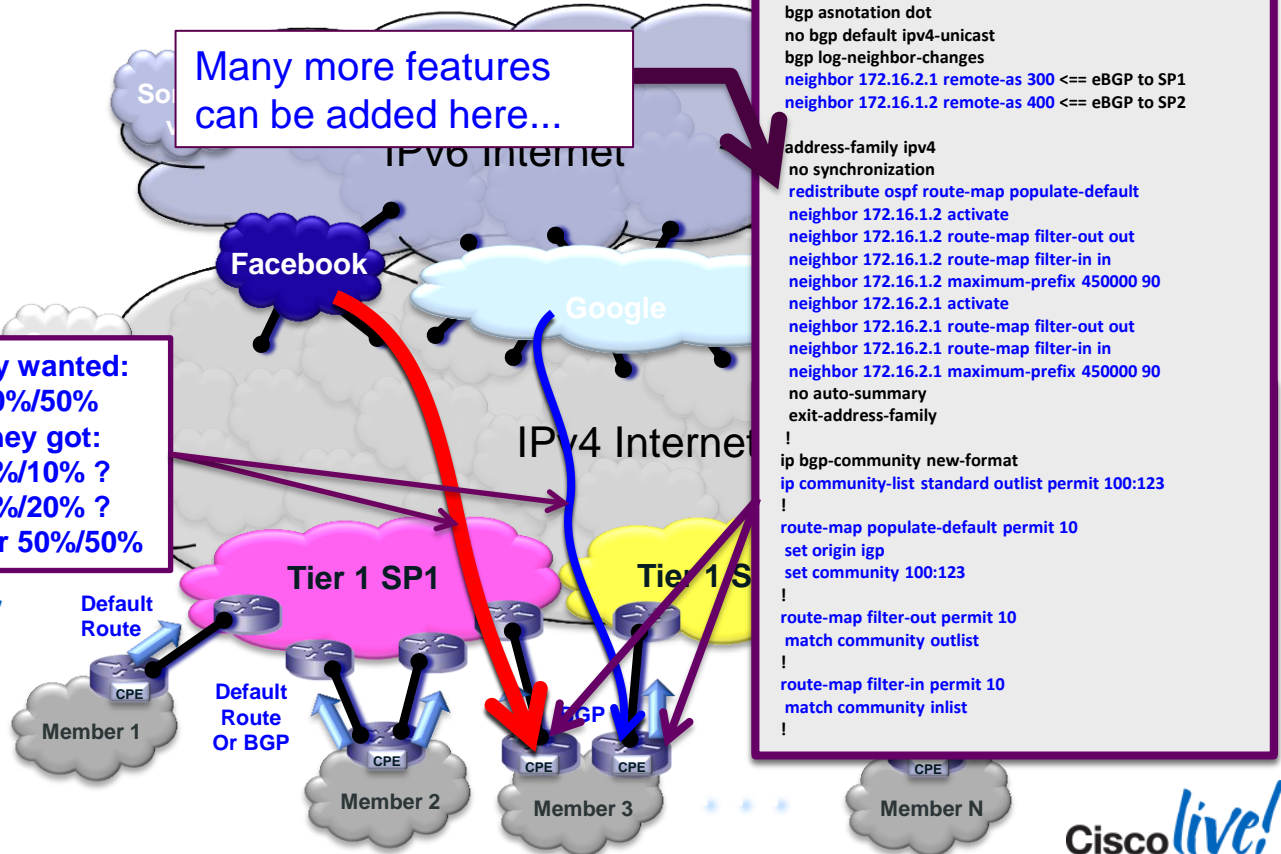
LISP Multihoming and Multi-AF

Efficient Multi-Homing and Multi-AF -- Customer Example

- Configuration complexity...
- Uneven multihoming load shares...

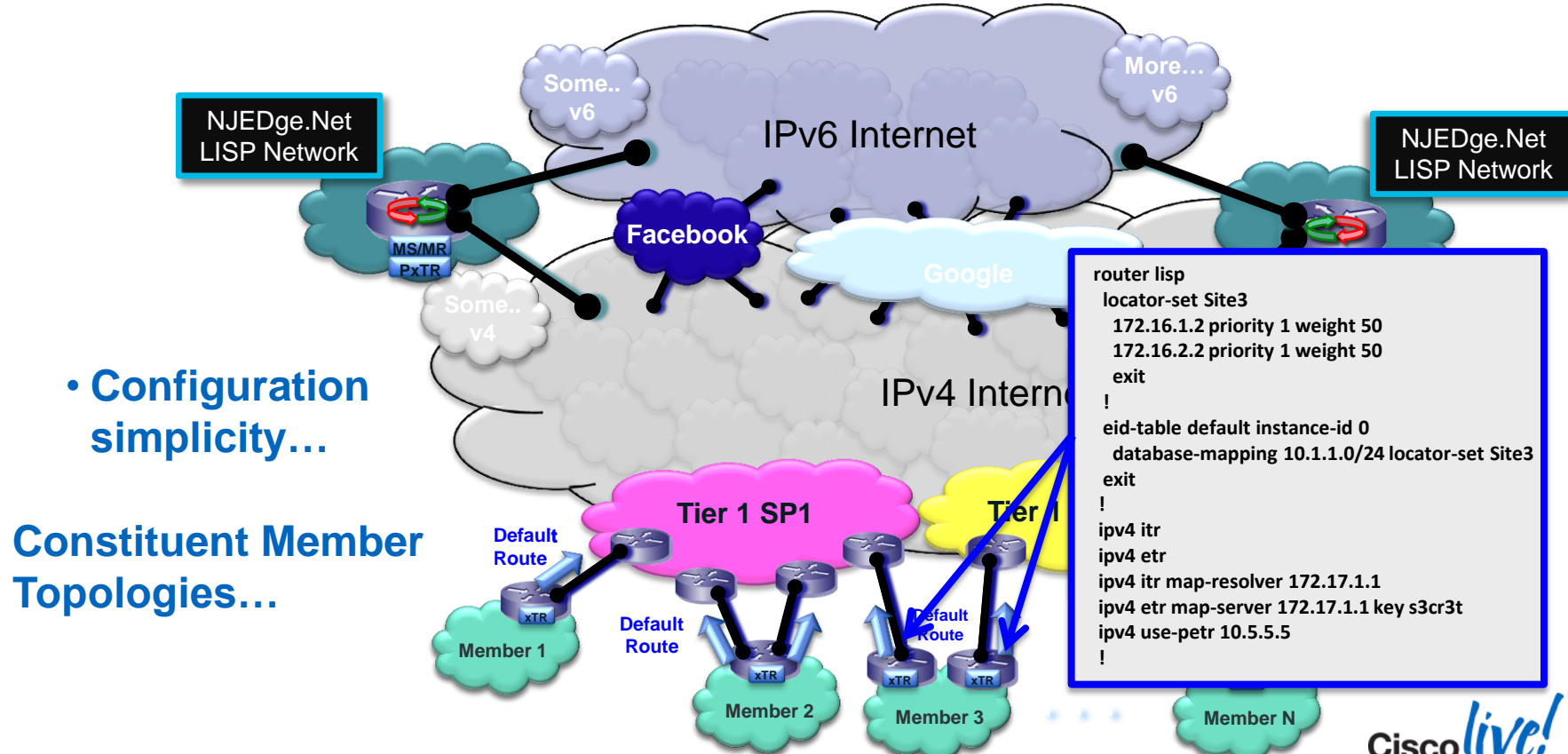
They wanted:
50%/50%
They got:
90%/10% ?
80%/20% ?
Never 50%/50%

Constituent Member Topologies...



LISP Multihoming and Multi-AF

Efficient Multi-Homing and Multi-AF -- Customer Example

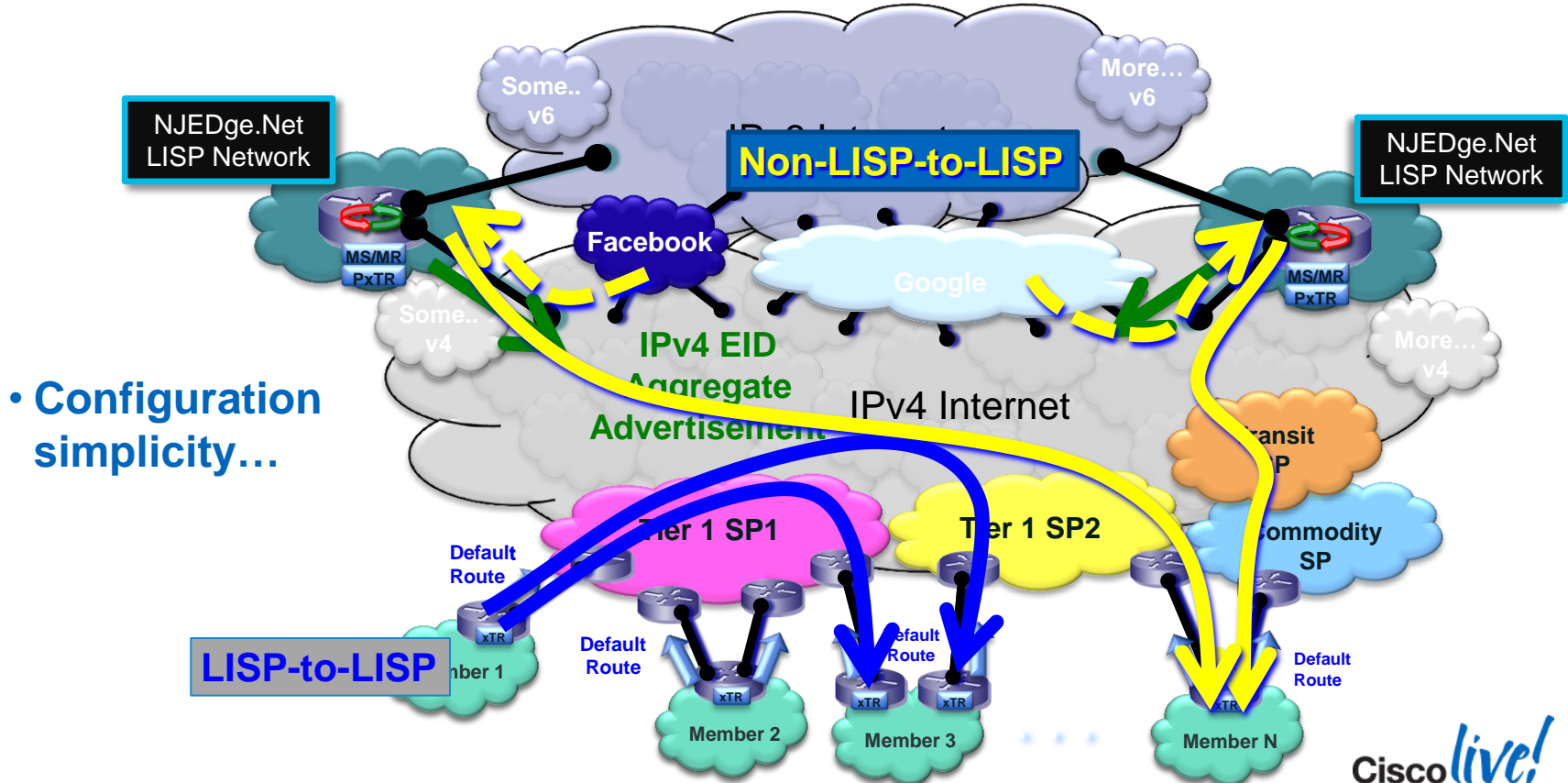


• Configuration simplicity...

Constituent Member Topologies...

LISP Multihoming and Multi-AF

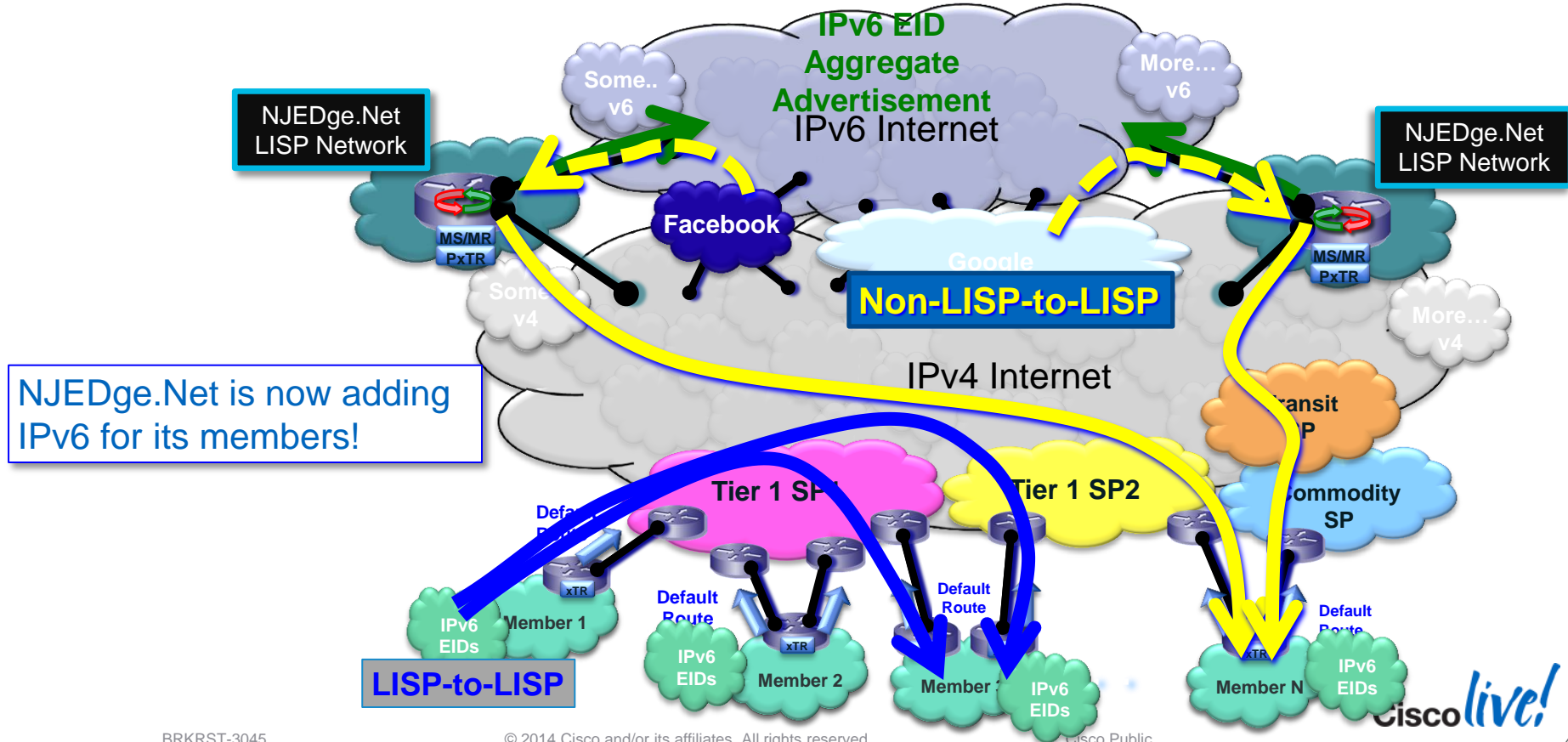
Efficient Multi-Homing and Multi-AF -- Customer Example



• Configuration simplicity...

LISP Multihoming and Multi-AF

Efficient Multi-Homing and Multi-AF -- Customer Example



LISP Use Cases :: Multihoming and Multi-AF

Customer Example :: NJEDge.Net

Key NJEDge.Net LISP Equipment

- ✓ ASR1Ks as MSMRs
- ✓ ASR9Ks as PxTRs (90G Internet capacity)

Key LISP Benefits

- ✓ No BGP to configure or manage
- ✓ No complex configurations
- ✓ Optimised Ingress load balancing
- ✓ Cost Savings by reducing OPEX and CAPEX
- ✓ LISP offers non disruptive transition approach which does not affect end system and allows for incremental deployment
- ✓ Disaster Recovery for Critical Applications introduces Increased Complexity

LISP Multihoming and Multi-AF

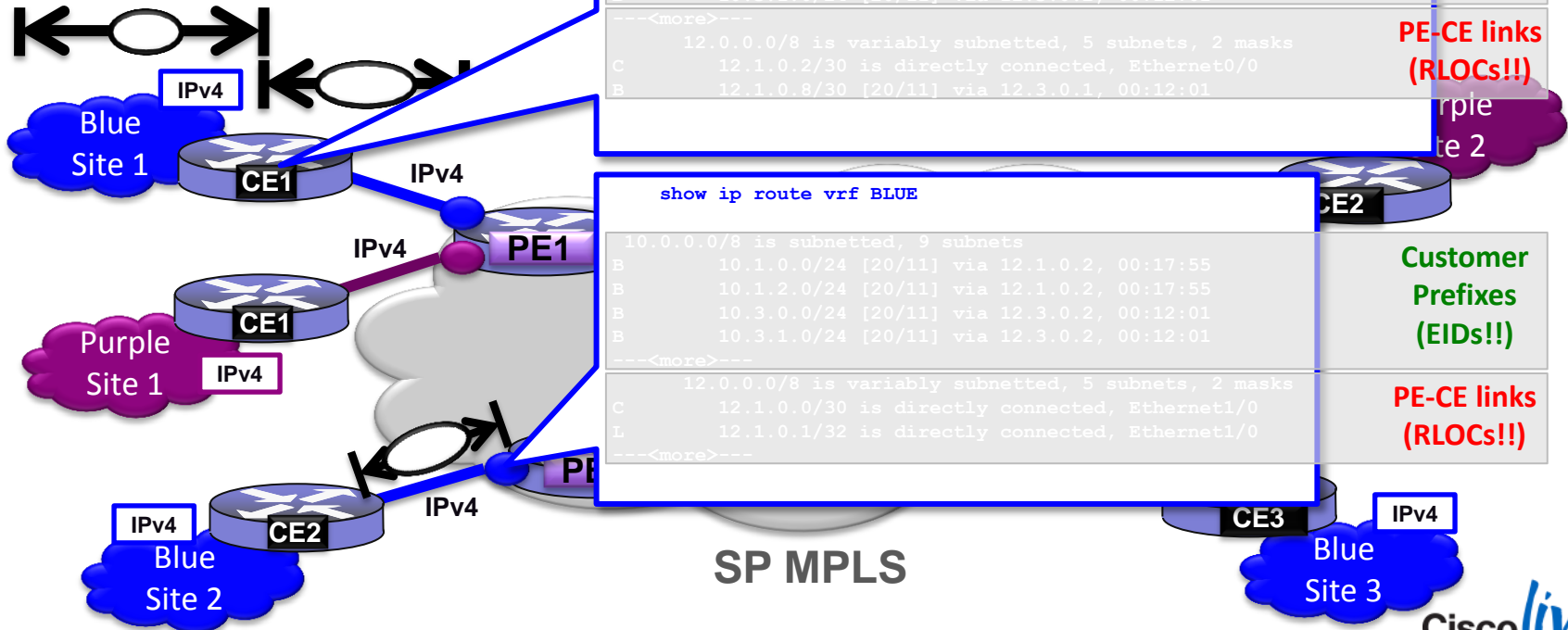
LISP and MPLS Integration

- LISP / MPLS results in an “ideal” deployment environment
 - Locator/ID split idealises a pure “RLOC core” and “EID overlay”
- Opportunities
 - IPv4 over MPLS via LISP
 - Use of LISP (v4-over-v4) removes Customer IPv4 Prefixes from MPLS
 - PE benefits ::
 - (a) (substantially) improved scaling
 - (b) reduced CPU load due to customer route advertisement churn
 - IPv6 over MPLS via LISP
 - Use of LISP (v6-over-v4) removes SP from Customer IPv6 config/mgmt
 - Immediate support :: even if not running LISP for IPv4
 - PE benefits ::
 - (a) no added v6 interface
 - (b) no added v6 eBGP peering
 - (c) no added IPv6 customer prefixes
 - Permits Inter-Departmental VPNs without additional PE VRFs

LISP Multihoming and Multi-AF

LISP and MPLS Integration

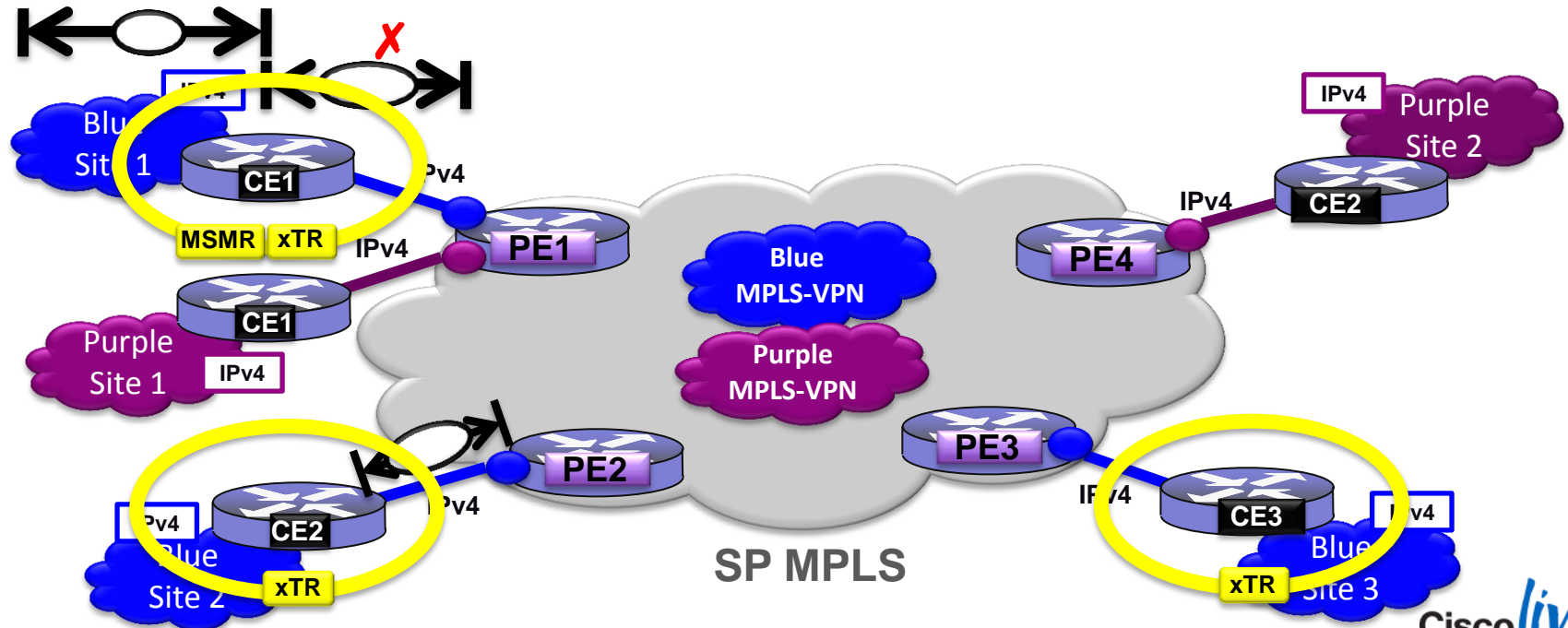
1: Existing IPv4 MPLS



LISP Multihoming and Multi-AF

LISP and MPLS Integration

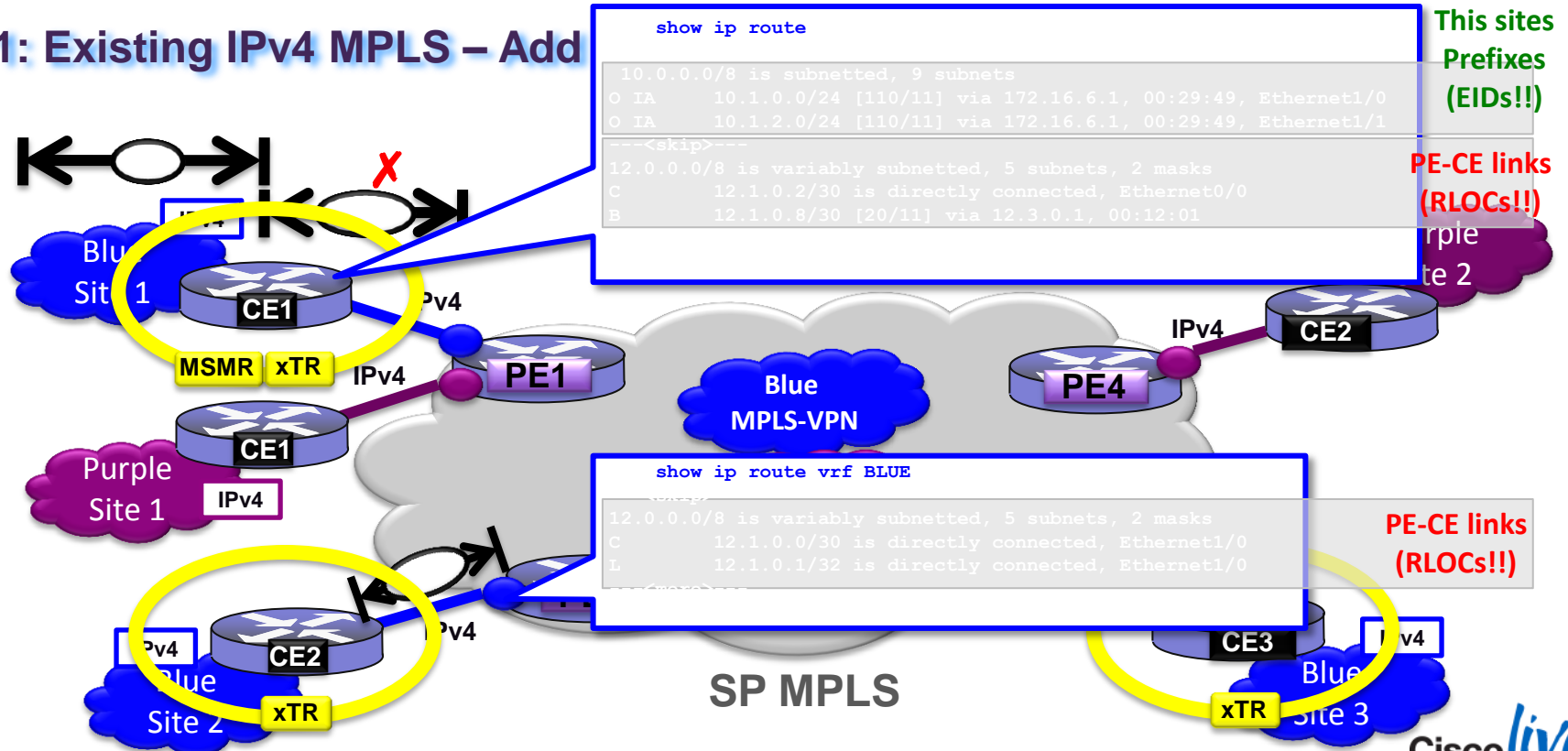
1: Existing IPv4 MPLS – Add LISP!



LISP Multihoming and Multi-AF

LISP and MPLS Integration

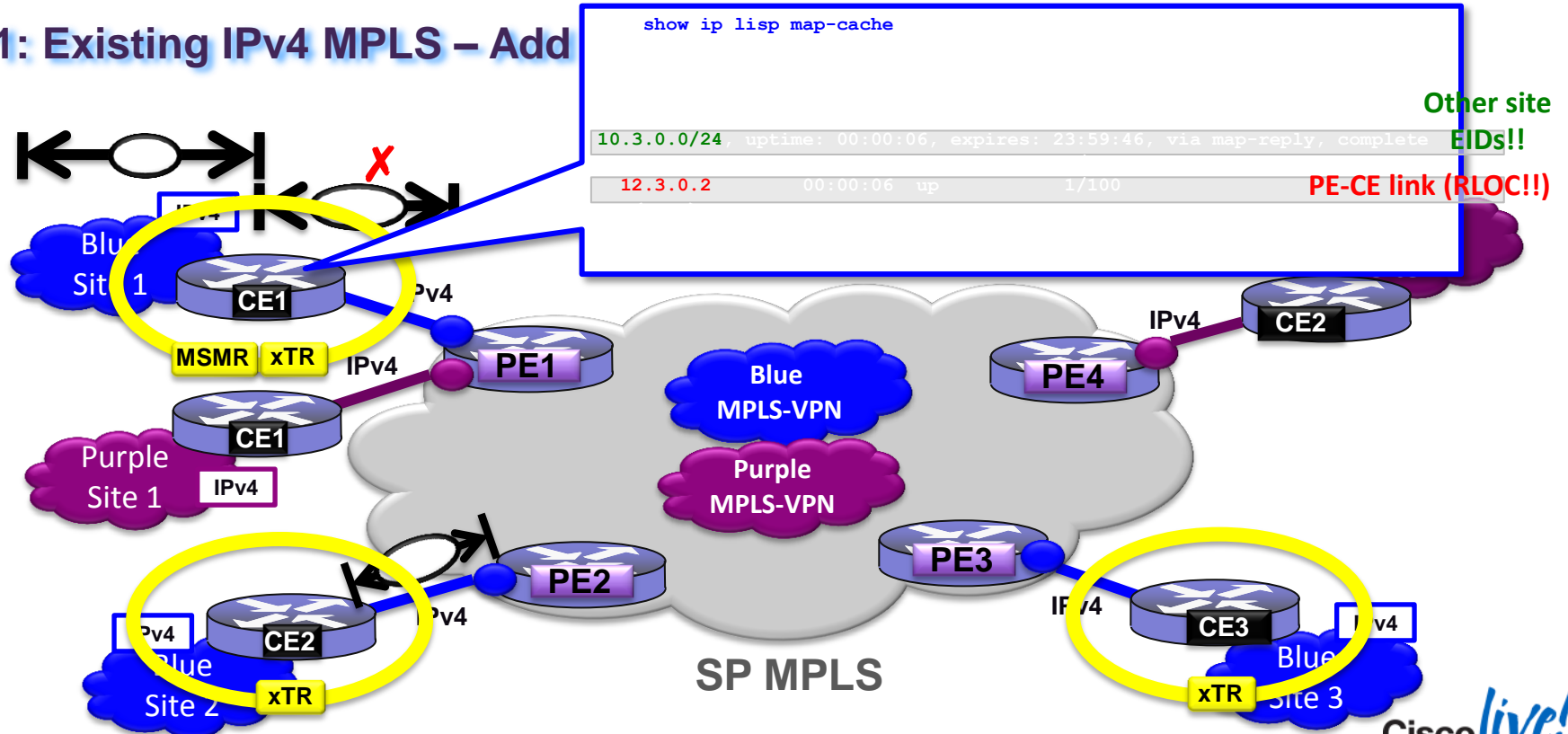
1: Existing IPv4 MPLS – Add



LISP Multihoming and Multi-AF

LISP and MPLS Integration

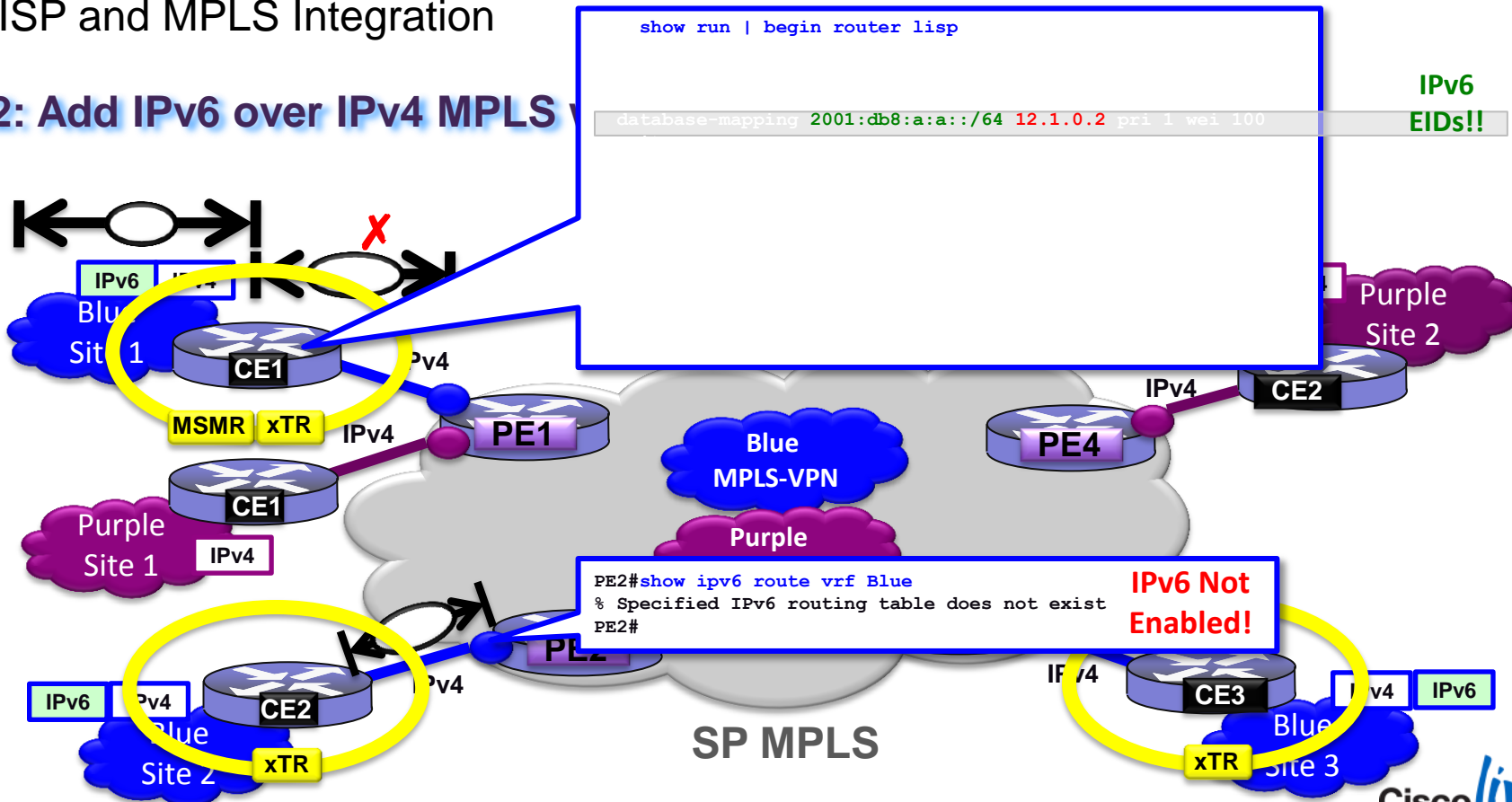
1: Existing IPv4 MPLS – Add



LISP Multihoming and Multi-AF

LISP and MPLS Integration

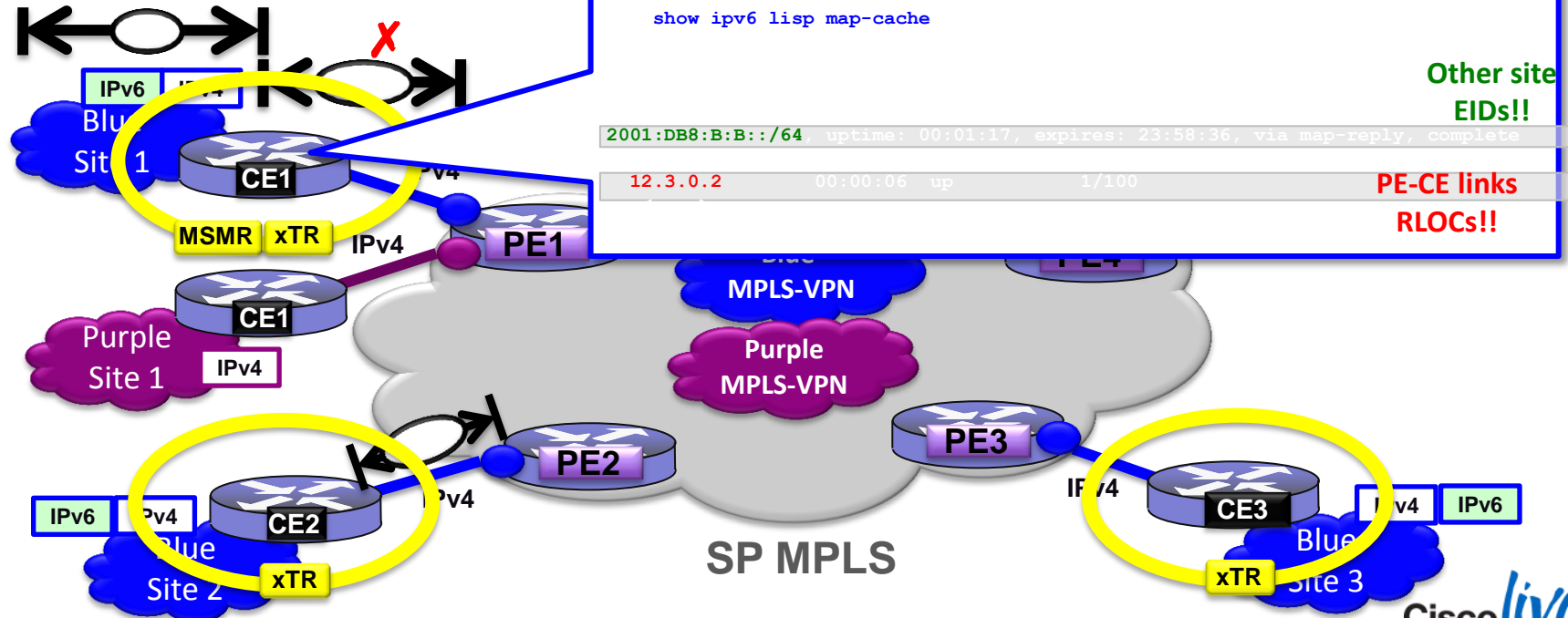
2: Add IPv6 over IPv4 MPLS



LISP Multihoming and Multi-AF

LISP and MPLS Integration

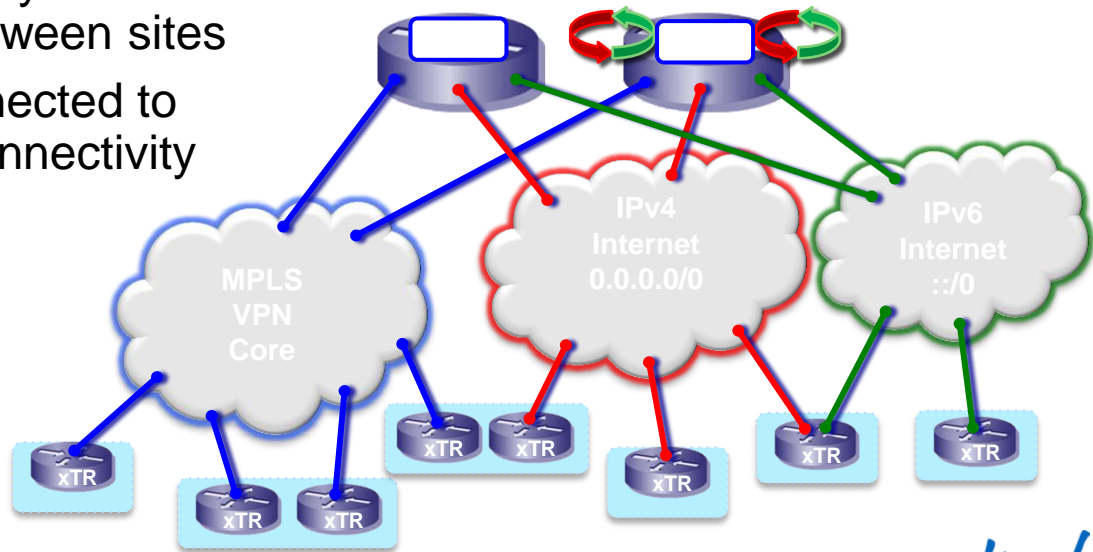
2: Add IPv6 over IPv4 MPLS



LISP Use Cases :: Additional “Useful” Features

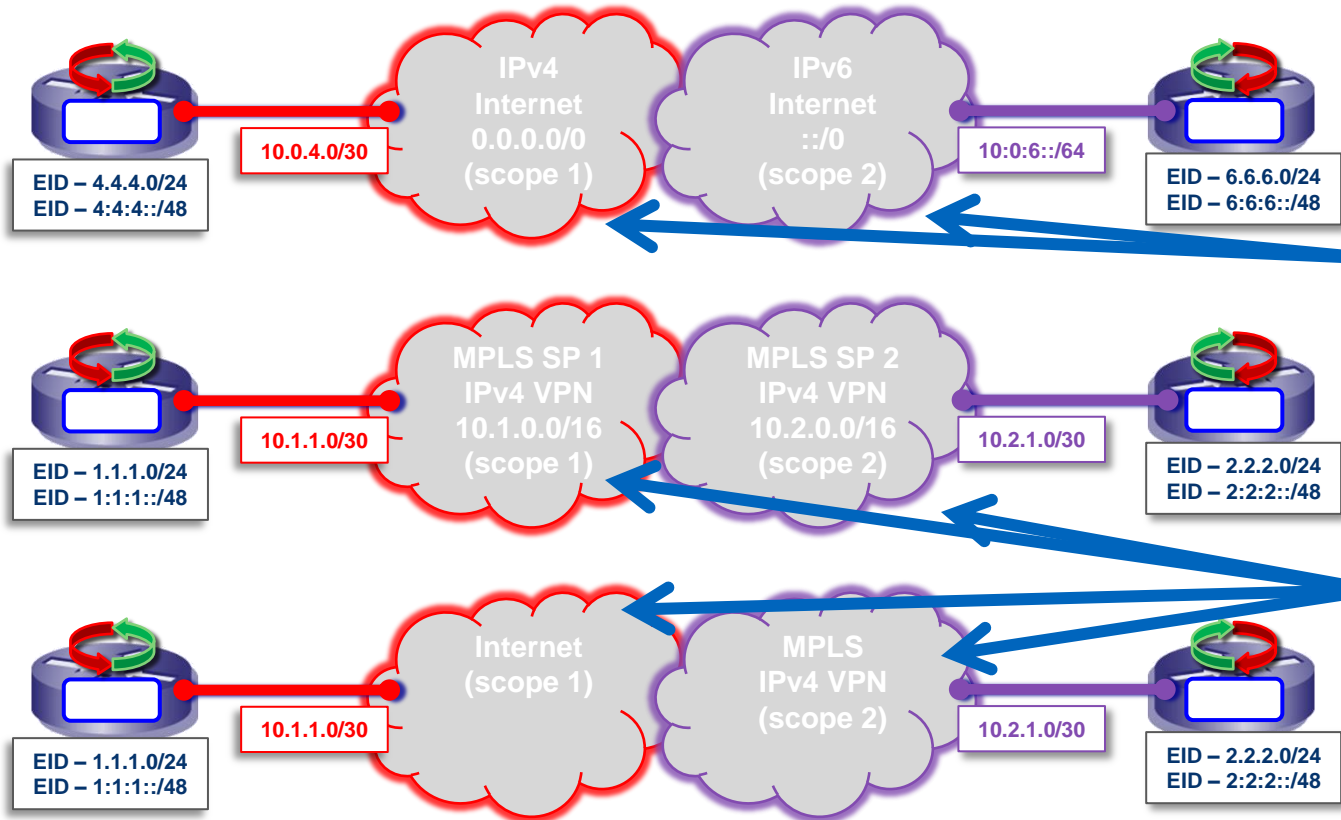
Disjointed Locator Space Support

- Locator/ID separation creates two namespaces: EIDs and RLOCs
 - EID space is the overlay of Enterprise prefixes
 - RLOC space is the underlay network connectivity
- The fundamental principal of any network is that connectivity must exist between sites
- LISP supports sites being connected to locator spaces that have no connectivity to each other!
 - In LISP, this is known as a “disjointed RLOC set”



LISP Use Cases :: Additional “Useful” Features

Disjointed Locator Space Support



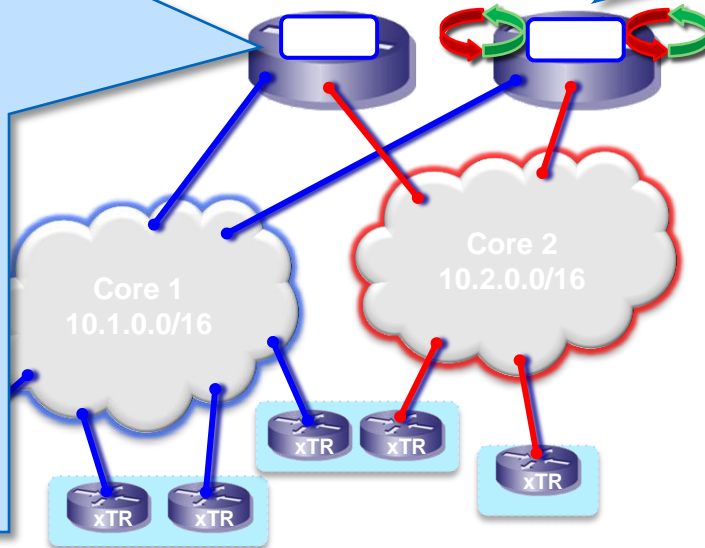
One obvious example of disjointed RLOC spaces is for IPv4 and IPv6 attached sites

The same situation occurs for distinct core networks of the same address family. Two MPLS VPN cores, for example, exhibit disjointed RLOC properties.

LISP Use Cases :: Additional “Useful” Features

Disjoint Locator Space Support

```
!  
router lisp  
locator-set rtr-set1  
 10.1.3.1 priority 1 weight 1  
 exit  
!  
locator-set rtr-set2  
 10.2.3.1 priority 1 weight 1  
 exit  
!  
locator-scope s1  
 rtr-locator-set rtr-set1  
 rloc-prefix 10.1.0.0/16  
 exit  
!  
locator-scope s2  
 rtr-locator-set rtr-set2  
 rloc-prefix 10.2.0.0/16  
 exit  
!  
---<etc.>---
```



```
!  
router lisp  
locator-set setALL  
 10.1.3.1 priority 1 weight 1  
 10.2.3.1 priority 1 weight 1  
 exit  
!  
map-request itr-rlocs setALL  
 eid-table default instance-id 0  
 map-cache 0.0.0.0/0 map-request  
 map-cache ::/0 map-request  
 exit  
!  
---<etc.>---
```

LISP Deployment Examples

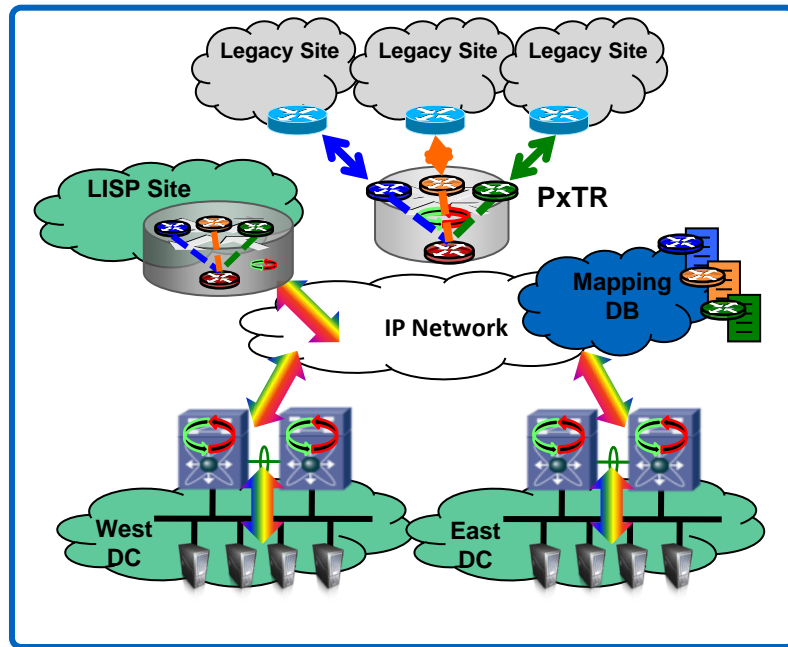
LISP Deployment Examples...

1. Efficient Multihoming and Multi-AF (IPv4 and IPv6)
2. Efficient Virtualisation and High-Scale VPNs
3. Data Centre/Host Mobility
4. LISP-Mobile Node

LISP VPN/Virtualisation

Efficient Virtualisation and High-Scale VPNs – Overview

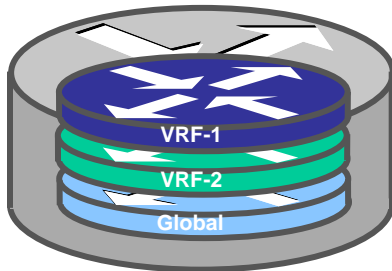
- **Needs:**
 - Integrated Segmentation
 - Global scale and interoperability
 - Minimal Infrastructure disruption
- **LISP Solution:**
 - 24-bit LISP Instance-ID segments control plane and data plane, with VRF binding to the Instance-ID
- **Benefits:**
 - Very high scale tenant segmentation
 - Global mobility + high scale segmentation integrated in single IP solution
 - IP-based “overlay” solution, transport independent
 - No Inter-AS complexity



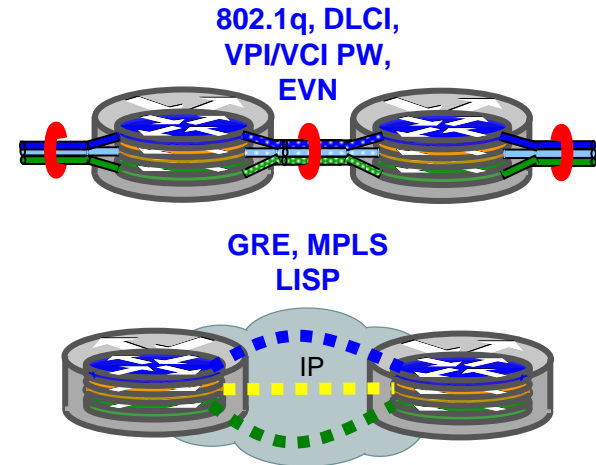
LISP Virtualisation/VPNs

Efficient Virtualisation/Multi-Tenancy Support – Concepts

- Virtualisation of the **DEVICE** level
 - Virtual Routing and Forwarding (VRF) tables segment Layer 3 routing tables
 - VRFs are used to virtualise the component resources
 - Virtualisation secures movement of traffic between networks and enhances security policy options



- Virtualisation of the **PATH** level
 - VRFs assist in path isolation
 - Single-hop (hop-by-hop)
 - Multi-hop (over-the-top)



LISP Virtualisation/VPNs

LISP Virtualisation/Multi-Tenancy Support – Concepts

- Recalling that... LISP is “**Locator/ID**” **separation**... and creates two namespaces: **EIDs** and **RLOCs**... LISP can virtualise both **EID** and **RLOC** namespaces, or both!
- Two models of operation are defined: **Shared** and **Parallel**
 - **Shared Model Virtualisation:**
 - Virtualises the EID namespaces
 - Binds an EID namespace privately defined using a VRF to an Instance-ID
 - Uses a common (shared) RLOC (locator) address space
 - The Mapping System is also part of the locator namespaces and is shared
 - **Parallel Model Virtualisation:**
 - Virtualises the RLOC (locator) namespaces
 - One or more EID instances may share a virtualised RLOC namespace
 - A Mapping System must also be part of each locator namespaces

LISP Virtualisation/VPNs

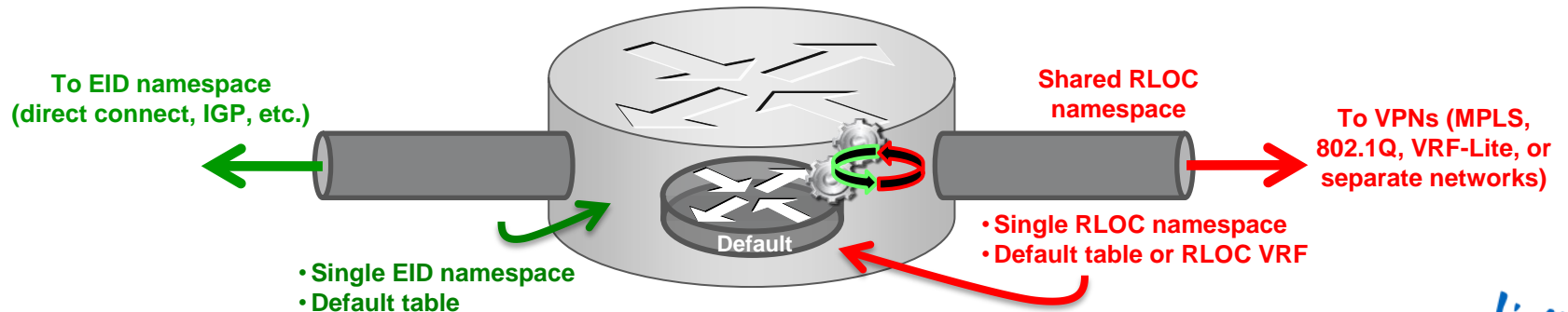
LISP Virtualisation/Multi-Tenancy Support – Concepts

- **RLOC virtualisation** is enabled in conjunction with **locator table** VRFs
- **EID virtualisation** uses **LISP Instance-IDs** in conjunction with **EID** VRFs
 - **Instance-IDs** maintain address space segmentation in control plane and data plane
 - **Instance-IDs** are numerical tags defined in LISP Canonical Address Format (LCAF)
 - IID: a 24-bit unstructured number
 - Data Plane: IID is included in LISP encapsulation header
 - Control Plane: IID is encoded with the EID in LCAF header

LISP Virtualisation/VPNs

LISP Virtualisation/Multi-Tenancy Support – Concepts

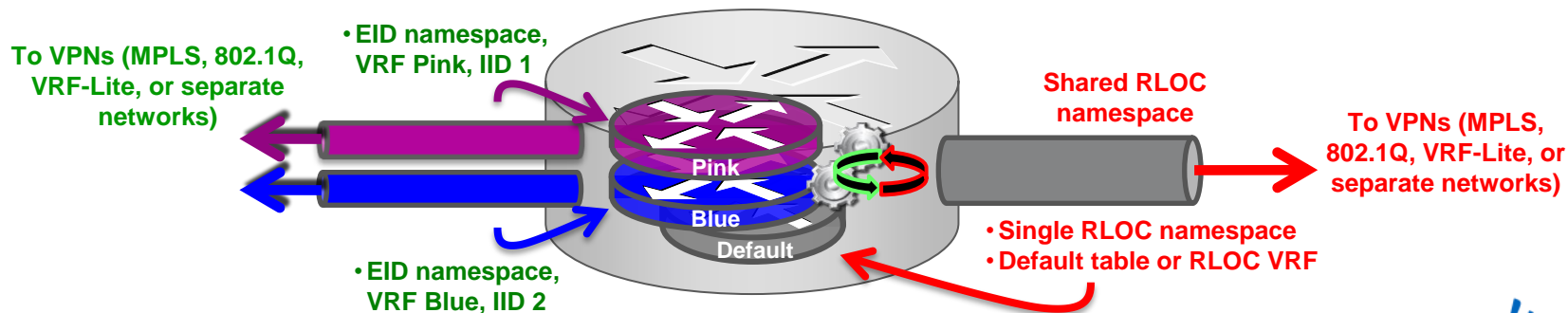
- Default (non-Virtualised) Model – at the device level
 - Conceptually, the Default Model is just a single Parallel Model instance
 - All EID lookups are also in the same single table – default
 - Thus, EIDs are associated with Instance-ID 0
 - All RLOC lookups are in a single table – default
 - The Mapping System is part of the locator address space



LISP Virtualisation/VPNs

LISP Virtualisation/Multi-Tenancy Support – Concepts

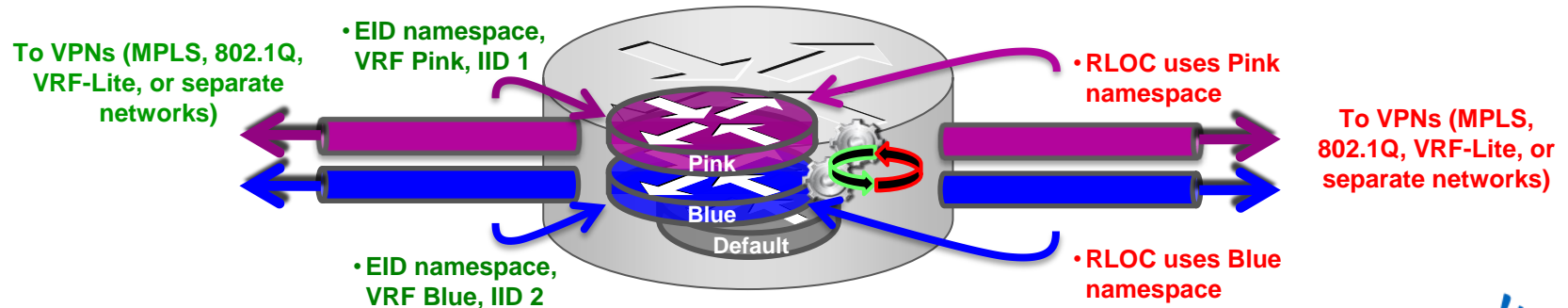
- Shared Model – at the device level
 - Multiple EID-prefixes are allocated privately using VRFs
 - EID lookups are in the VRF associated with an Instance-ID
 - All RLOC lookups are in a single table – (default/global or RLOC VRF)
 - The Mapping System is part of the locator address space and is shared



LISP Virtualisation/VPNs

LISP Virtualisation/Multi-Tenancy Support – Concepts

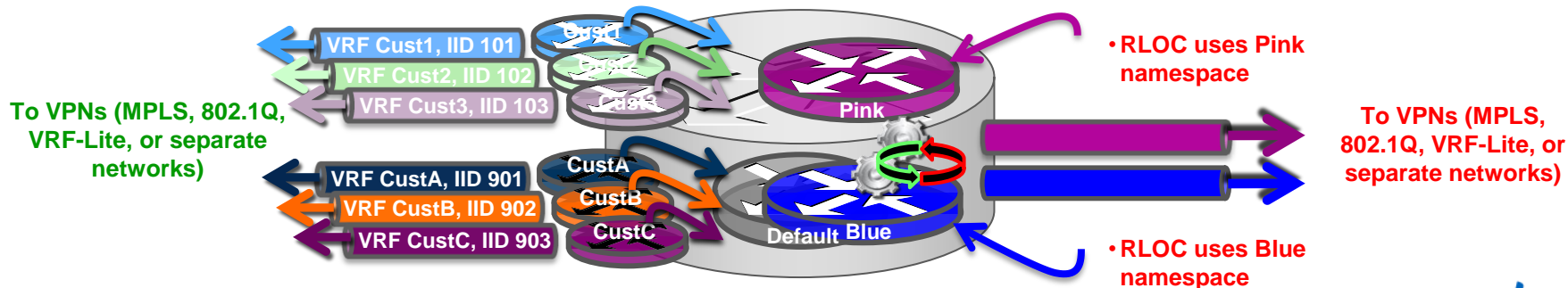
- Parallel Model – at the device level
 - Multiple EID-prefixes are allocated privately using VRFs
 - EID lookups are in the VRF associated with an Instance-ID
 - RLOC lookups are in the VRF associated with the locator table
 - A Mapping System must be part of each locator address space



LISP Virtualisation/VPNs

LISP Virtualisation/Multi-Tenancy Support – Concepts

- Shared and Parallel Models Combined – at the device level
 - Multiple “Shared Model” instantiations combined with Multiple “Parallel Model” instantiations
 - Multiple EID VRFs bound to a single RLOC VRF
 - Multiple RLOC VRFs on the same device



LISP VPN/Virtualisation

Efficient Virtualisation and High-Scale VPNs – Overview

All VPNs share a set of common requirements

1. Encapsulation:

- Virtualisation
 - EID prefix virtualisation
 - Tied to EID VRFs
 - Locators can be virtualised too

2. Site to Site Routing:

- Spoke to spoke connectivity
- Optional local Internet offload (split-tunnel)
- No IGP required to branch sites!

LISP VPNs Routing? or Tunnelling? -- It's BOTH!

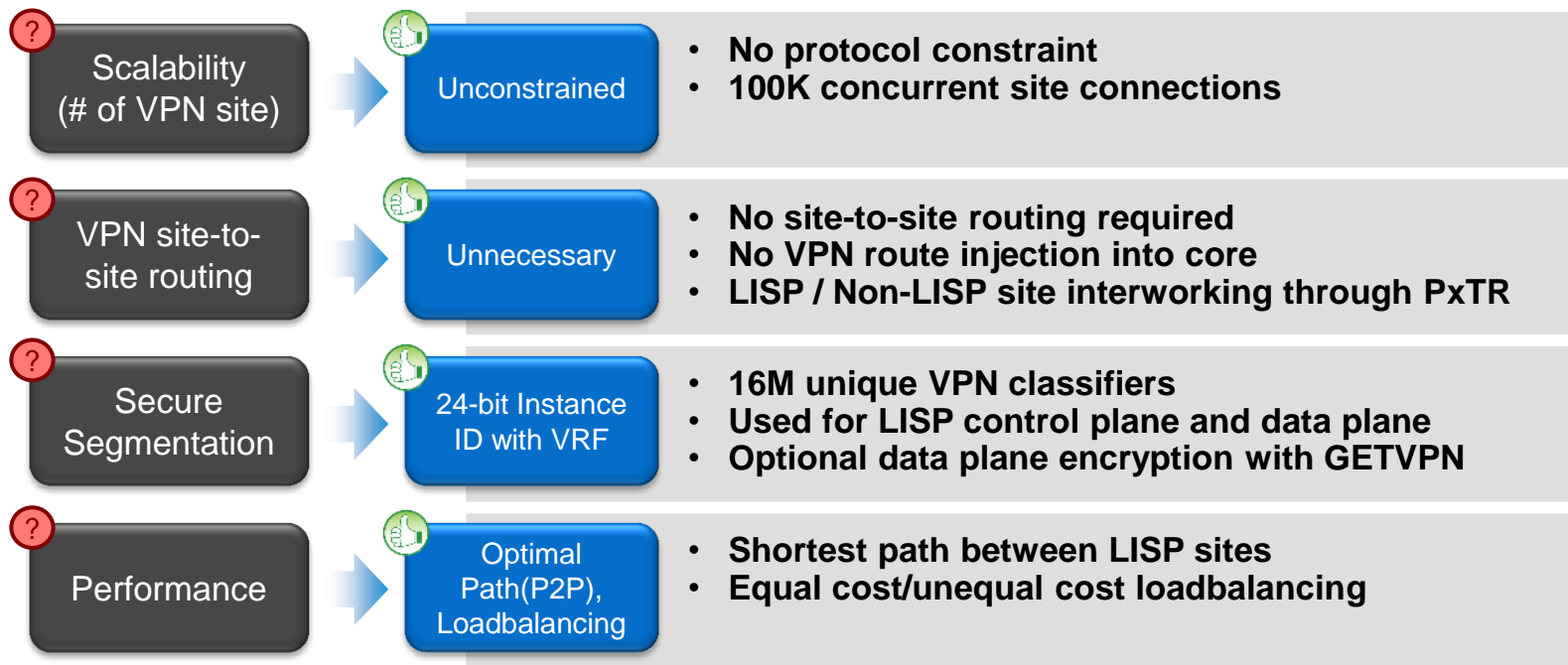
3. Security: Built-in and Add-on

- Built-in security mechanisms
- LISP Works with any crypto scheme
 - **Locators** or **EIDs** can be encrypted
- LISP-SEC for control plane security

LISP VPN/Virtualisation

Efficient Virtualisation and High-Scale VPNs – Overview

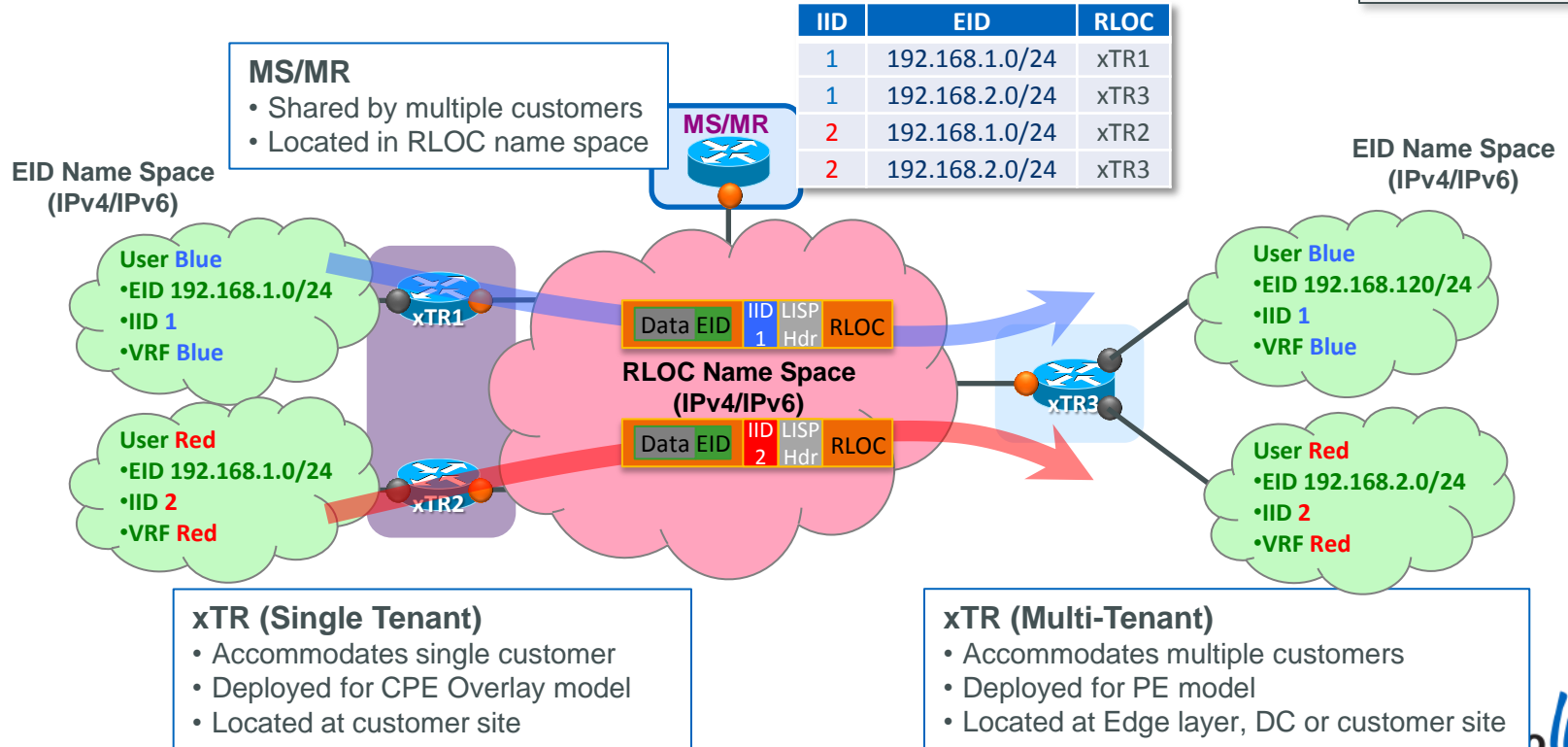
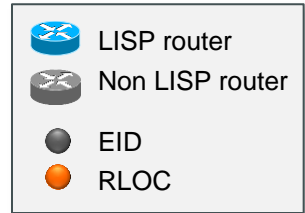
LISP – Inherently scalability and virtualisation, rapidly deployable



LISP VPN/Virtualisation

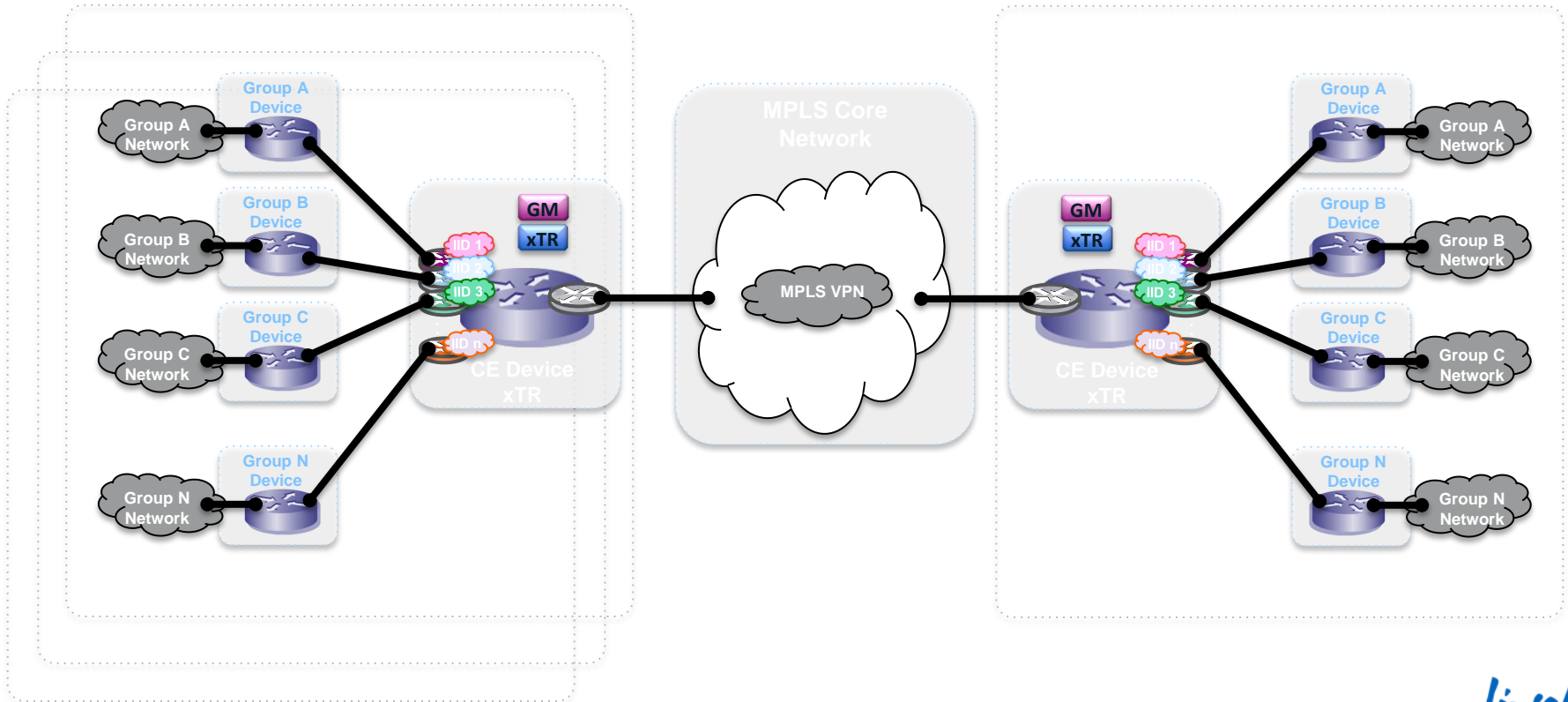
Efficient Virtualisation and High-Scale VPNs – Overview

Generalised LISP Shared Model deployment



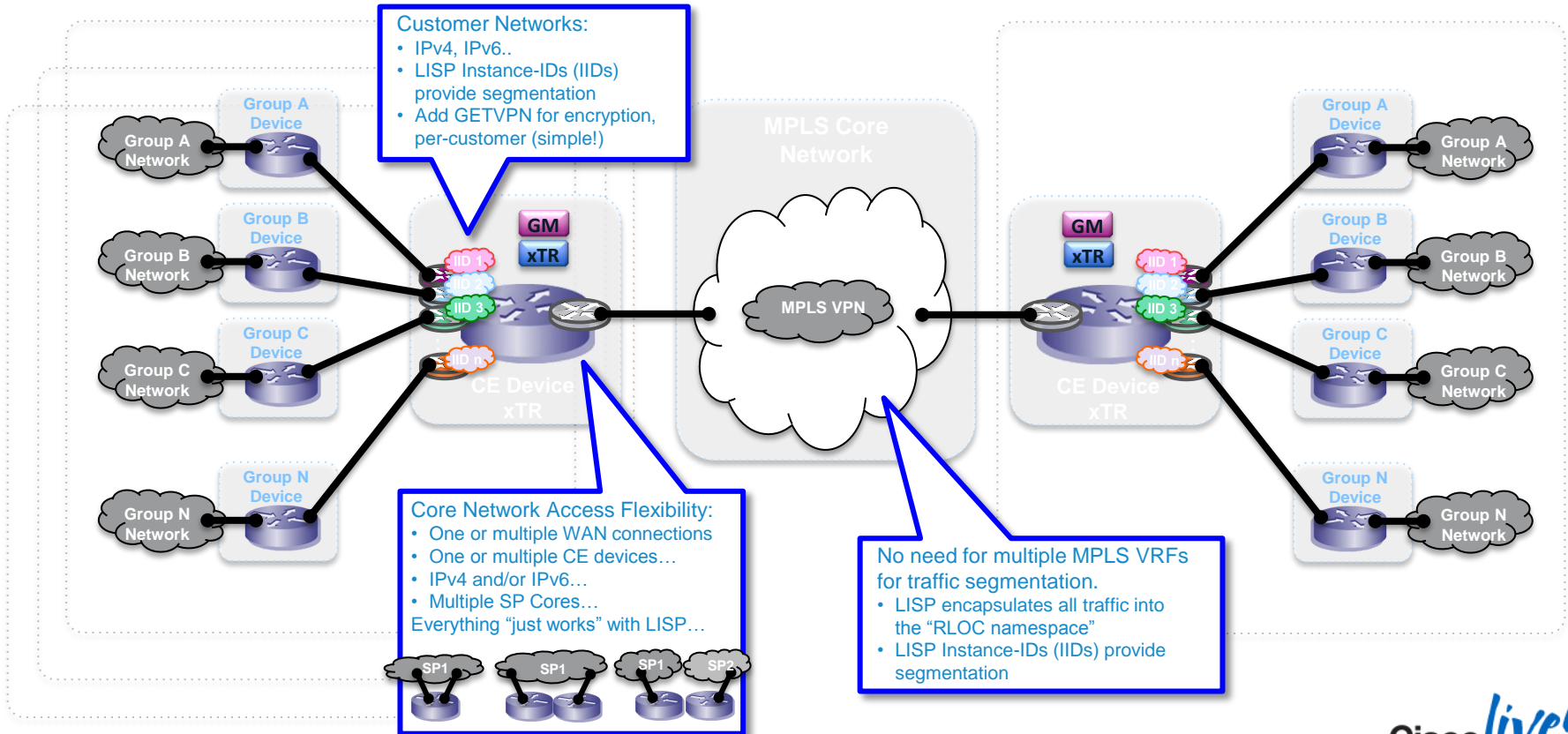
LISP Use Cases :: Virtualisation/VPNs

Customer Example :: US State Government (Multi-tenancy)



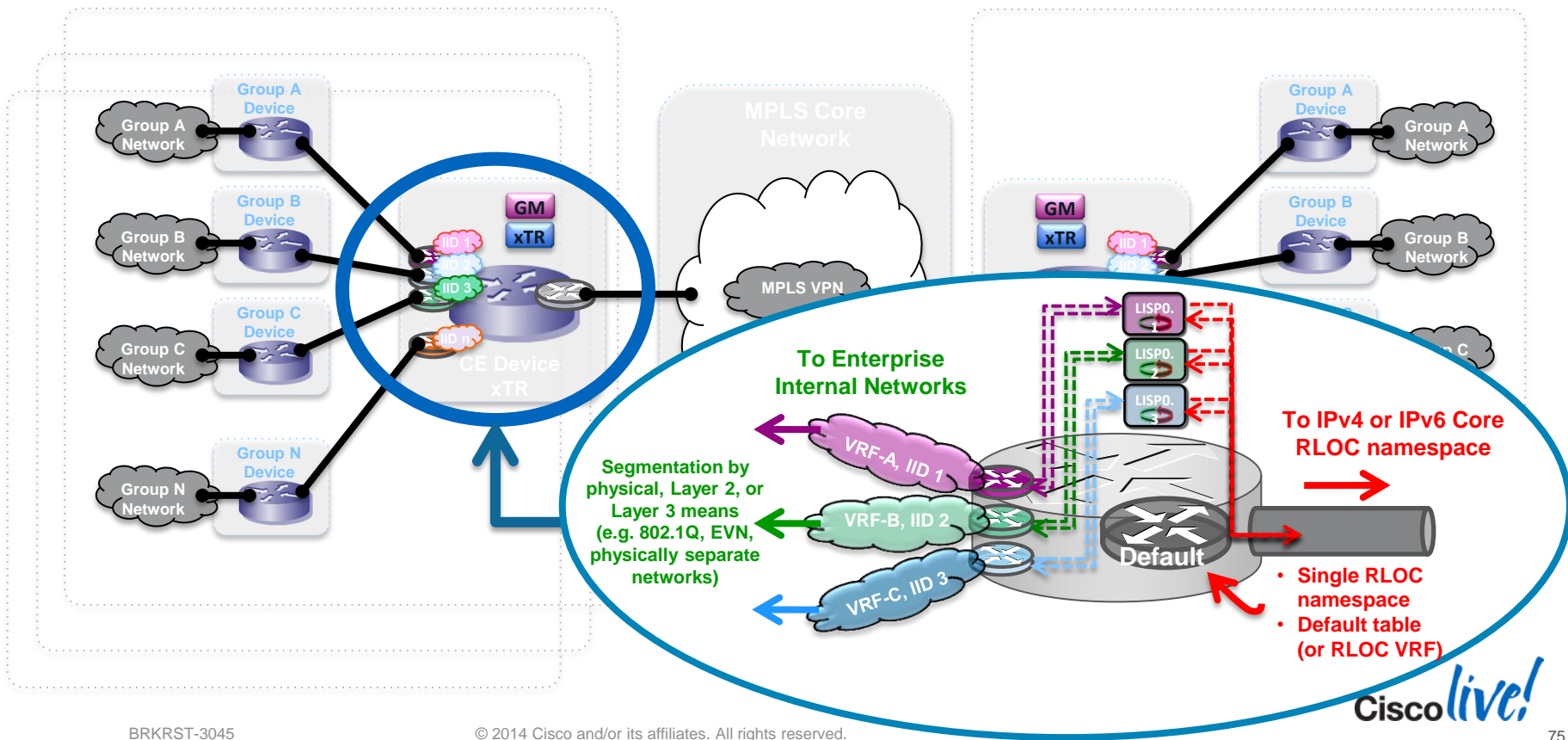
LISP Use Cases :: Virtualisation/VPNs

Customer Example :: US State Government (Multi-tenancy)



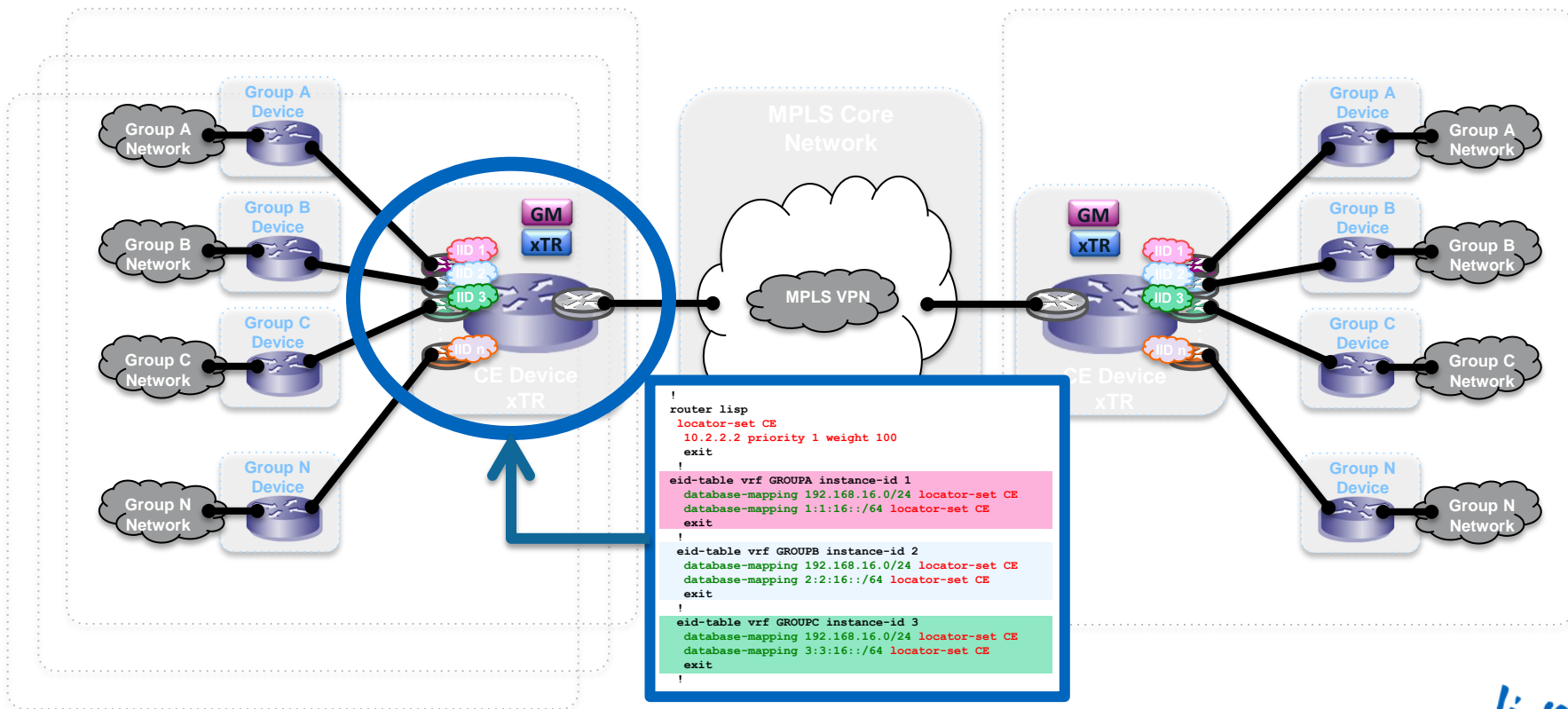
LISP Use Cases :: Virtualisation/VPNs

Customer Example :: US State Government (Multi-tenancy)



LISP Use Cases :: Virtualisation/VPNs

Customer Example :: US State Government (Multi-tenancy)



LISP Virtualisation/VPNs

LISP Virtualisation/Multi-Tenancy Support – Concepts

- LISP and encryption (IOS)
 - Recalling that... LISP is “Locator/ID” separation... and creates two namespaces: **EIDs** and **RLOCs**
 - LISP provides two ways to apply a crypto map

Use-Case		Vanilla IPsec	GETVPN	Comments
LISP Default Model	crypto-map on RLOC	✓	✓	LISP encap first, then encryption based on RLOC
	crypto-map on LISP0	✓	✓	Encryption first based on EID, then LISP encap
LISP Virtualisation	crypto-map on RLOC	✓	✓	LISP encap first, then encryption based on RLOC
	crypto-map on LISP0.x	✓	✓	Encryption first based on EID, then LISP encap

lisp.cisco.com

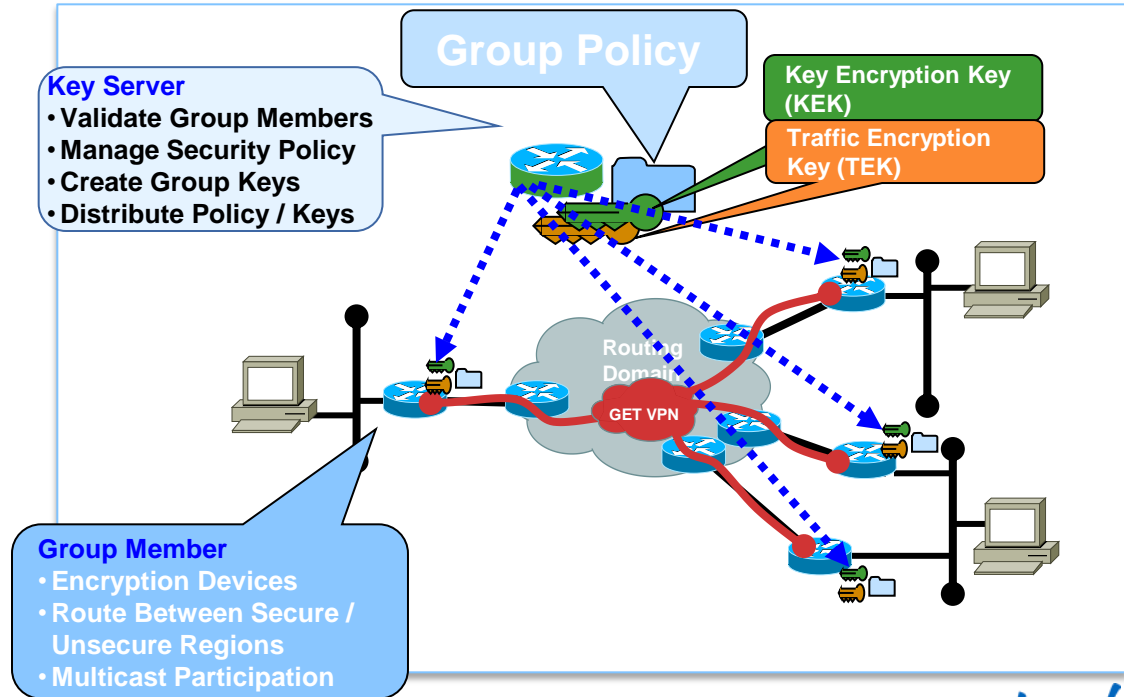
LISP Virtualisation/VPNs

LISP Virtualisation/Multi-Tenancy Support – Concepts

- Group Domain of Interpretation (GDOI) RFC 6407 – adding encryption

- **GDOI**

- RFC 6407
- “Stateless” IPsec
- Traffic encryption keys computed on Key Server, distributed to all Group Members
- Better scaling than vanilla IPsec

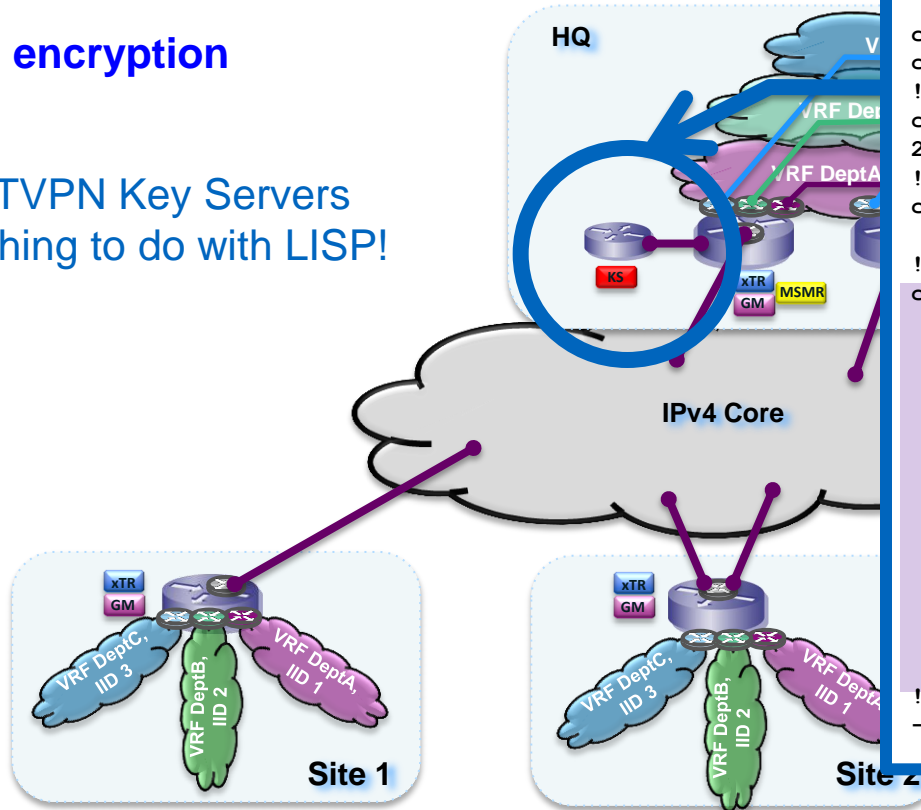


LISP VPN/Virtualisation

Efficient Virtualisation and High-Scale VPNs – Overview

3. Add encryption

- GETVPN Key Servers
- Nothing to do with LISP!



```

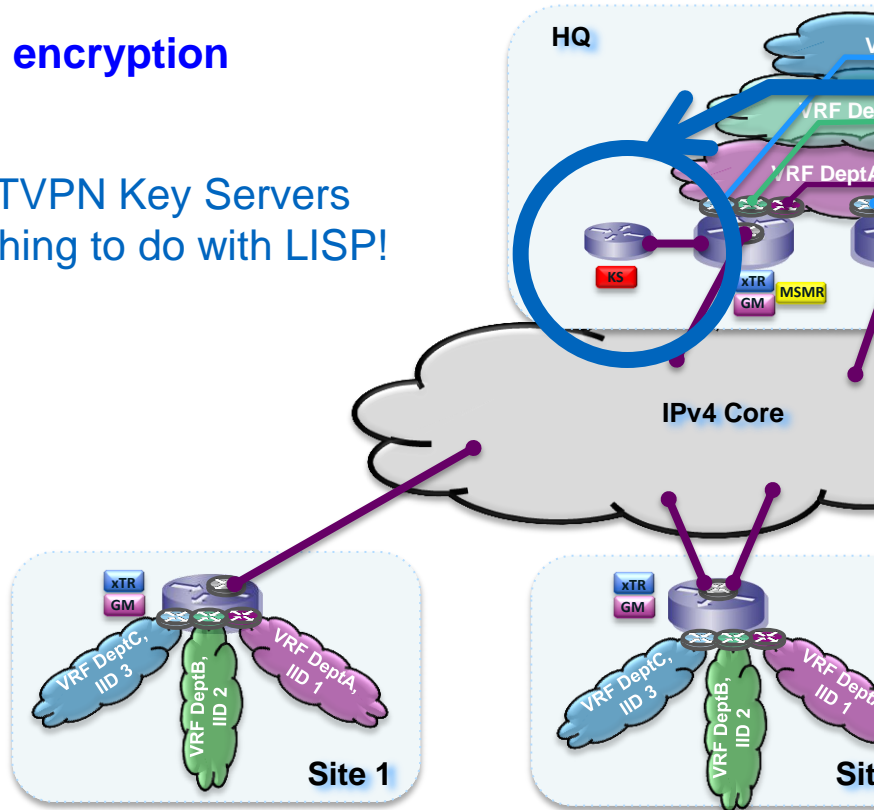
!
crypto isakmp policy 10
  encr aes 256
  authentication pre-share
  group 16
crypto isakmp key FOO address 0.0.0.0
crypto isakmp keepalive 15 periodic
!
crypto ipsec transform-set GDOI-TRANS esp-aes
256 esp-sha512-hmac
!
crypto ipsec profile GDOI-PROFILE
  set transform-set GDOI-TRANS
!
crypto gdoi group V4GROUP-0001
  identity number 10001
  server local
  rekey retransmit 60 number 2
  rekey authentication mypubkey rsa GET-KEYS1
  rekey transport unicast
  sa ipsec 1
  profile GDOI-PROFILE
  match address ipv4 GETVPN-0001
  replay time window-size 5
  address ipv4 192.168.18.2
  redundancy
  local priority 100
  peer address ipv4 192.168.19.2
!
----<cont.>----
  
```

LISP VPN/Virtualisation

Efficient Virtualisation and High-Scale VPNs – Overview

3. Add encryption

- GETVPN Key Servers
- Nothing to do with LISP!



```

! ---<cont.>---
!
crypto gdoi group ipv6 V6GROUP-0003
  identity number 20003
  server local
  rekey retransmit 60 number 2
  rekey authentication mypubkey rsa GET-KEYS3
  rekey transport unicast
  sa ipsec 1
  profile GDOI-PROFILE
  match address ipv6 GETVPN6-0003
  replay time window-size 5
  address ipv4 192.168.18.2
  redundancy
  local priority 100
  peer address ipv4 192.168.19.2
!

ip access-list extended GETVPN-0001
  permit ip any any
ip access-list extended GETVPN-0002
  permit ip any any
ip access-list extended GETVPN-0003
  permit ip any any
!

ipv6 access-list GETVPN6-0001
  permit ipv6 any any
!

ipv6 access-list GETVPN6-0002
  permit ipv6 any any
!

ipv6 access-list GETVPN6-0003
  permit ipv6 any any
!

```

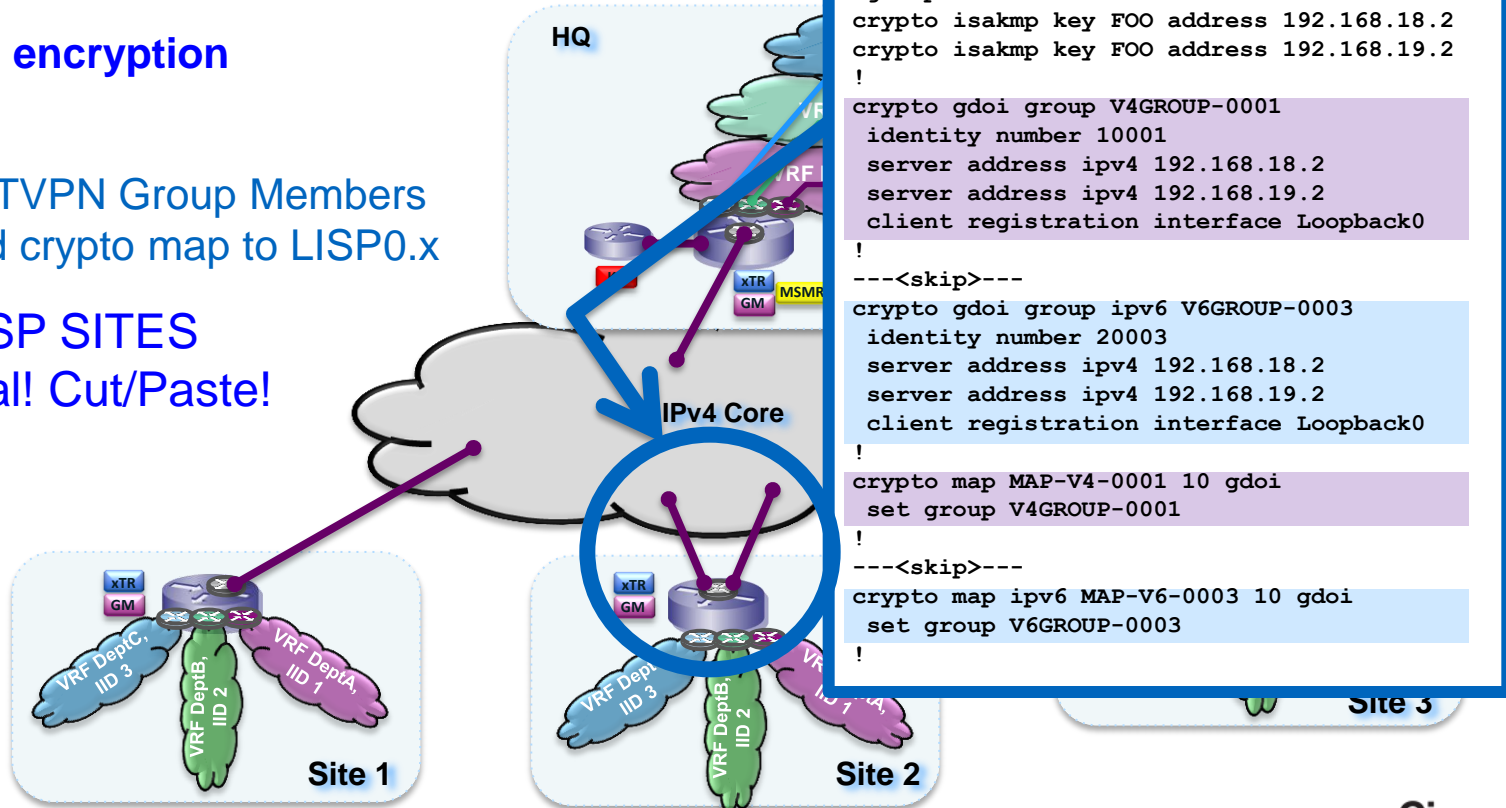

LISP VPN/Virtualisation

Efficient Virtualisation and High-Scale VPNs –

3. Add encryption

- GETVPN Group Members
- Add crypto map to LISP0.x

ALL LISP SITES
identical! Cut/Paste!



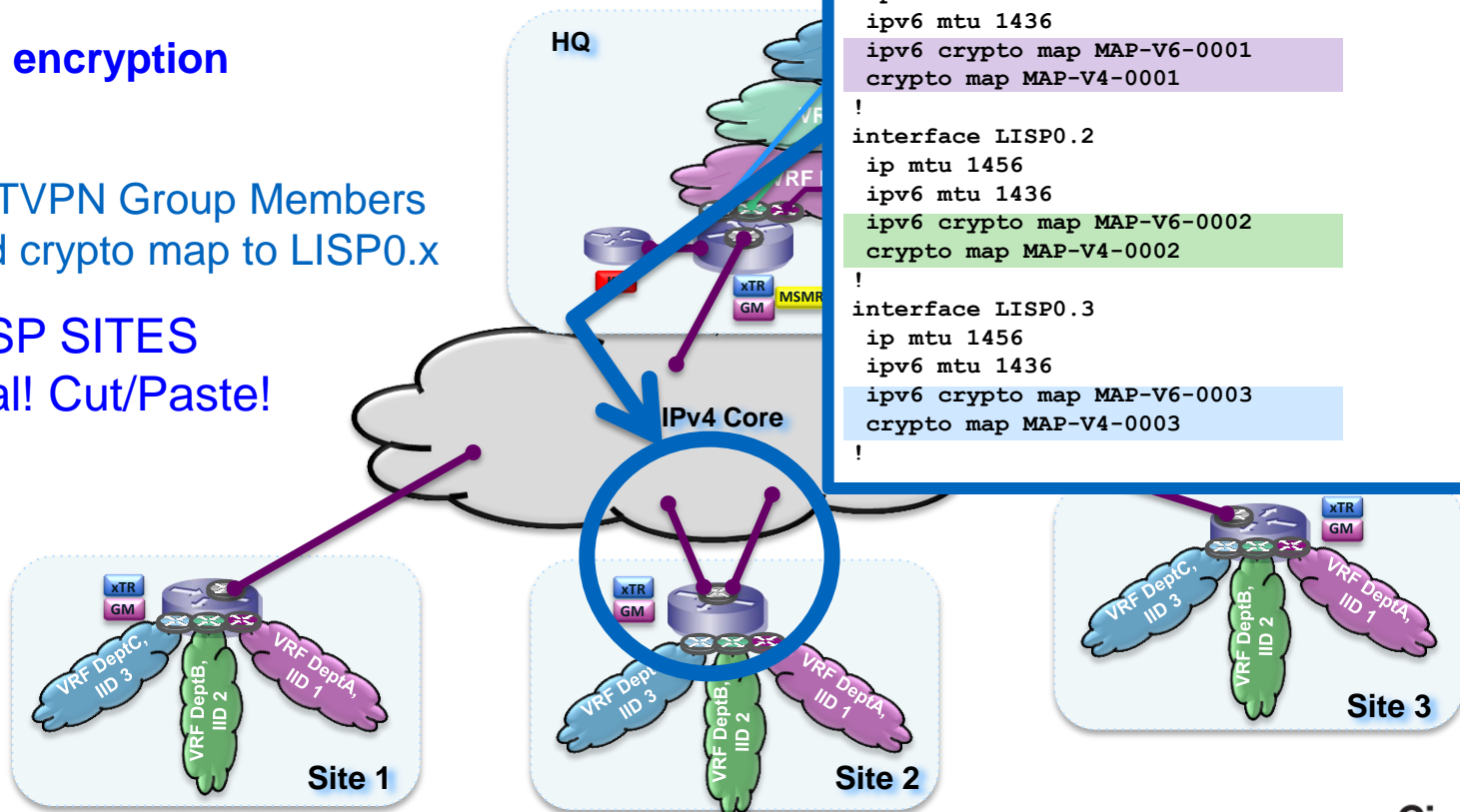
LISP VPN/Virtualisation

Efficient Virtualisation and High-Scale VPNs –

3. Add encryption

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ALL LISP SITES
identical! Cut/Paste!



LISP VPN/Virtualisation

Efficient Virtualisation and High-Scale VPNs – Overview

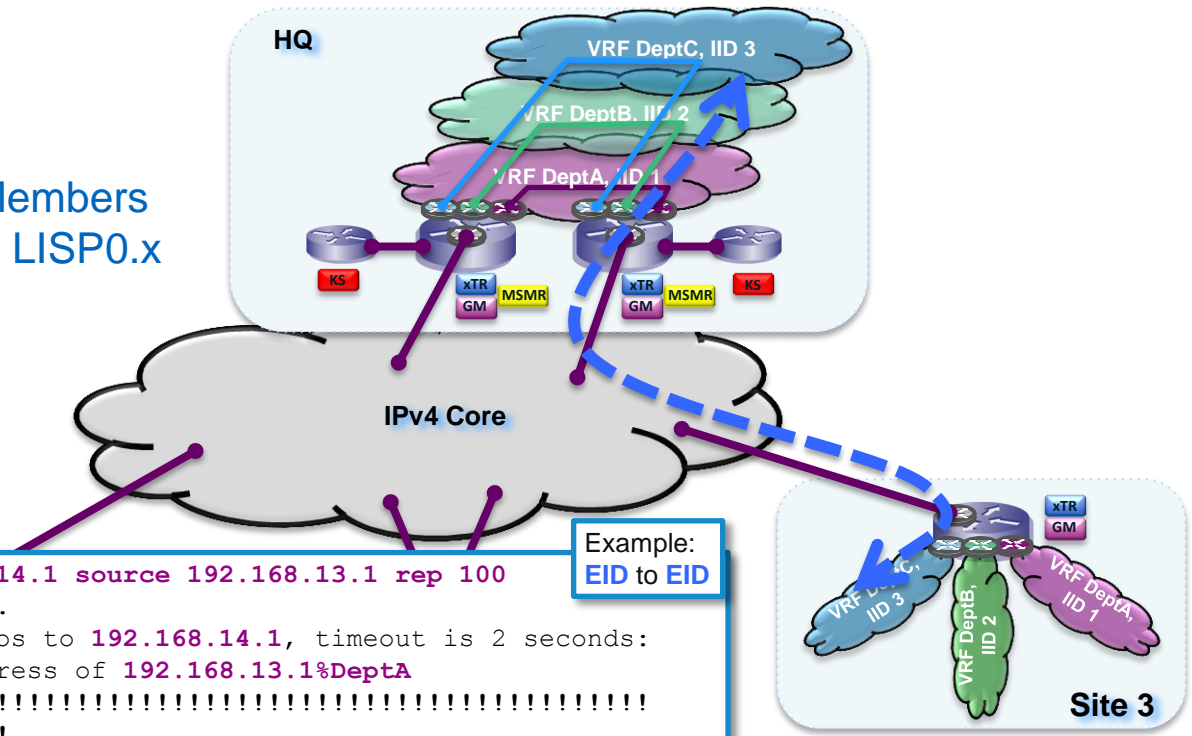
3. Add encryption

- GETVPN Group Members
- Add crypto map to LISP0.x

Verification...

```
Site3#ping vrf DeptA 192.168.14.1 source 192.168.13.1 rep 100
Type escape sequence to abort.
Sending 10, 100-byte ICMP Echos to 192.168.14.1, timeout is 2 seconds:
Packet sent with a source address of 192.168.13.1%DeptA
!!!!!!!!!!!!!!!!!!!!!!!!!!!!!!!!!!!!!!!!!!!!!!!!!!!!!!!!!!!!!!!!!!!!!!!!!!!!
!!!!!!!!!!!!!!!!!!!!!!!!!!!!!!!!!!!!!!!!!!!!!!!!!!!!!!!!!!!!!!!!!!!!!!!!!!!!
Success rate is 100 percent (100/100), round-trip min/avg/max = 5/6/12 ms
Site3#
```

Example:
EID to EID



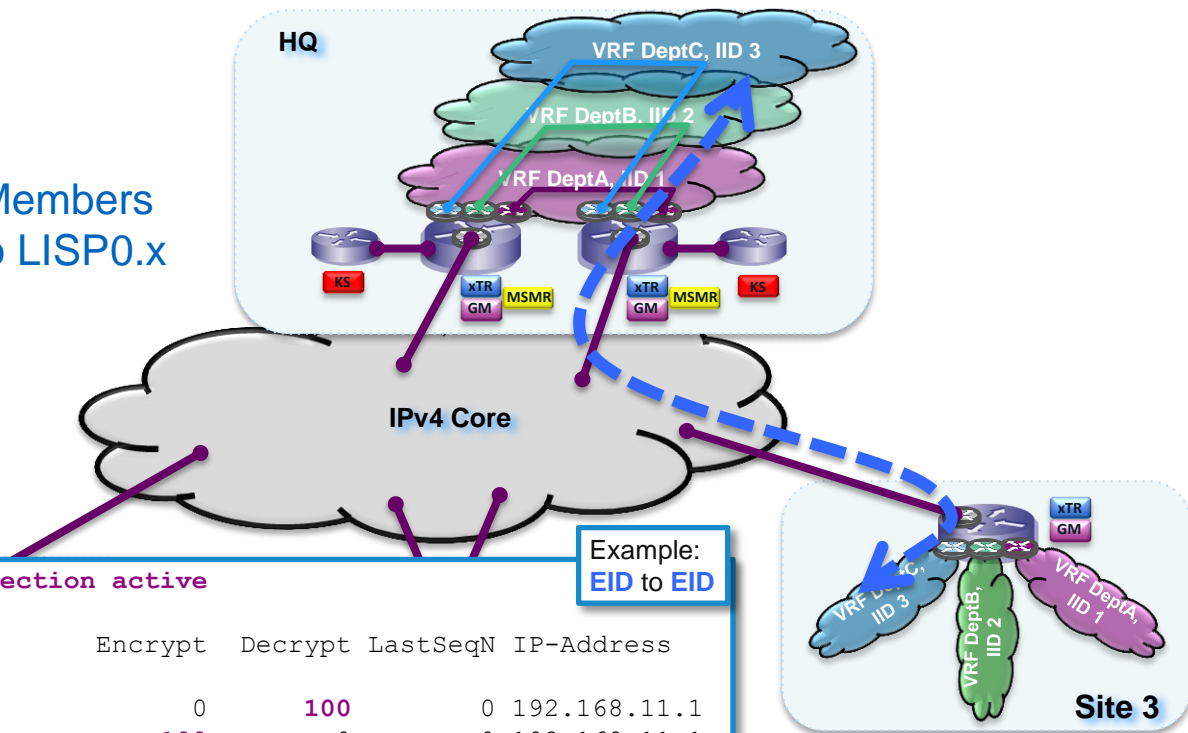
LISP VPN/Virtualisation

Efficient Virtualisation and High-Scale VPNs – Overview

3. Add encryption

- GETVPN Group Members
- Add crypto map to LISP0.x

Verification...



```
Site3#show crypto engine connection active
```

```
Crypto Engine Connections
```

ID	Type	Algorithm	Encrypt	Decrypt	LastSeqN	IP-Address
---	<skip>	---				
143	IPsec	AES256+SHA512	0	100	0	192.168.11.1
144	IPsec	AES256+SHA512	100	0	0	192.168.11.1

```
---
```

```
Site3#
```

LISP Deployment Examples

LISP Deployment Examples...

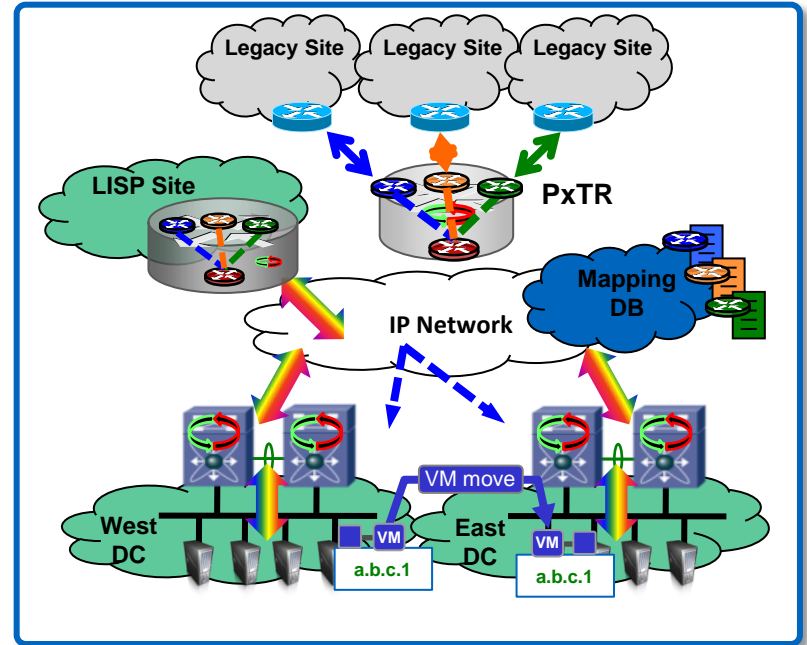
1. Efficient Multihoming and Multi-AF (IPv4 and IPv6)
2. Efficient Virtualisation and High-Scale VPNs
3. Data Centre/Host Mobility
4. LISP-Mobile Node

LISP Data Centre/Host Mobility

Data Centre/Host Mobility – Overview

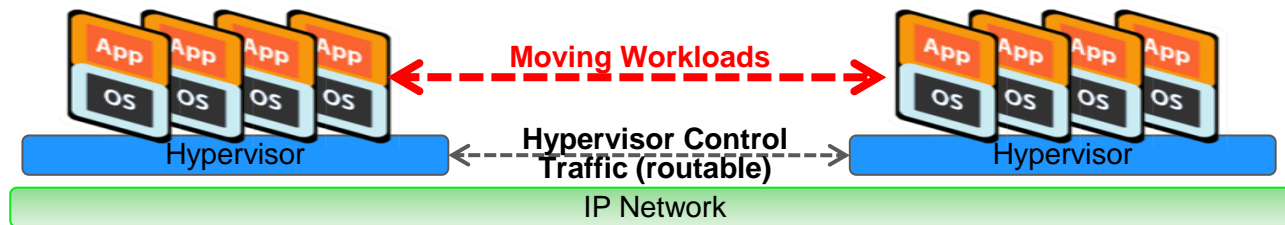
- **Needs:**
 - VM-Mobility extending subnets and across subnets
 - Move detection, dynamic EID-to-RLOC mappings, traffic redirection
- **LISP Solution:**
 - OTV + LISP for VM-moves in extended subnets
 - LISP for VM-moves across subnets
- **Benefits:**
 - VM OS agnostic, seamless, integrated, global workload mobility
 - Direct Path (no triangulation)
 - Connections survive across moves
 - No routing re-convergence, no DNS updates
 - Global Scalability (cloud bursting)
 - ARP elimination

Data Centre/Host Mobility

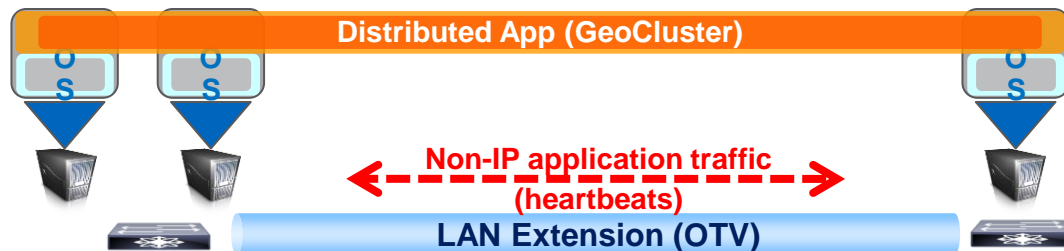


Moving vs. Distributing Workloads

Why do we really need LAN Extensions?



- **Move workloads** with IP mobility solutions: LISP Host Mobility
 - IP preservation is the real requirement (LAN extensions not mandatory)



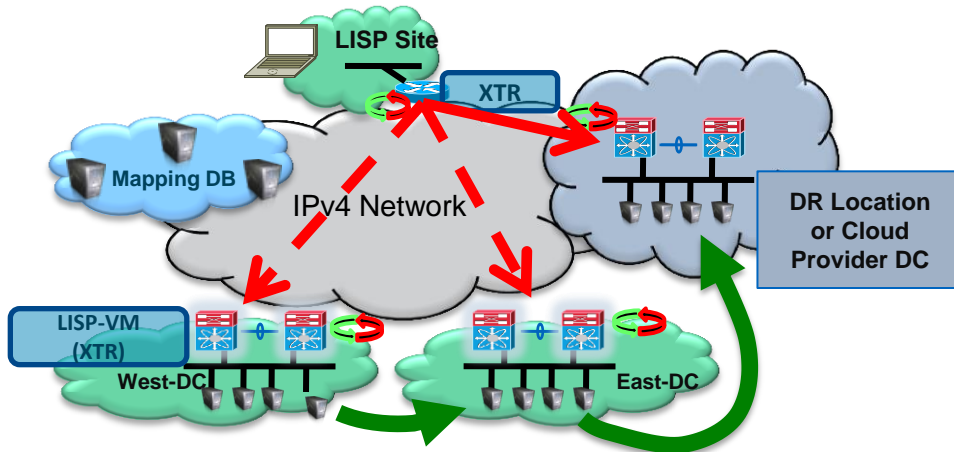
- **Distribute workloads** with LAN extensions
 - Application High Availability with Distributed Clusters

Host-Mobility Scenarios

Two Mobility Scenarios

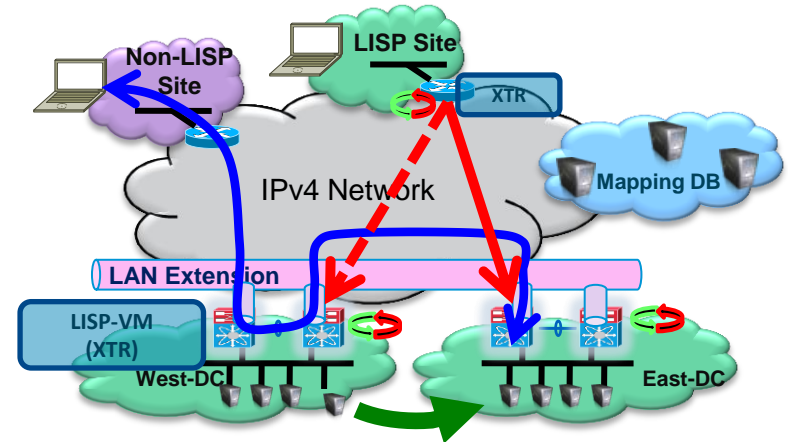
LISP Host Mobility Config Guide:
<http://lisp.cisco.com>

Moves Without LAN Extension



- IP Mobility Across Subnets
 - Disaster Recovery
 - Cloud Bursting
- Application Members In One Location

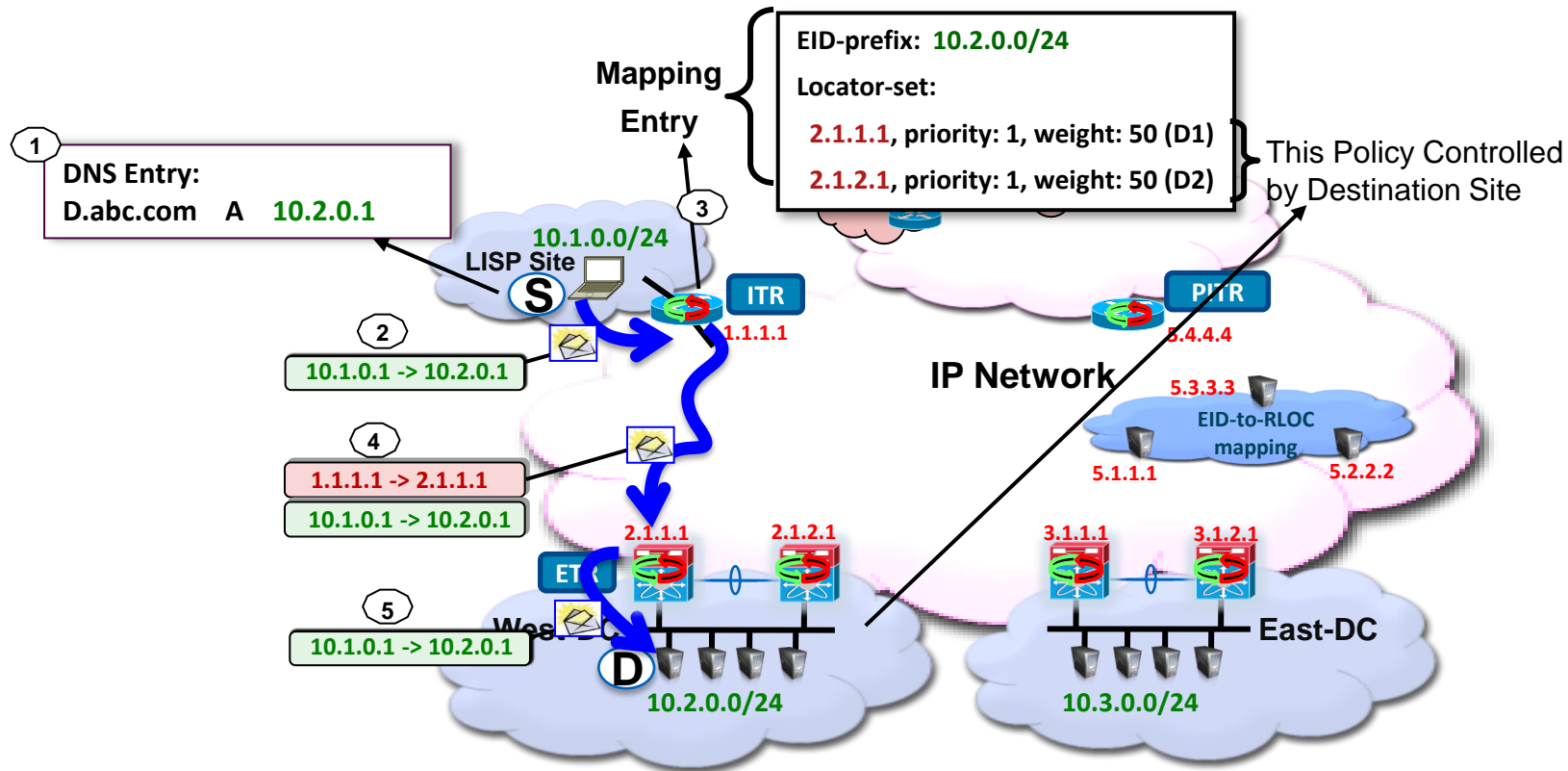
Moves With LAN Extension



- Routing for Extended Subnets
 - Active-Active Data Centres
 - Distributed Data Centres
- Application Members Distributed
 - Broadcasts across sites

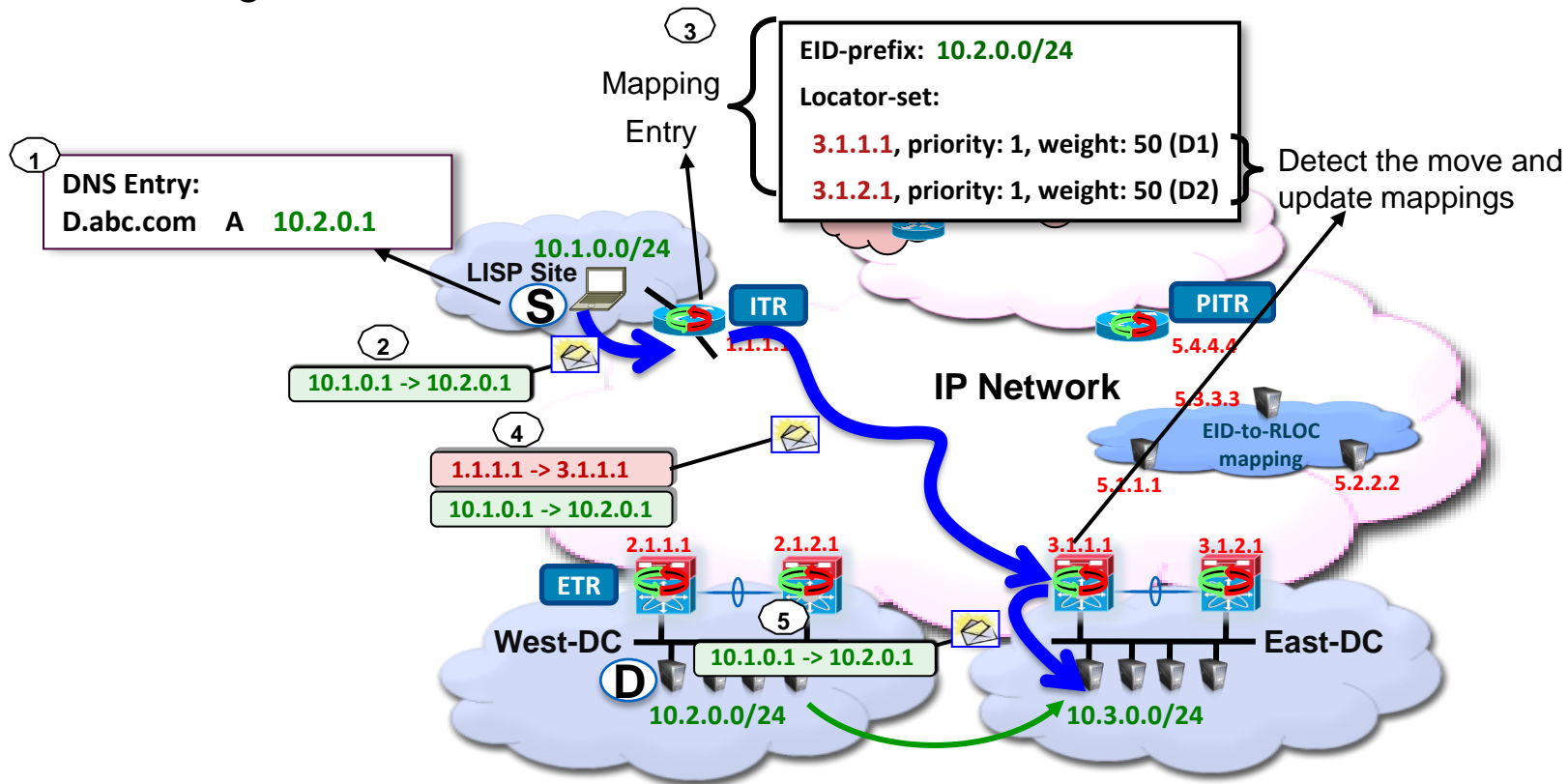
A LISP Packet Walk

Before Moving the Host



A LISP Packet Walk

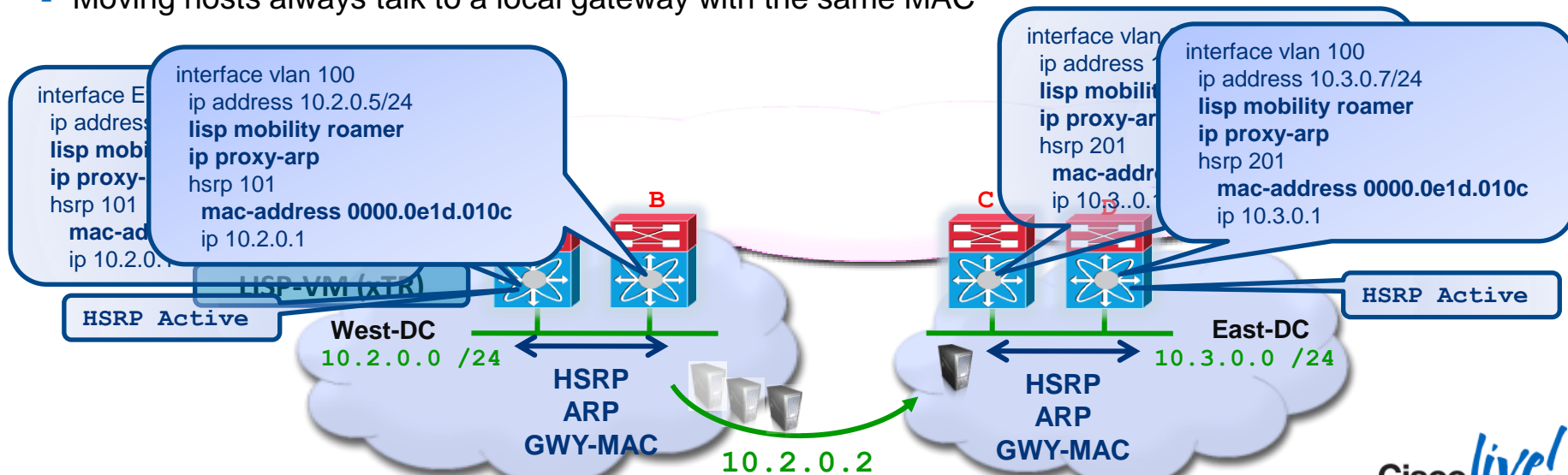
After Moving the Host



LISP Host-Mobility – First Hop Routing

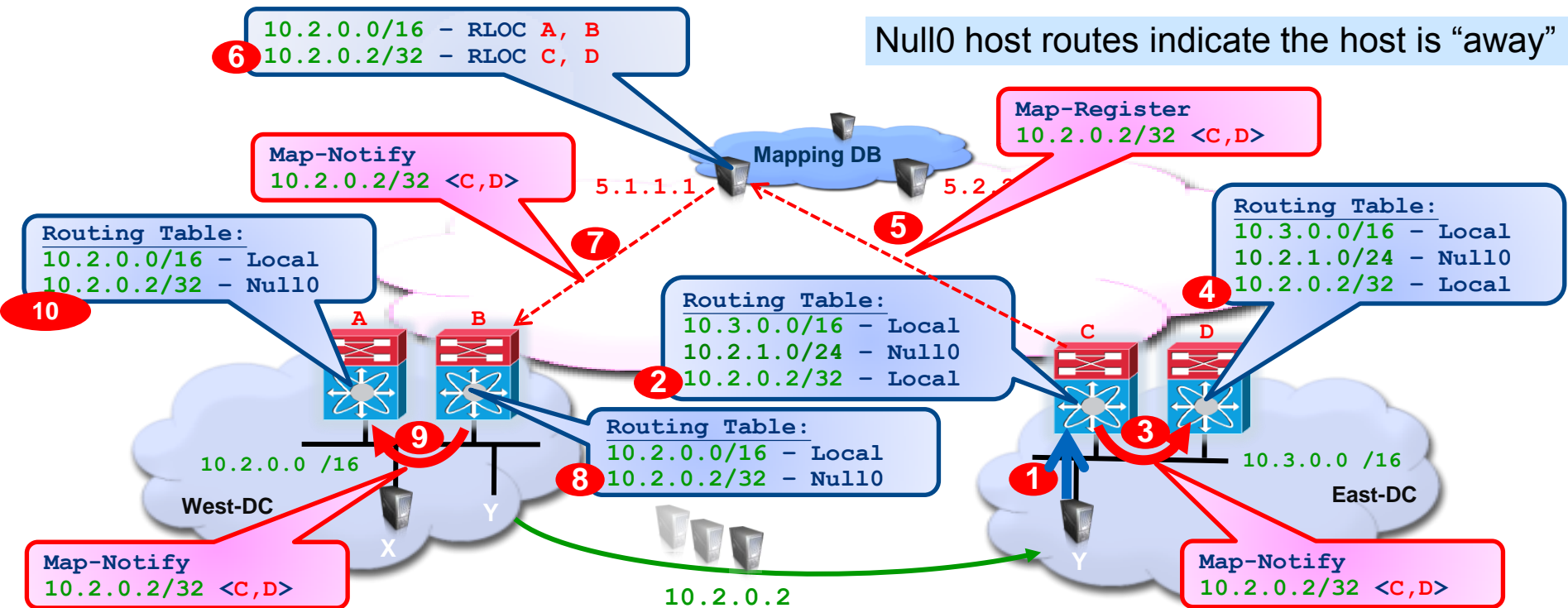
No LAN Extension

- SVI (Interface VLAN x) and HSRP configured as usual
 - Consistent GWY-MAC configured across all dynamic subnets
- The lisp mobility <dyn-eid-map> command enables proxy-arp functionality on the SVI
 - The LISP-VM router services first hop routing requests for both local and roaming subnets
- Moving hosts always talk to a local gateway with the same MAC



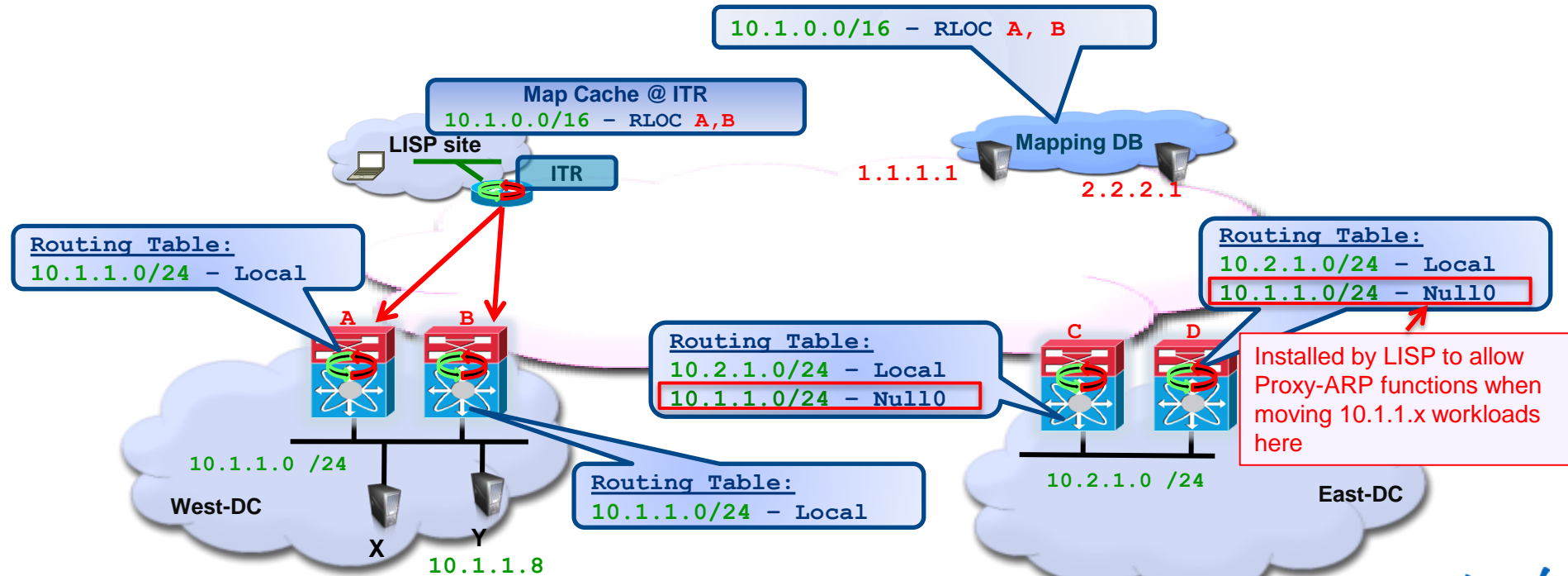
Host-Mobility and Multi-homing

ETR Updates – Across LISP Sites



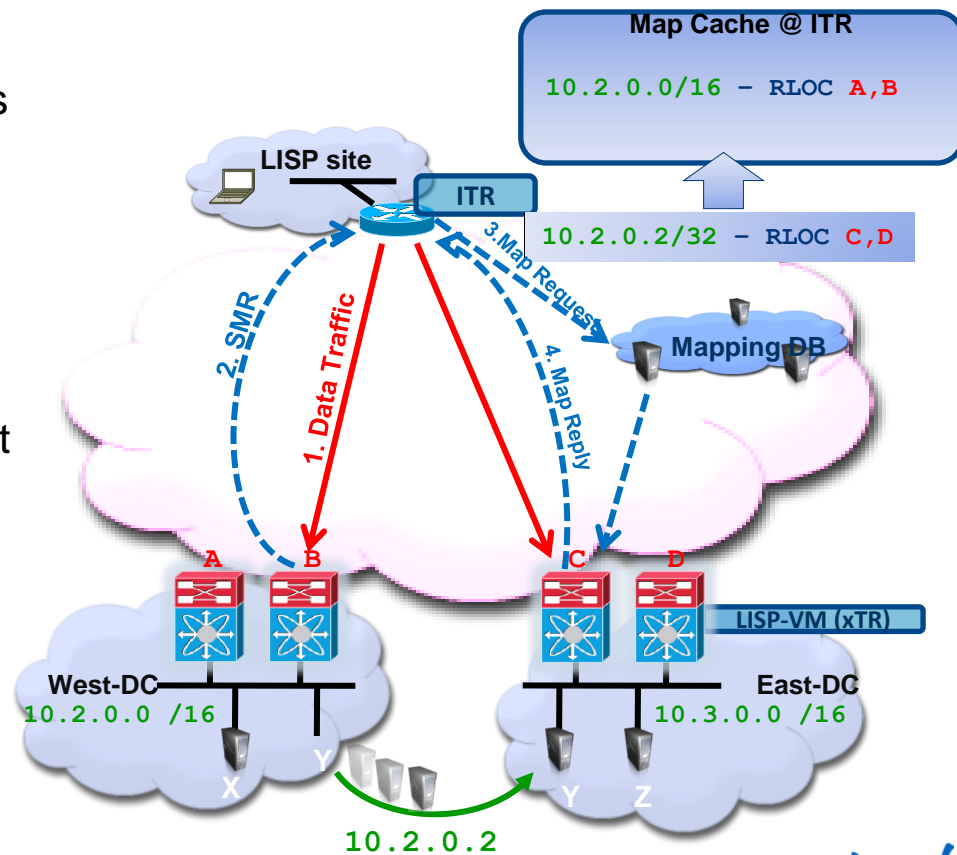
LISP Mobility Across LISP Sites

Client-server communication established without the need to discover the workloads in the “home subnet” in West-DC



Refreshing the Map Caches

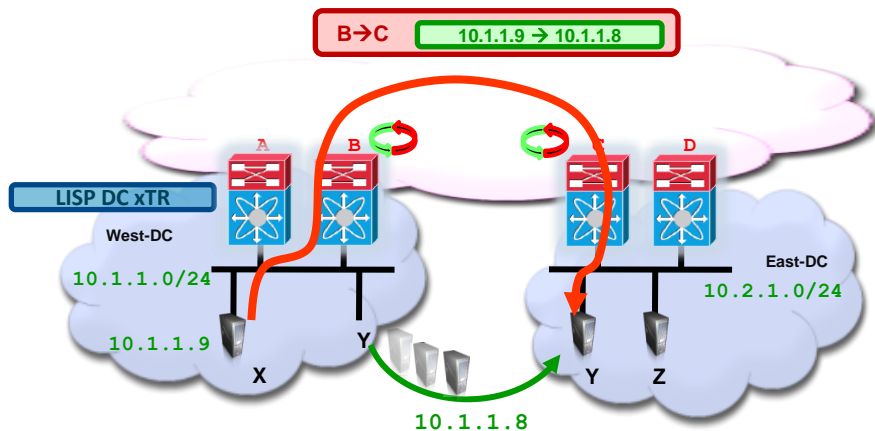
- ITRs and PITRs with cached mappings continue to send traffic to the old locators
 - The old xTR knows the host has moved (Null0 route)
- Old xTR sends Solicit Map Request (SMR) messages to any encapsulators sending traffic to the moved host
- The ITR then initiates a new map request process
- An updated map-reply is issued from the new location
- The ITR Map Cache is updated
 - Traffic is now re-directed
 - SMRs are an important integrity measure to avoid unsolicited map responses and spoofing



On-subnet Server-Server Traffic

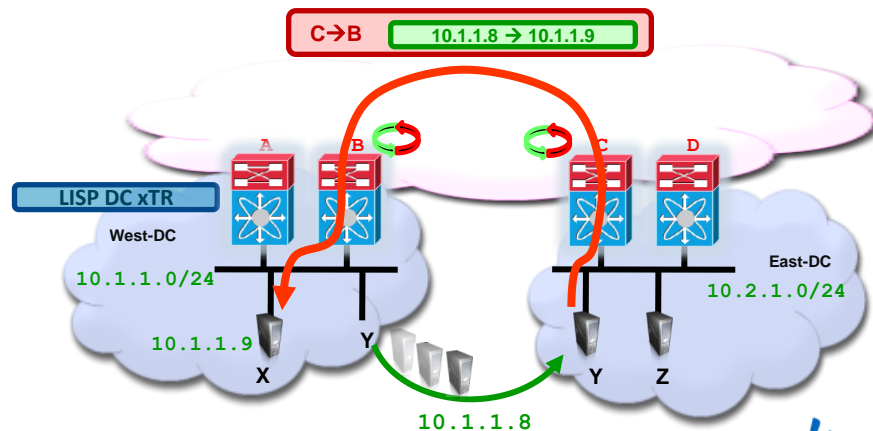
West-to-East

- X ARPs for Y, /32 Null0 entry for Y triggers proxy-ARP on West-DC xTRs to ensure traffic is steered there
 - Note: entry for Y in X ARP cache is cleared by GARP message originated by West-DC XTRs
- Traffic to Y is **LISP encapsulated**



East-to-West

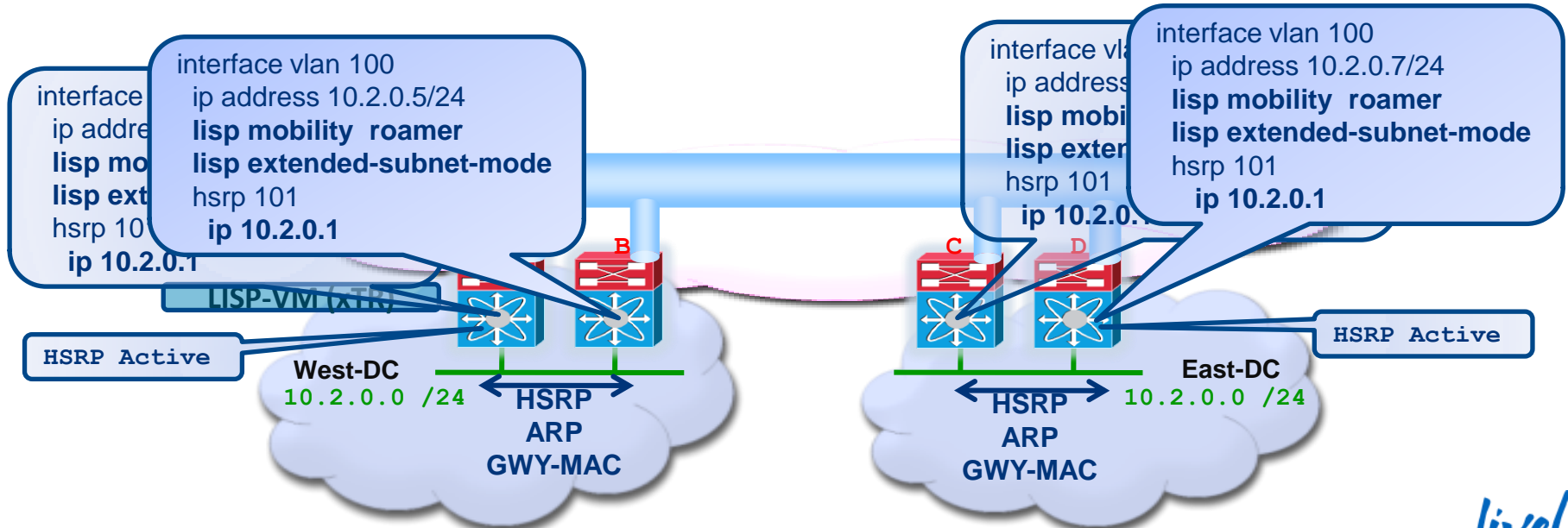
- Y ARPs for X, /24 Null0 entry for the 'home subnet' triggers proxy-ARP on East DC xTRs to ensure traffic is steered there
 - Note: assumption is that ARP cache on Y is refreshed after the move
- Traffic to X is **LISP encapsulated**



LISP Host-Mobility – First Hop Routing

With Extended Subnets

- Consistent GWY-IP and GWY-MAC configured across all sites
 - Consistent HSRP group number across sites → consistent GWY-MAC
- Servers can move anywhere and always talk to a local gateway with the same IP/MAC



Host-Mobility and Multi-homing

ETR Updates – Extended Subnets

Null0 host routes indicate the host is “away”

10.2.0.0 /24 is the dyn-EID

6
10.2.0.0/16 – RLOC A, B
10.2.0.2/32 – RLOC C, D

Map-Register
10.2.0.2/32 <C,D>

4
Routing Table:
10.2.0.0/16 – Local
10.2.0.0/24 – Null0
10.2.0.2/32 – Null0

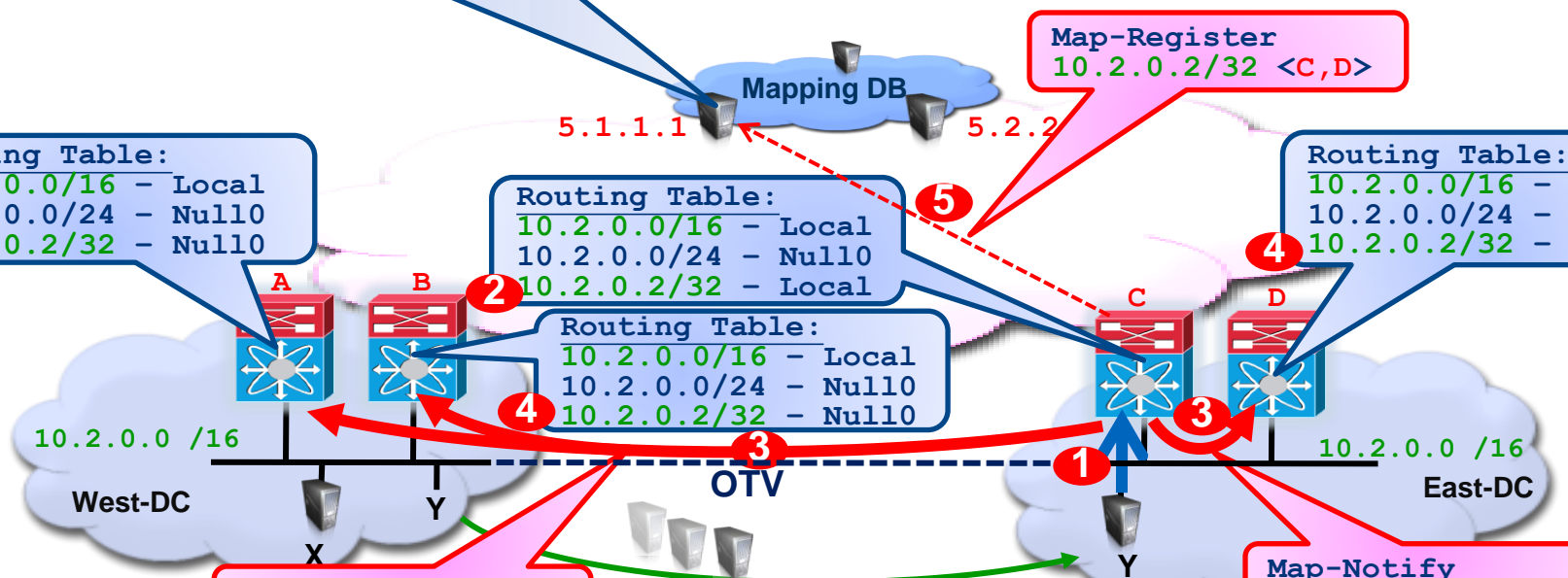
2
Routing Table:
10.2.0.0/16 – Local
10.2.0.0/24 – Null0
10.2.0.2/32 – Local

4
Routing Table:
10.2.0.0/16 – Local
10.2.0.0/24 – Null0
10.2.0.2/32 – Local

Routing Table:
10.2.0.0/16 – Local
10.2.0.0/24 – Null0
10.2.0.2/32 – Null0

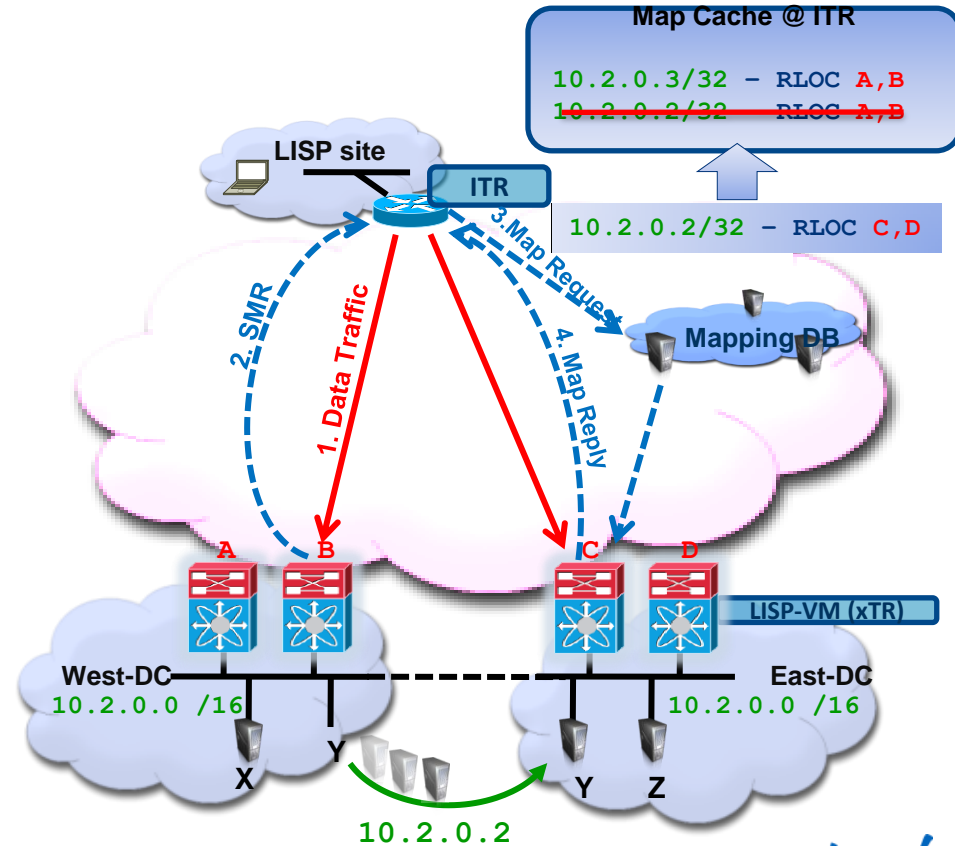
Map-Notify
10.2.0.2/32 <C,D>

Map-Notify
10.2.0.2/32 <C,D>



Refreshing the Map Caches

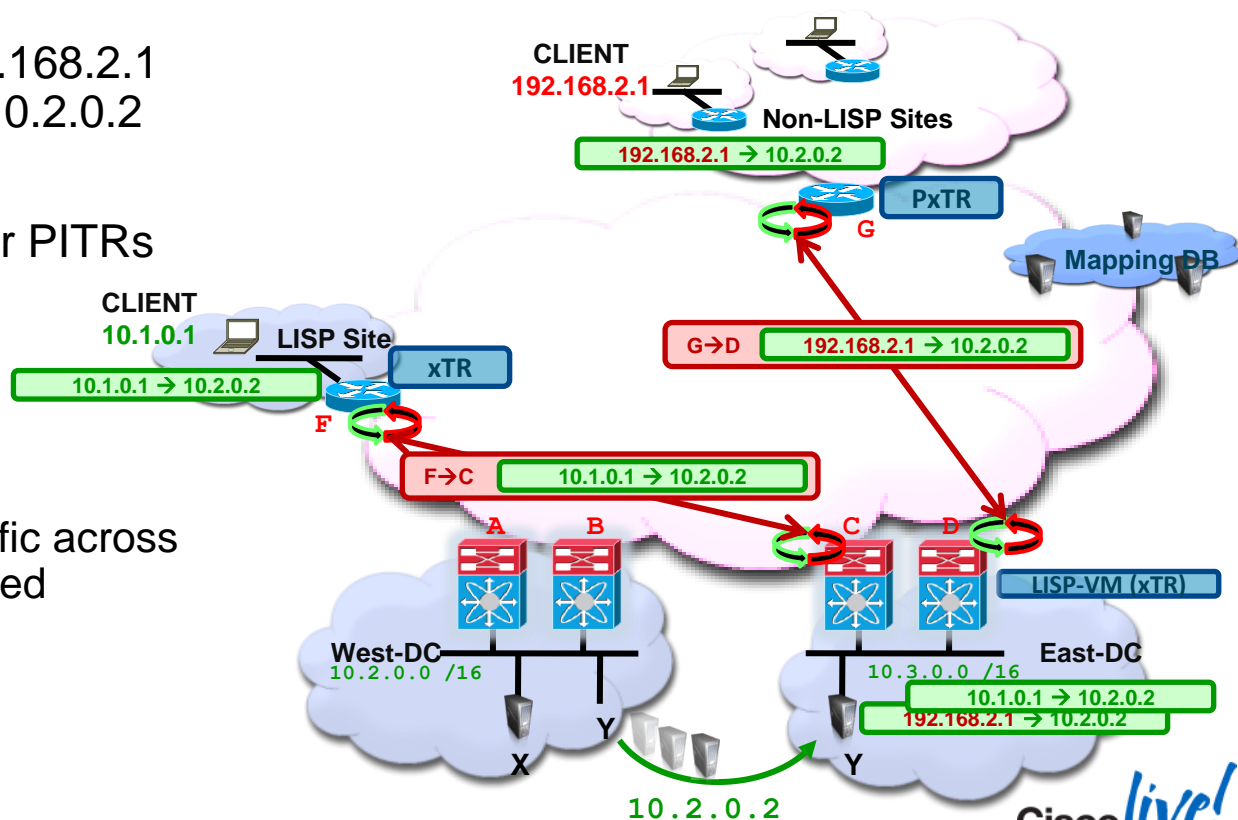
1. ITRs and PITRs with cached mappings continue to send traffic to the old locators
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5. The ITR Map Cache is updated
 - Traffic is now re-directed
 - SMRs are an important integrity measure to avoid unsolicited map responses and spoofing



Off-subnet Client-Server Traffic

All Off-Subnet/Off-Site Traffic Is LISP Encapsulated

- Clients (192.168.0.1 & 192.168.2.1) communicate with Server 10.2.0.2
- Client-server traffic is LISP encapsulated at the ITRs or PITRs
 - Client-to-server:
 - to ETRs C or D
 - Server-to-client:
 - to ETR (F) for LISP sites
 - to PETR (G) for non-LISP sites
- Server-Server off-subnet traffic across sites is also LISP encapsulated

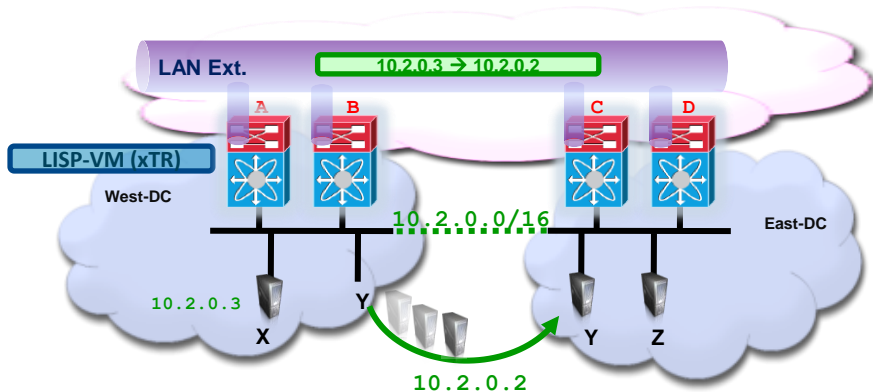


On-subnet Server-Server Traffic

On Subnet Traffic Across L3 Boundaries

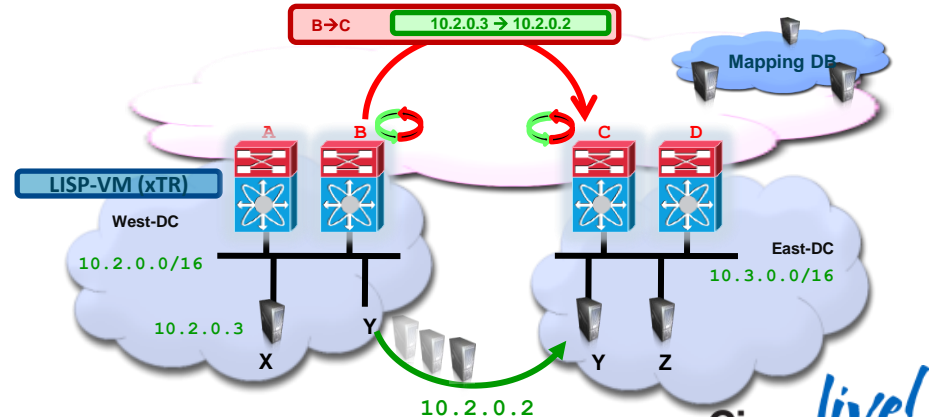
With LAN Extension

- Live moves and cluster member dispersion
- Traffic between X & Y uses the **LAN Extension**
- Link-local-multicast handled by the LAN Extension



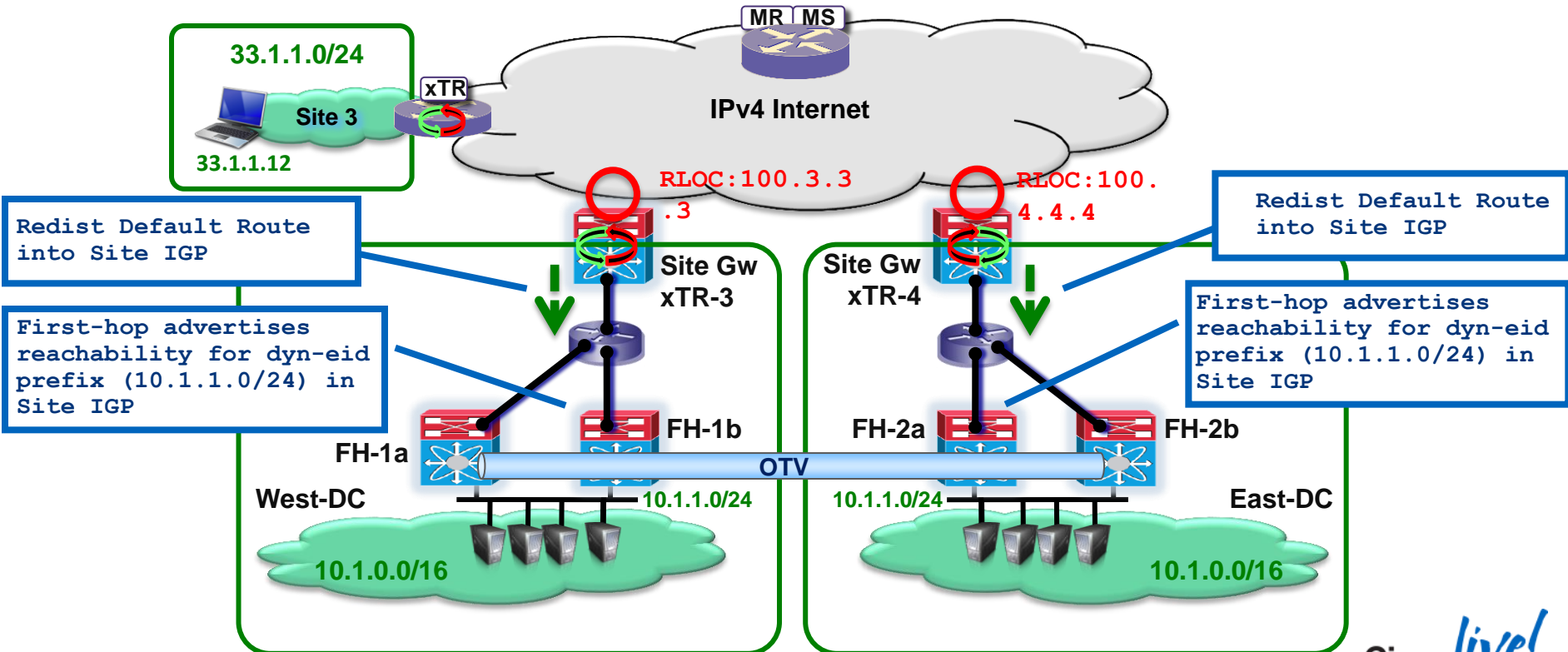
Without LAN Extensions

- Cold moves, no application dispersion
- X- Y traffic is sent to the LISP-VM router & **LISP encapsulated**
- Need LAN extensions for link-local multicast traffic



LISP – New Features

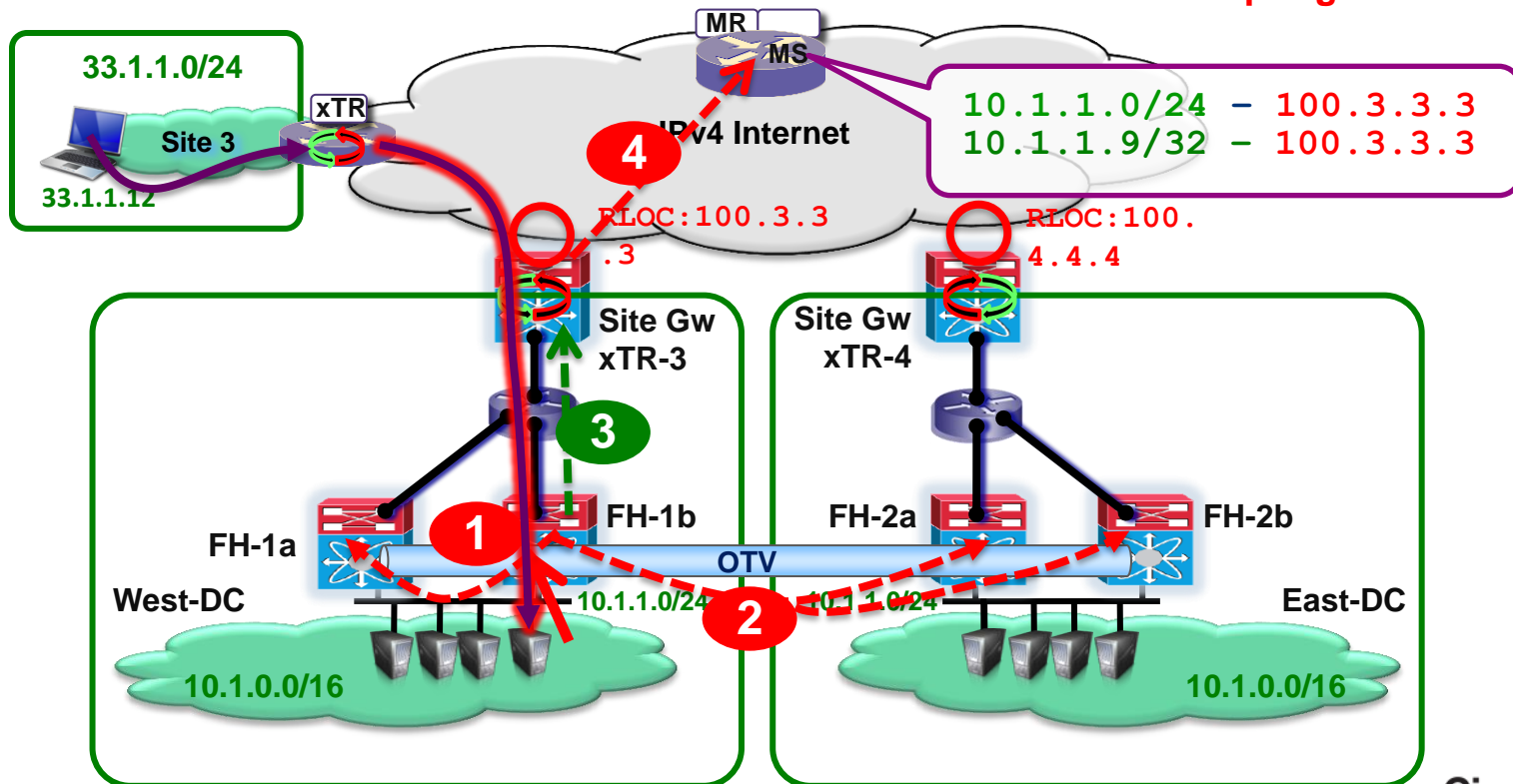
LISP Multi-Hop Mobility – Extended Subnet Mode (ESM)



LISP – New Features

LISP Multi-Hop Mobility – Extended Subnet Mode (ESM)

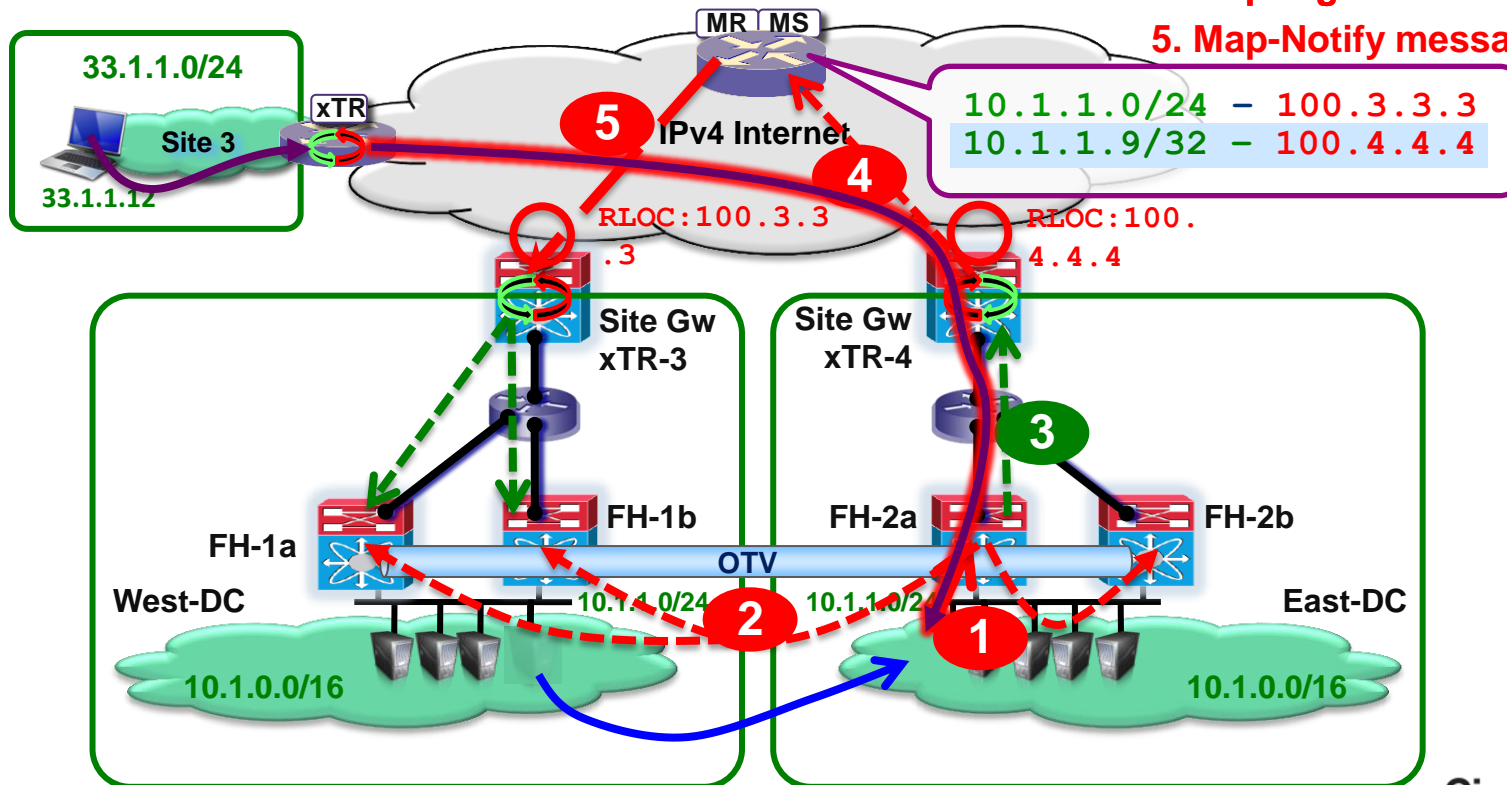
1. Dyn-EID detection
2. Multicast-map-notify
3. Eid-notify to Site Gw xTR
4. Map-register message



LISP – New Features

LISP Multi-Hop Mobility – Extended Subnet Mode (ESM)

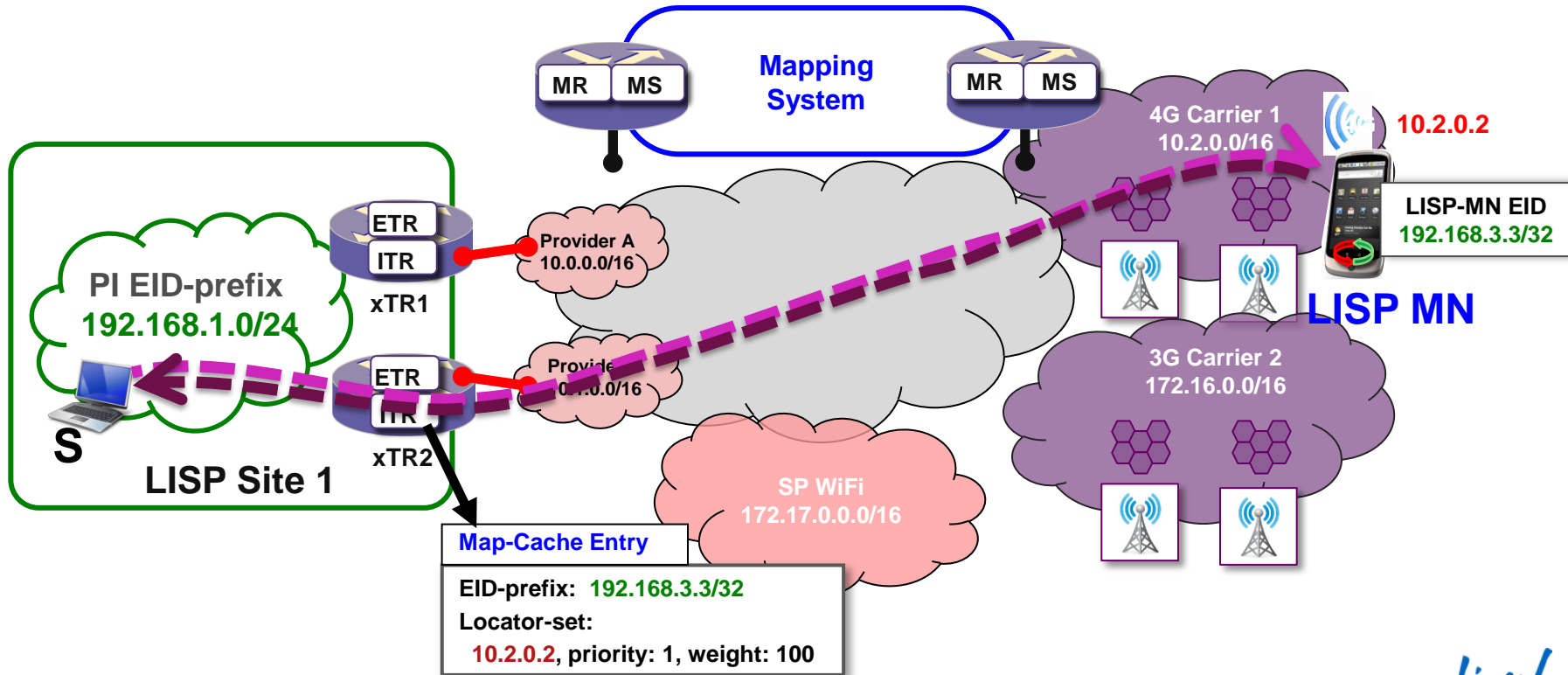
1. Dyn-EID detection
2. Multicast-map-notify
3. Eid-notify to Site Gw xTR
4. Map-register message
5. Map-Notify message



LISP Mobile Node

Session Continuity While Roaming!

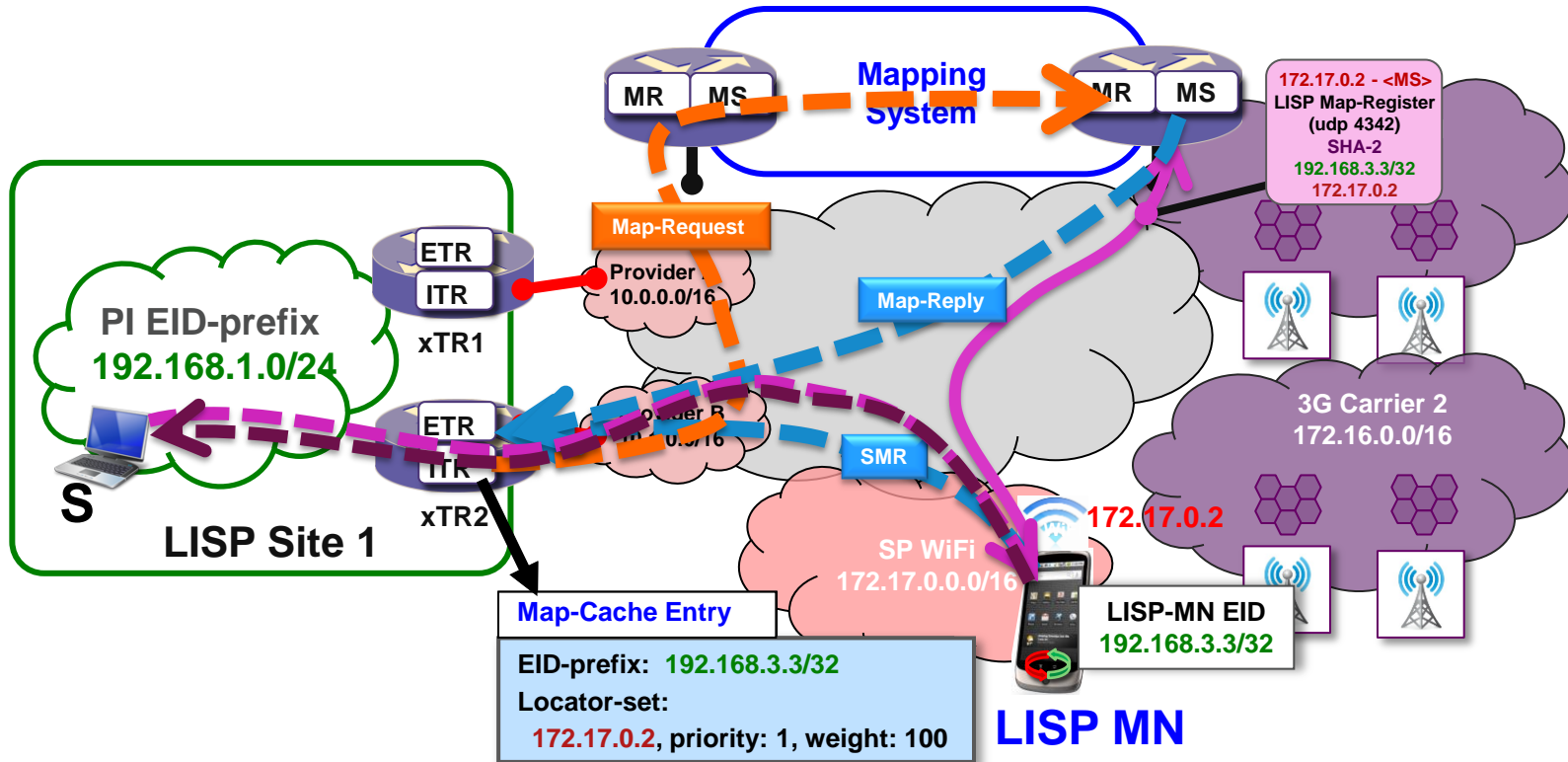
LISP-MN Mobility: Any Network, Anytime, Anywhere...



LISP Mobile Node

Session Continuity While Roaming!

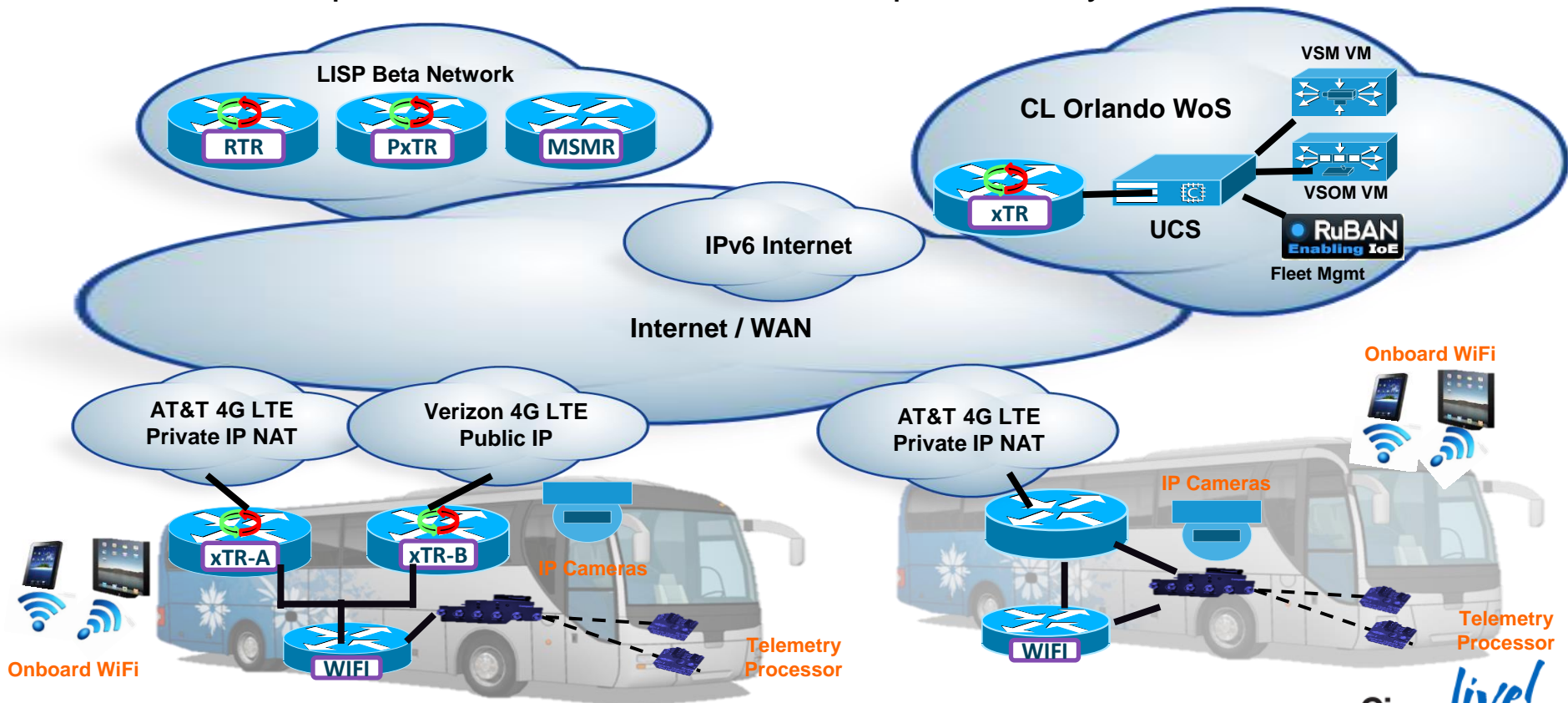
LISP-MN Mobility: Any Network, Anytime, Anywhere...



LISP Use Cases :: LISP Mobility

35 Buses Operational Throughout the Event

Customer Example :: Cisco Live US 2013 Transportation System



LISP Mobile Node Embedded Hardware

Open Source LISP Software

Linksys WRT160NL



Architecture	MIPS Atheros AP81
CPU	400 Mhz Atheros 9130-BC1E
Flash	8 MB cFeon EN25P64
RAM	32 MB Samsung K4H561638J
Ethernet	100 Mbps RTL8306SD
Wireless	Atheros 9102 802.11 b/g/n (integrated)
Serial / JTAG	Yes / Yes
USB	Yes 1x 2.0

Netgear WNDR3700 v2



Architecture	MIPS Atheros AR7161
CPU	680 Mhz Atheros 9130-BC1E
Flash	16 MB Macronix MX25L12845EWI-10G
RAM	64 MB 2 x Nanya NT5DS16M16CS-5T
Ethernet	1 Gbps RTL8366SR
Wireless	Atheros AR9223 802.11b/g/n + Atheros AR9220 802.11a/n
Serial / JTAG	Yes / Yes
USB	Yes 1x 2.0

LISP Mobile Node

LISP-MN Mobility: NAT Traversal Overview/Data Plane...

- Website: <http://lispmob.org/>
- GitHub: <https://github.com/LISPmob/>
- Mailing lists:
 - announce@lispmob.org
 - devel@lispmob.org
 - users@lispmob.org
- IRC: [#lispmob](#) channel on Freenode
- Twitter: <https://twitter.com/LISPmob>



The screenshot shows the LISPmob website homepage. The header features the LISPmob logo and the tagline "an open-source LISP implementation for Linux, Android and OpenWRT". A navigation menu includes links for Home, Downloads, Getting Involved, Documentation, Articles & Media, About, and Contact Us. The main content area is titled "What is LISPmob" and contains a detailed description of the project. Below the text, there are three columns representing different platforms: Linux (with a laptop icon), Android (with a smartphone icon), and OpenWrt (with a wireless router icon). Each column has a "Download" button. At the bottom, there are sections for "You can receive support by:" (with links to documentation, mailing lists, and contact) and "Follow us on social media:" (with icons for Twitter and Facebook).



LISP Status

LISP Status

LISP RFCs and Drafts...

IETF LISP WG: <http://tools.ietf.org/wg/lisp/>

RFCs	
Locator/ID Separation Protocol (LISP) base document	RFC 6830
LISP Map Server	RFC 6833
LISP Interworking	RFC 6832
LISP Multicast	RFC 6831
LISP Internet Groper	RFC 6835
LISP Map Versioning	RFC 6834
LISP+ALT	RFC 6836
LISP MIB	RFC 7052

Draft	Target
LISP Canonical Address Format (draft-ietf-lisp-lcaf-03)	Active Working Group Document
LISP Deployment (draft-ietf-lisp-deployment-11)	Active Working Group Document
LISP SEC (draft-ietf-lisp-sec-05)	Active Working Group Document
LISP DDT (draft-fuller-lisp-ddt-01)	Active Working Group Document
LISP Mobile Node (draft-meyer-lisp-mn-09)	Related Working Group Document
LISP NAT-Traversal (draft-ermagan-lisp-nat-traversal)	Related Working Group Document
LISP GPE (draft-lewis-lisp-gpe)	Related Working Group Document

LISP Status

LISP Deployments - International LISP Beta Network...

■ LISP Community Operated:

- More than **5+** years of operation...
- More than **~600** Sites, **35** countries...

■ Interoperable LISP implementations:

- Cisco

- IOS (ISR, ISRG2, 7200) and IOS-XE (ASR1K)
- Cisco IOS-XR (CRS3, ASR9K (beta))
- Cisco NX-OS (N7K, C200)

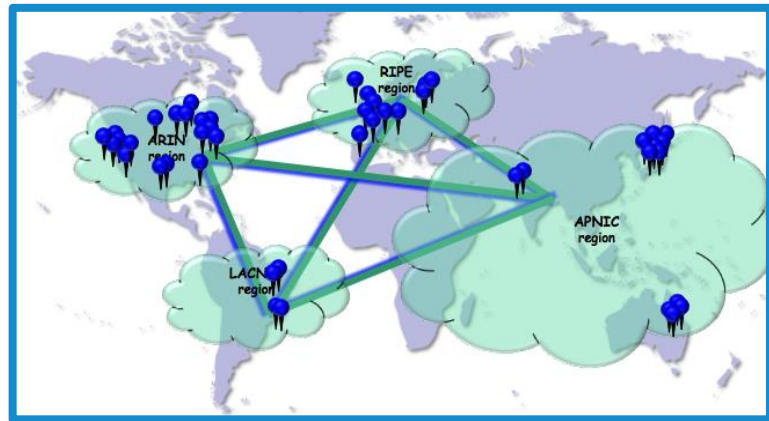
- AVM "FRITZ!Box"



- OpenWrt

- Open Source

- FreeBSD: OpenLISP
- Linux: Aless, LISPmob, OpenWrt
- Android (Gingerbread)



Plus some others... ;-)

LISP Status

LISP Software – Available Releases...

- Cisco Releases (<http://lisp.cisco.com>)

	NX-OS	IOS	IOS-XE	IOS-XR
Software	First Available: 12/2009 Current Main: 6.1(4a) or 6.2(2a)	First Available: 12/2009 Current Main: 15.4(1)T Current Eng: 15.3(3)XB12	First Available: 03/2010 Current Main: 15.3(3)S Current Eng: 15.3(3)S1xb	First Available: 03/2012 Current Main: 4.3.2
Platforms	Nexus 7000 M1-32 linecard	ISR (1800/2800/3800) ISRG2 (800/1900/2900/3900) Catalyst 6500	ASR1K CSR1000V	CRS 3 ASR9k
Features	Roles: ITR/ETR/MS/MR/PITR/PETR AF: EID-v4/v6, RLOC-v4 Virtualisation: Shared/Parallel Mobility: ASM/ESM OTV Multicast: yes	Roles: ITR/ETR/MS/MR/PITR/PETR AF: EID-v4/v6, RLOC-v4/v6 Virtualisation: Shared/Parallel Mobility: ASM/ESM Multicast: roadmap March 2014	Roles: ITR/ETR/MS/MR/PITR/PETR AF: EID-v4/v6, RLOC-v4/v6 Virtualisation: Shared/Parallel Mobility: ASM/ESM OTV Multicast: roadmap Nov 2014	Roles: PITR/PETR AF: EID-v4/v6, RLOC-v4 Virtualisation: Shared/Parallel Mobility: roadmap Multicast: roadmap March 2014



LISP Summary

LISP References

LISP Information

- **LISP Information**

Cisco LISP Site	http://lisp.cisco.com (IPv4 and IPv6)
Cisco LISP Marketing Site	http://www.cisco.com/go/lisp/
LISP Beta Network Site	http://www.lisp4.net or http://www.lisp6.net
LISP DDT Root	http://www.ddt-root.org
IETF LISP Working Group	http://tools.ietf.org/wg/lisp/

- **LISP Mailing Lists**

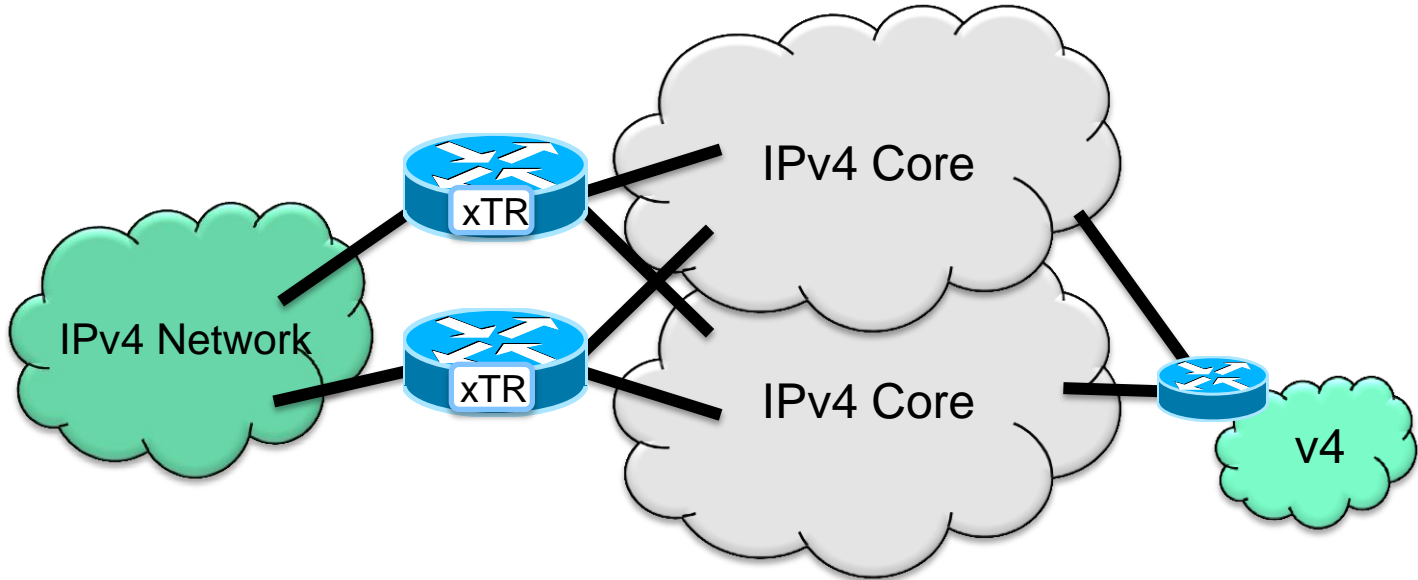
Cisco LISP Questions	lisp-support@cisco.com
IETF LISP Working Group	lisp@ietf.org
LISP Interest (public)	lisp-interest@puck.nether.net
LISPMob Questions	users@lispmob.org



LISP Summary

Part of the LISP Solution Space

1. Multihoming
2. IPv6 Transition
3. Virtualisation/VPN
4. Mobility

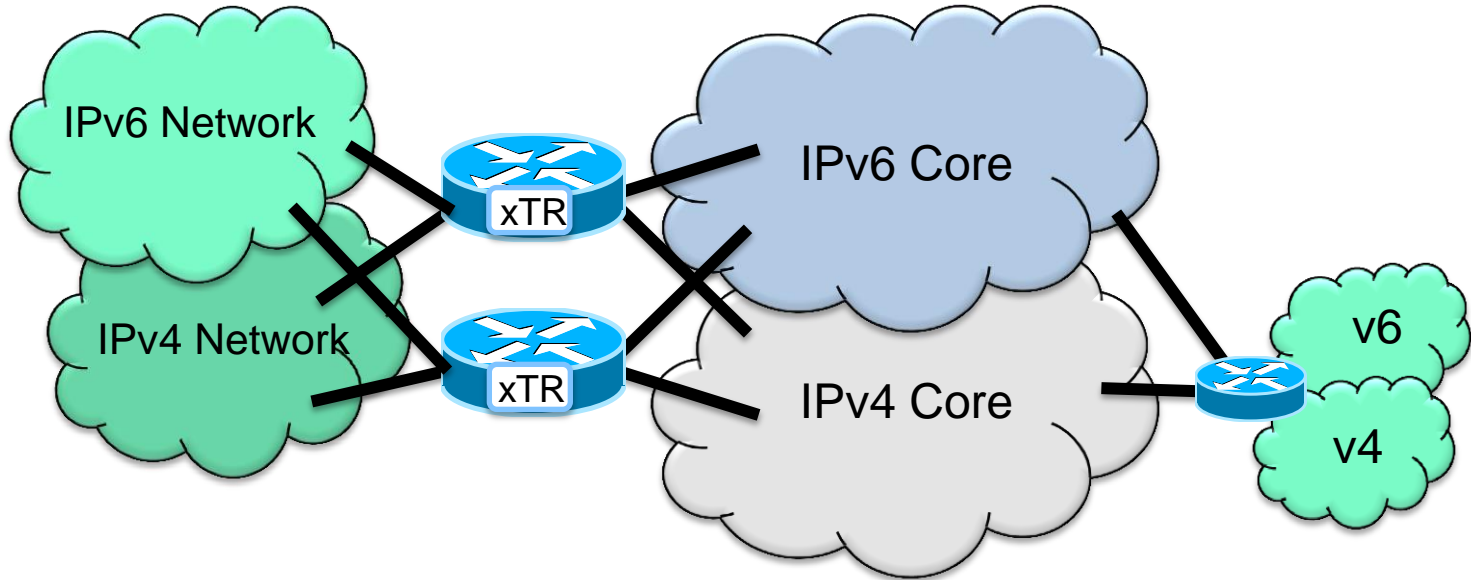


LISP is an Architecture...

LISP Summary

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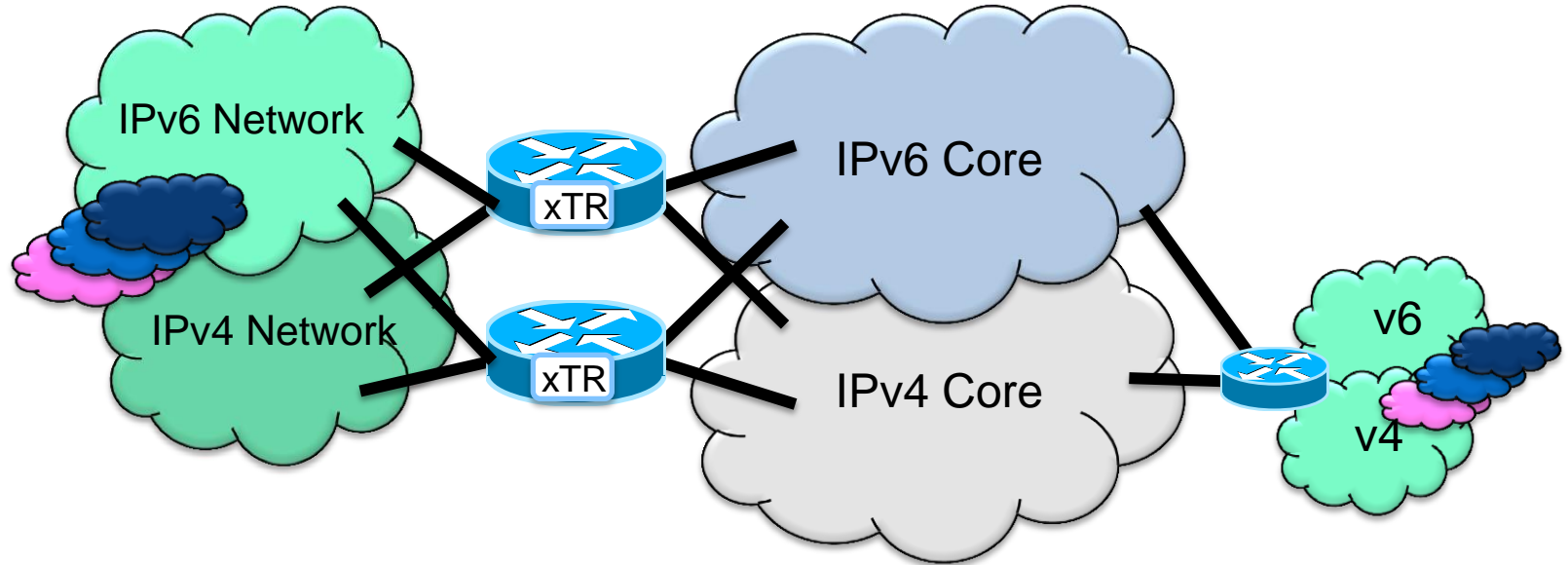


LISP is an Architecture...

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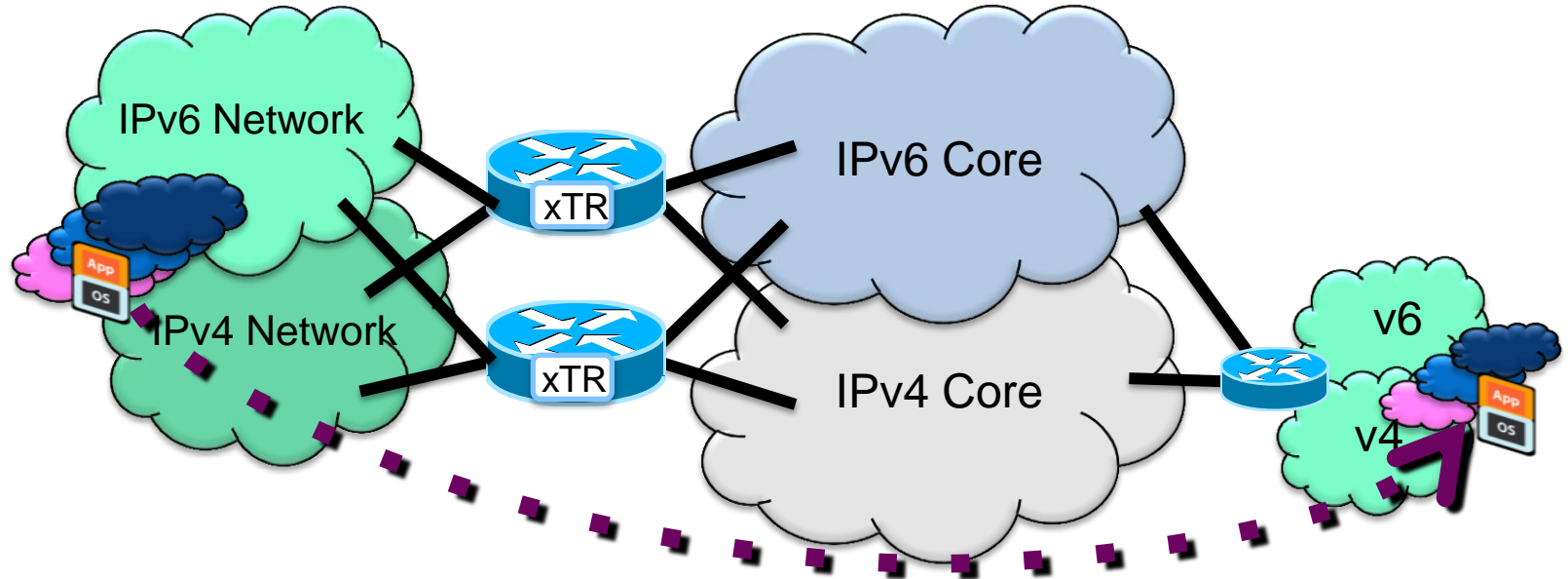


LISP is an Architecture...

LISP Summary

Part of the LISP Solution Space

1. Multihoming
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LISP is an Architecture...

LISP Overview

LISP :: A Routing Architecture – Not a Feature

- Uses pull vs. push routing
 -
 -
- An over-the-top technology
 - Address Family agnostic
 - Incrementally deployable
 - End systems can be unaware of LISP
- Deployment simplicity
 - No host changes
 - Minimal CPE changes
 - Some new core infrastructure components
- LISP use-cases are complimentary
 - Simplified multi-homing with Ingress traffic Engineering; no need for BGP
 - Address Family agnostic support
 - Virtualisation support
 - End-host mobility without renumbering
- Enables IP Number Portability
 - Never change host IP's; No renumbering costs
 - No DNS changes; “name == EID” binding
 - Session survivability
- An Open Standard
 - Being developed in the IETF (RFC 6830-6836)
 - No Cisco Intellectual Property Rights



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